

# Architecture of ML Systems

## 04 Adaptation, Fusion, and JIT

**Matthias Boehm**

Graz University of Technology, Austria  
Computer Science and Biomedical Engineering  
Institute of Interactive Systems and Data Science  
BMVIT endowed chair for Data Management

Last update: Mar 27, 2020

# Announcements/Org

## #1 Video Recording

- Link in **TeachCenter** & **TUbe** (lectures will be public)
- Streaming: <https://tugraz.webex.com/meet/m.boehm>



## #2 Course Administration AMLS

- COVID-19 precautions **March 11 – April 19**
- **Project selection** by **Apr 03** (see **Lecture 02**)  
→ Would a grace period until Apr 25 help?
- Discussion current status project selection



## #3 SystemDS v0.2.0 Release

- Released **Mar 24 2020**,  
with **219 commits**
- Merge into Apache SystemML today

Latest release

v0.2.0

07dd56a

## SystemDS 0.2.0 (March 24, 2020)

corepointer released this 2 days ago · 3 commits to master since this release

### Acknowledgements

Thanks to Enrique Barba Roque, Sebastian Baunsgaard, Matthias Boehm, Mark Dokter, Lukas Erlbacher, Kevin Innerebner, Florian Klezin, Valentin Leutgeb, Arnab Phani, Benjamin Rath, Svetlana Sagadeeva, Afan Secic, Shafaq Siddiqi, Thomas Wedenig, Sebastian Wrede for their support in the creation of the release of SystemDS 0.2.0.

# Agenda

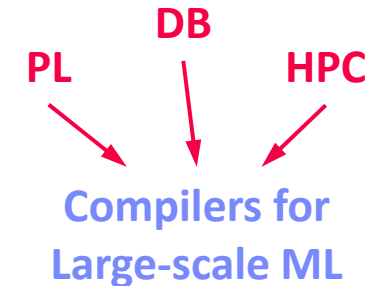
- Motivation and Terminology
- Runtime Adaptation
- Operator Fusion & JIT Compilation
- Discussion Programming Projects

# Motivation and Terminology

# Recap: Linear Algebra Systems

## ■ Comparison Query Optimization

- Rule- and cost-based rewrites and operator ordering
- Physical operator selection and query compilation
- Linear algebra / other ML operators, DAGs, control flow, sparse/dense formats



## ■ #1 Interpretation (operation at-a-time)

- Examples: [R](#), [PyTorch](#), [Morpheus](#) [PVLDB'17]

## ■ #2 Lazy Expression Compilation (DAG at-a-time)

- Examples: [RIOT](#) [CIDR'09], [TensorFlow](#) [OSDI'16], [Mahout Samsara](#) [MLSystems'16]
- Examples w/ control structures: [Weld](#) [CIDR'17], [OptiML](#) [ICML'11], [Emma](#) [SIGMOD'15]

## ■ #3 Program Compilation (entire program)

- Examples: [SystemML](#) [ICDE'11/PVLDB'16], [Julia](#), [Cumulon](#) [SIGMOD'13], [Tupeware](#) [PVLDB'15]

## Optimization Scope

```

1: X = read($1); # n x m matrix
2: y = read($2); # n x 1 vector
3: maxi = 50; lambda = 0.001;
4: intercept = $3;
5: ...
6: r = -(t(X) %*% y);
7: norm_r2 = sum(r * r); p = -r;
8: w = matrix(0, ncol(X), 1); i = 0;
9: while(i < maxi & norm_r2 > norm_r2_trgt)
10: {
11:   q = (t(X) %*% X %*% p) + lambda * p;
12:   alpha = norm_r2 / sum(p * q);
13:   w = w + alpha * p;
14:   old_norm_r2 = norm_r2;
15:   r = r + alpha * q;
16:   norm_r2 = sum(r * r);
17:   beta = norm_r2 / old_norm_r2;
18:   p = -r + beta * p; i = i + 1;
19: }
20: write(w, $4, format="text");

```

# Major Compilation/Runtime Challenges

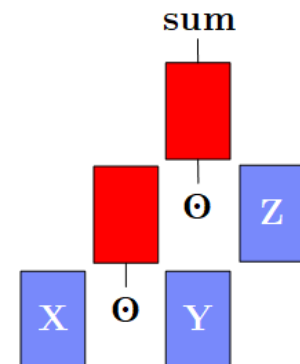
## ■ #1 Unknown/Changing Sizes

- Sizes inference crucial for cost-estimation and validity constraints (e.g., rewrites)
- **Tradeoff:** optimization scope vs size inference effort
- **Challenge:** Unknowns → conservative fallback plans

$Y = \text{foo}(X)$   
 $Z = Y[I_x, ]$   
**#  $\text{nrow}(Z)?$**

## ■ #2 Operator Runtime Overhead

- Operators great for **programmability**, size inference, simple compilation, and **efficient kernel implementations** (sparse, dense, compressed)
- **Tradeoff:** general-purpose vs specialization
- **Challenges:** intermediates, parallelization, complexity of operator combinations



# Terminology Ahead-of-Time / Just-in-Time

## ■ Ahead-of-Time Compilation

- Originating from compiled languages like C, C++
- #1 **Program compilation** at different abstraction levels
- #2 **Inference program compilation** & packaging



## ■ Just-In-Time Compilation (at runtime for specific data/HW)

- Originating from JIT-compiled languages like Java, C#
- #1 **Lazy expression evaluation** + optimization
- #2 Program/function compilation **with recompilation**



## ■ Excursus: Java JIT

- #1 Start w/ Java bytecode interpretation by JVM → **fast startup**
- #2 **Tiered JIT compile** (cold, warm, hot, very hot, scorching) → **performance**
- Trace statistics (frequency, time) at method granularity
- Note: -XX:+PrintCompilation

PL

# Terminology Runtime Adaptation & JIT

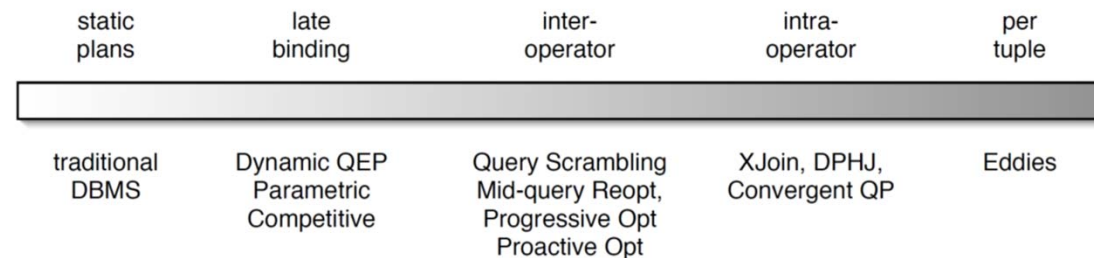
## ■ Excursus: Adaptive Query Processing

[Amol Deshpande, Joseph M. Hellerstein, Shankar Raman: Adaptive query processing: why, how, when, what next. **SIGMOD 2006**]



### ■ Spectrum of Adaptivity

DB



## ■ Excursus: Query Execution Strategies

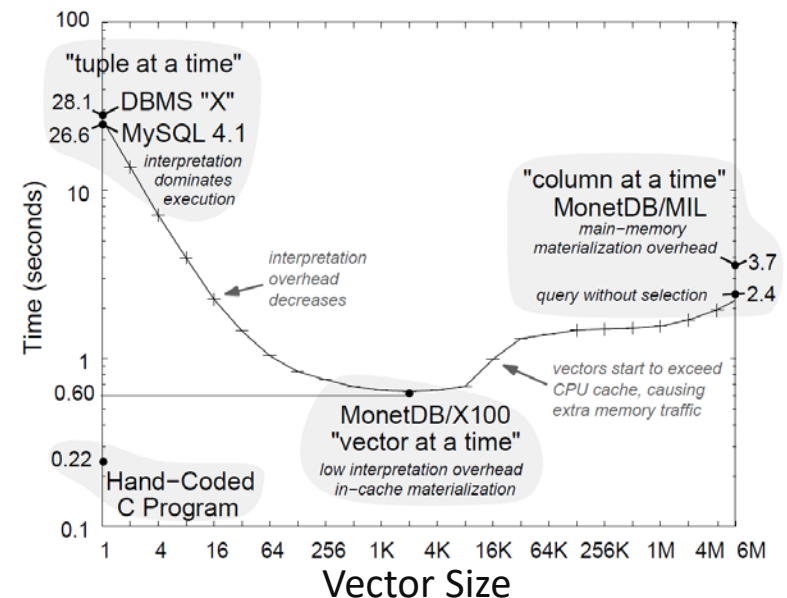
- #1 Volcano Iterator Model
- #2 Materialized Intermediates
- #3 Vectorized (Batched) Execution
- #4 Query Compilation
- Similar: Loop fusion, fission, tiling

DB

HPC



[Peter A. Boncz, Marcin Zukowski, Niels Nes: MonetDB/X100: Hyper-Pipelining Query Execution. **CIDR 2005**]





# Runtime Adaptation

ML Systems w/ Optimizing Compiler



# Issues of Unknown or Changing Sizes

## ■ Problem of unknown/changing sizes

- **Unknown or changing** sizes and sparsity of intermediates  
These unknowns lead to very **conservative fallback plans** (distributed ops)

## ■ #1 Control Flow

- Branches and loops
- Complex function call graphs
- User-Defined Functions

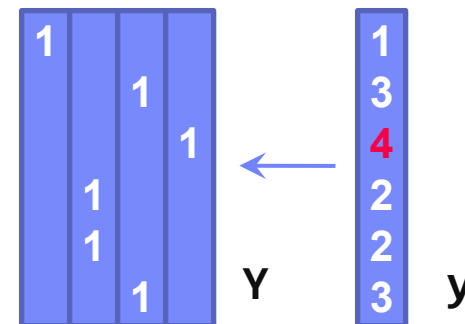
```
X = read('/tmp/X.csv');
if( intercept )
  X = cbind(X, matrix(1,nrow(X),1));
Z = foo(X) + X; # size of + and Z?
```

## ■ #2 Data-Dependencies

- Data-dependent operators  
(e.g., table, rmEmpty, aggregate)
- Computed size expressions

```
d = dout[, (t-2)*M+1:(t-1)*M];
cur_Q = matrix (0, 1, 2*ncur);
cur_S = matrix (0, 1, ncur*dist);
```

```
Y = table(seq(1,nrow(X)), y);
grad = t(X) %*% (P - Y);
```



**Ex.: Multinomial  
Logistic Regression**

# Issues of Unknown or Changing Sizes, cont.

## ■ #3 Changing Dims and Sparsity

- Iterative feature selection workloads
- Changing dimensions or sparsity
- ➔ Same code with different data

## ■ #4 API Limitations

- Precompiled scripts/programs (inputs unavailable)

## ■ (#5 Compiler Limitations)

➔ **Dynamic recompilation techniques** as robust fallback strategy

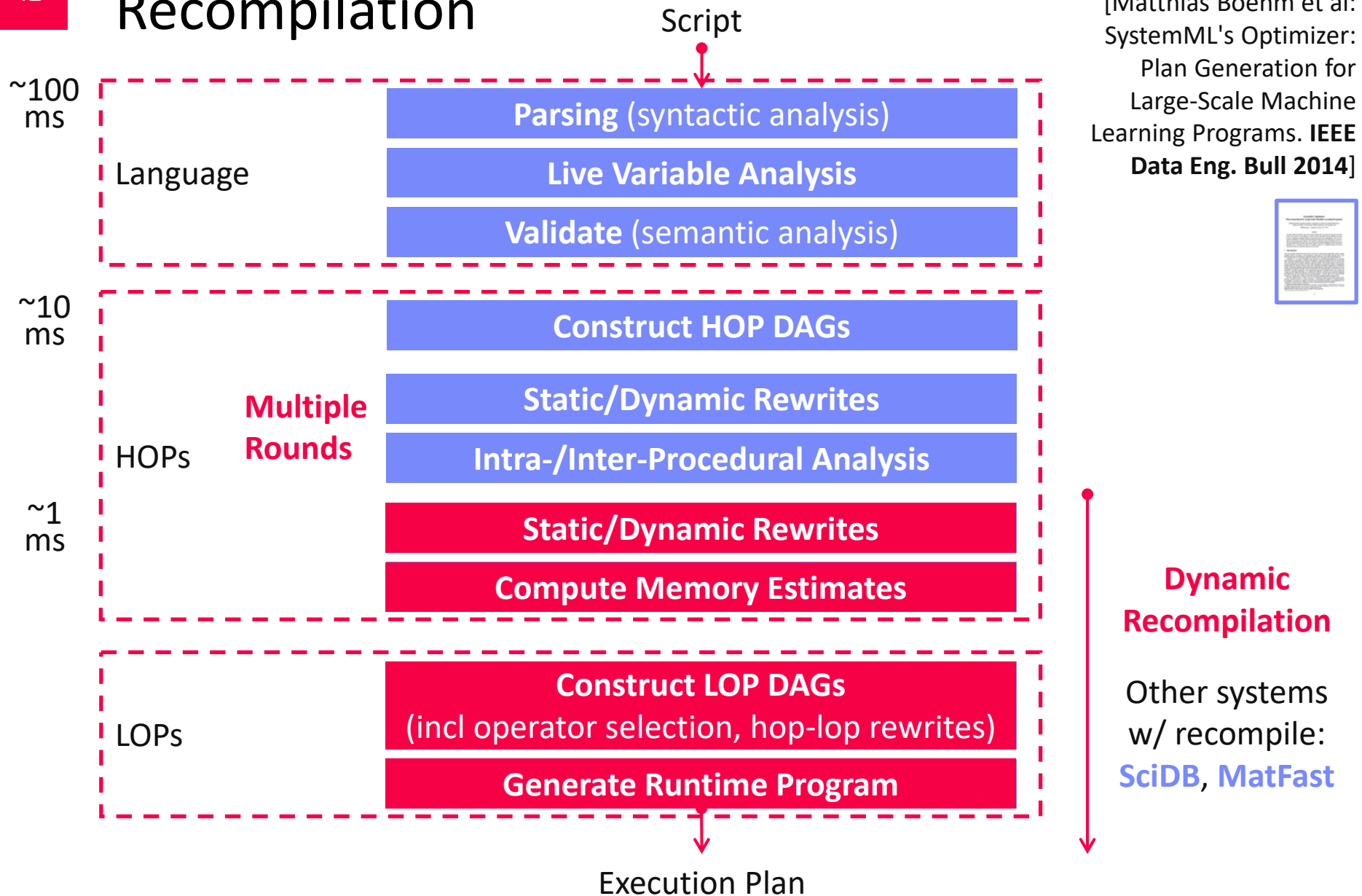
- Shares goals and challenges with adaptive query processing
- However, ML domain-specific techniques and rewrites

### Ex: Stepwise LinReg

```
while( continue ) {  
    parfor( i in 1:n ) {  
        if( !fixed[1,i] ) {  
            Xi = cbind(Xg, X[,i])  
            B[,i] = lm(Xi,y)  
        }  
    }  
    # add best to Xg (AIC)  
}
```

12

# Recompilation

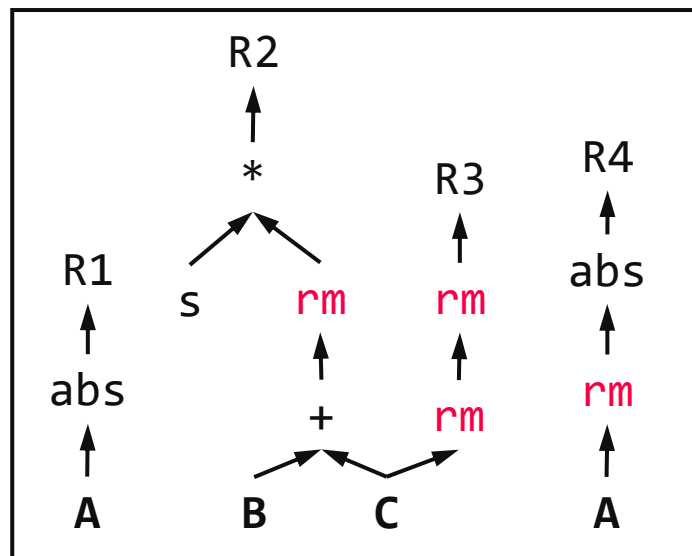


# Dynamic Recompilation

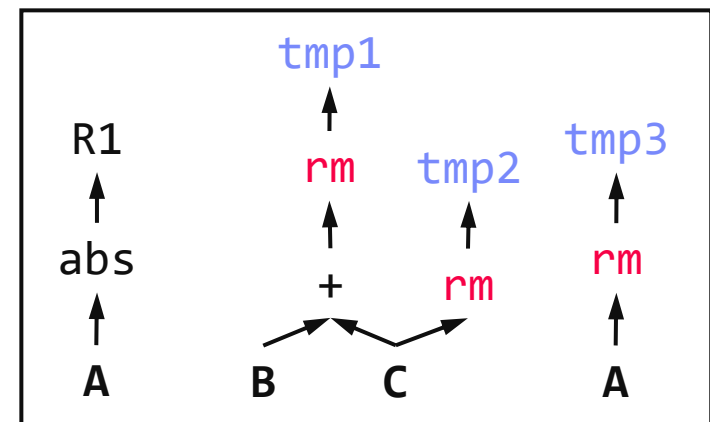
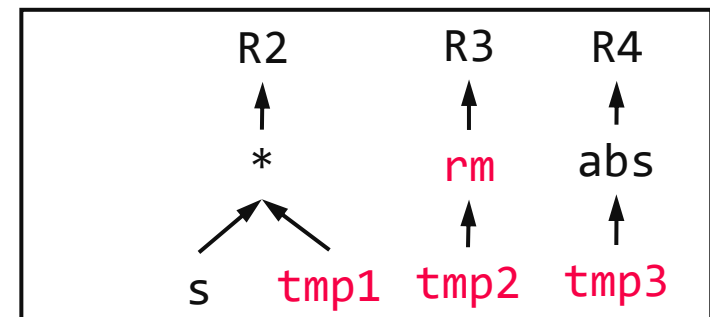
## ■ Compile-time Decisions

- **Split HOP DAGs for recompilation:** prevent unknowns but keep DAGs as large as possible; split after reads w/ unknown sizes and specific operators
- **Mark HOP DAGs for recompilation:** Spark due to unknown sizes / sparsity

Control flow → statement blocks  
 → **initial recompilation granularity**



➡  
 (recursive  
 rewrite)

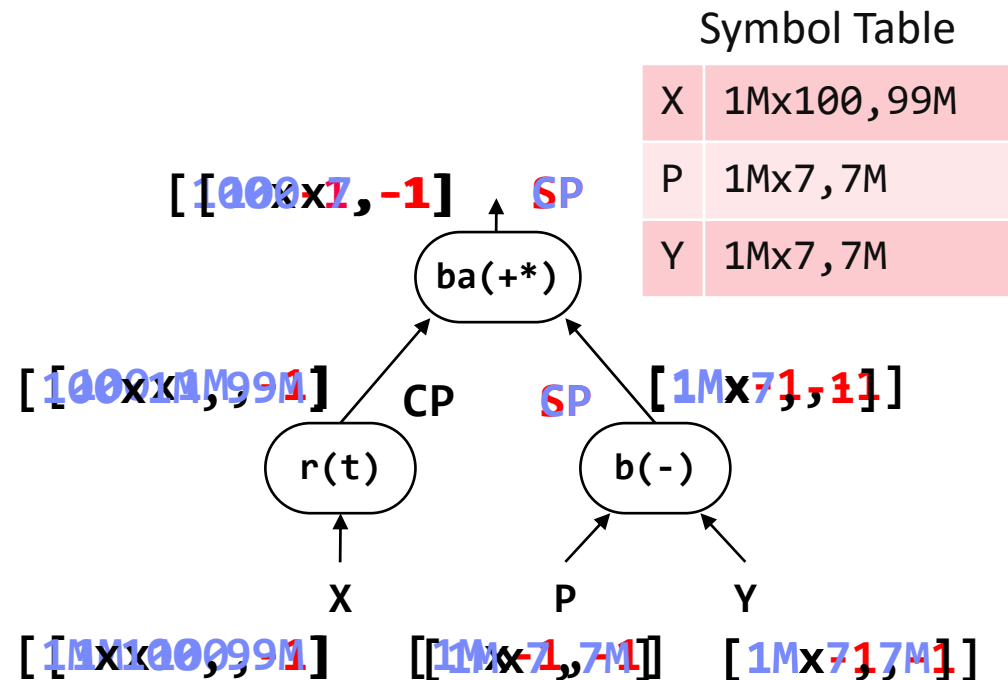


rm .. removeEmpty(X, [margin="rows",select=I])

# Dynamic Recompilation, cont.

- **Dynamic Recompilation at Runtime** on recompilation hooks  
(last level program blocks, predicates, recompile once functions)

- Deep Copy DAG
- Replace Literals
- Update DAG Statistics
- Dynamic Rewrites
- Recompute Memory Estimates
- [Codegen]
- Generate Runtime Instructions



# Dynamic Recompilation, cont.

## ■ Recompile Once Functions

- Unknowns due to inconsistent or unknown call size information
- IPA marks functions as “recompile once”, if it contains loops
- **Recompile the entire function on entry + disable unnecessary recompile**

```
foo = function(Matrix[Double] A)
  # recompiled w/ size of A
  return (Matrix[Double] C)
{
  C = rand(nrow(A),1) + A;
  while(...)
    C = C / rowSums(C) * s
}
```

## ■ Recompile parfor Loops

- Unknown sizes and iterations
- **Recompile parfor loop on entry + disable unnecessary recompile**
- Create independent DAGs for individual parfor workers

```
while( continue ) {
  parfor( i in 1:n ) {
    if( !fixed[1,i] ) {
      Xi = cbind(Xg, X[,i])
      B[,i] = lm(Xi,y)
    }
  }
  # add best to Xg (AIC)
}
```

# Operator Fusion & JIT Compilation (aka Code Generation)

Many State-of-the-Art ML Systems,  
especially for DNNs and numerical computation









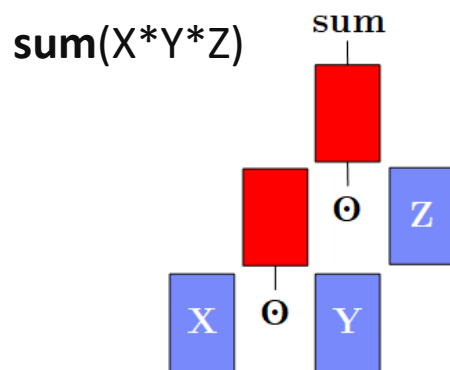
# Motivation: Fusion

[Matthias Boehm et al.: On Optimizing Operator Fusion Plans for Large-Scale ML in SystemML. **PVLDB 2018**]

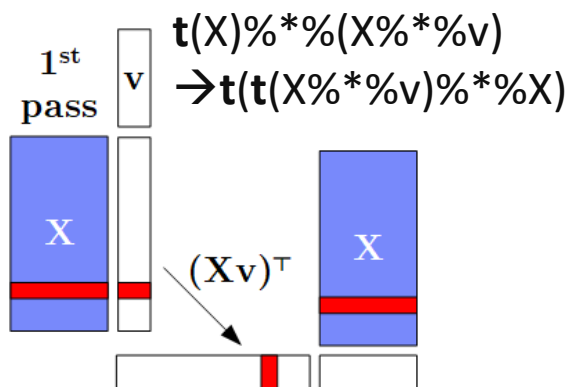


- **Data Flow Graphs (better data access)**
  - DAGs of linear algebra (LA) operations and statistical functions
  - Materialized intermediates → **ubiquitous fusion opportunities**

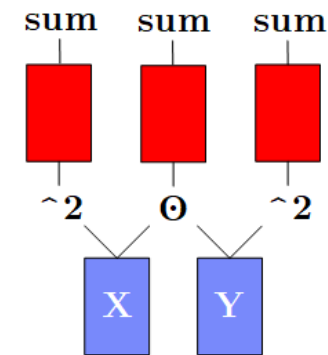
a) Intermediates



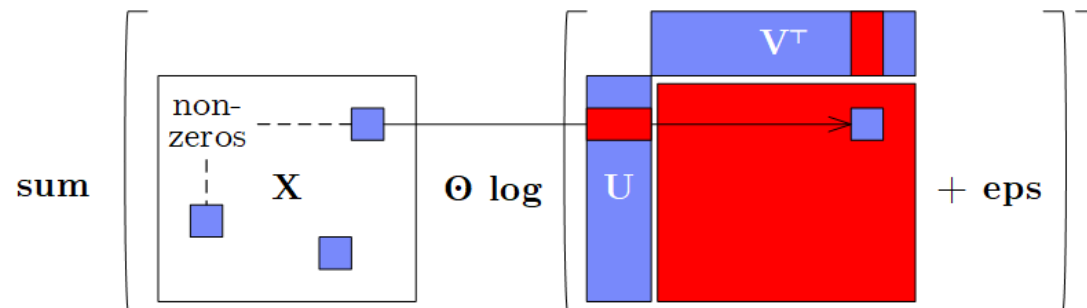
b) Single-Pass



c) Multi-Aggregates

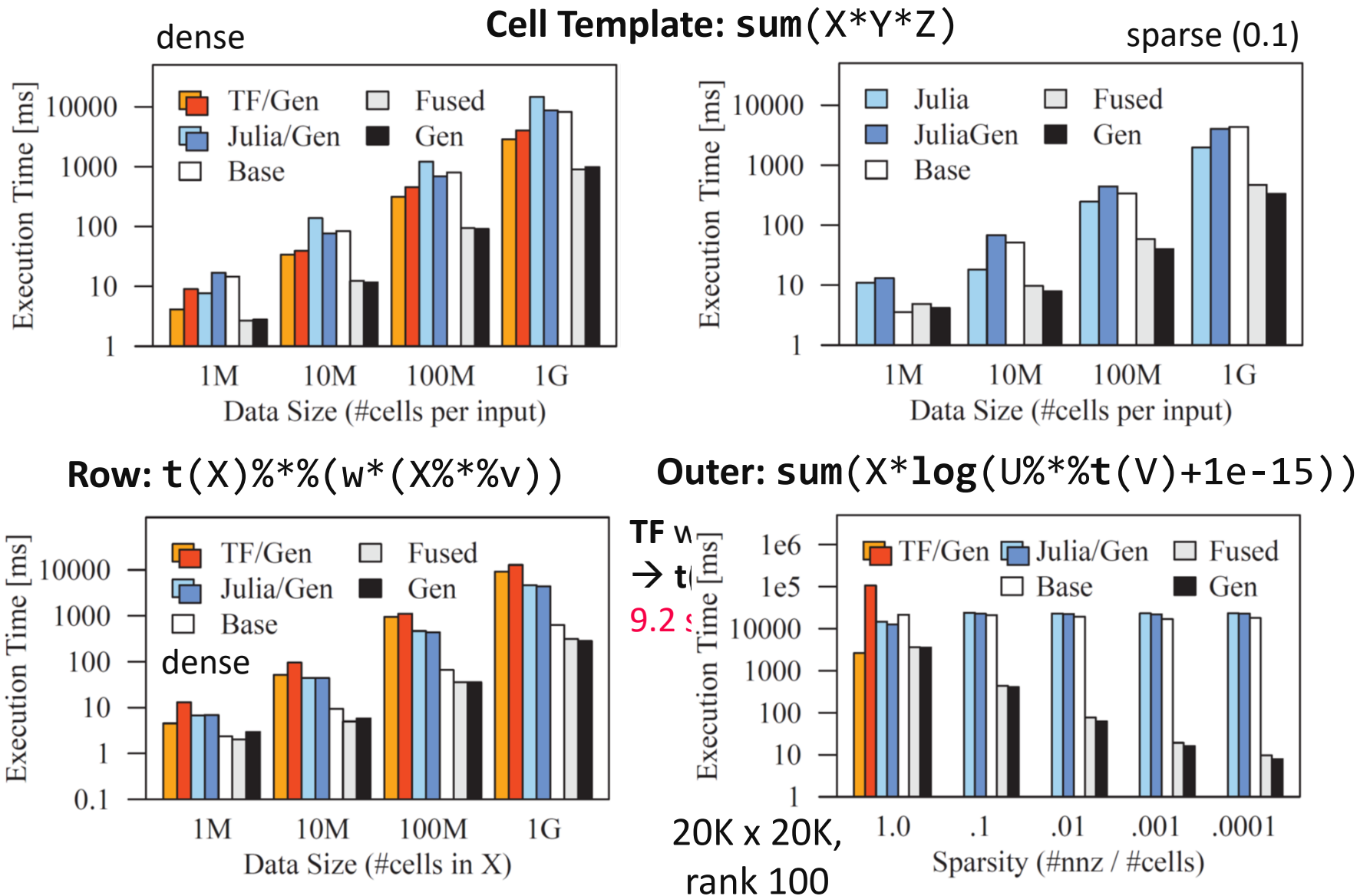


d) Sparsity Exploitation



# Motivation: Fusion, cont.

**Beware:** SystemML 1.0,  
Julia 0.6.2, TensorFlow 1.5



# Motivation: Just-In-Time Compilation

## ■ Operator Kernels (**better code**)

- Specialization opportunities: data types, shapes, and operator graphs
- Heterogeneous hardware: CPUs, GPUs, FPGAs, ASICs x architectures

## ■ #1 CPU Architecture

- Specialize to available instructions sets
- Register allocation and assignment, etc

**Examples:** x86-64,  
sparc, amd64, arm, ppc

## ■ #2 Heterogeneous Hardware

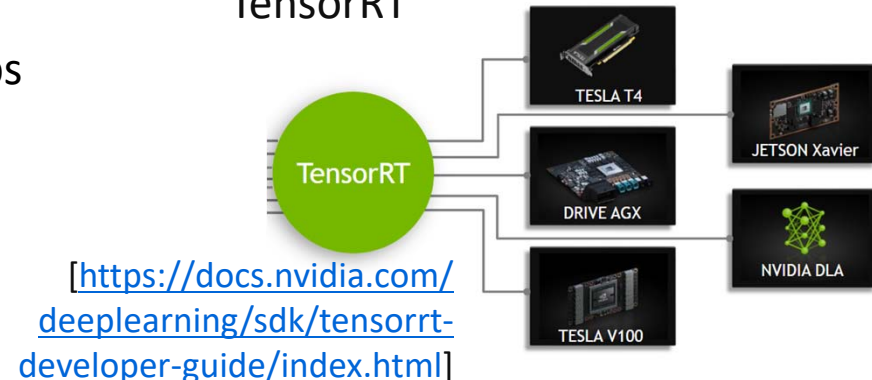
- JIT compilation for custom-build ASICs with HW support for ML ops
- Different architectures of devices

## ■ #3 Custom ML Program

- Operator graphs and sizes

**Example: NVIDIA**  
TensorRT

**GPU Platforms**

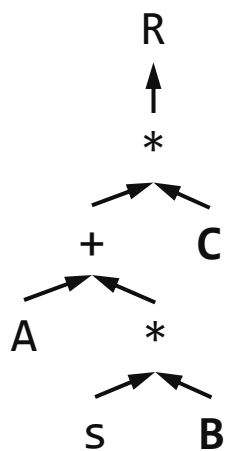


# Operator Fusion Overview

## Related Research Areas

- DB: **query compilation**
- HPC: **loop fusion, tiling, and distribution** (**NP complete**)
- ML: **operator fusion** (dependencies given by data flow graph)

## Example Operator Fusion



```
for( i in 1:n )
  tmp1[i,1] = s * B[i,1];
for( i in 1:n )
  tmp2[i,1] = A[i,1] + tmp1[i,1];
for( i in 1:n )
  R[i,1] = tmp2[i,1] * C[i,1];
```



```
for( i in 1:n )
  R[i,1] = (A[i,i] + s*B[i,1]) * C[i,1];
```

## Memory Bandwidth:

L1 core: 1TB/s

L3 socket: 400GB/s

Mem: 100 GB/s

[\[https://software.intel.com/en-us/articles/memory-performance-in-a-nutshell\]](https://software.intel.com/en-us/articles/memory-performance-in-a-nutshell)

# Evolution of Operator Fusion in ML Systems

## ■ 1<sup>st</sup> Gen: **Handwritten** Fused Operators

- [BLAS (since 1979): e.g.,  $\alpha * X + Y \rightarrow \text{AXPY}$ ]
- Rewrites: e.g.,  $A+B+C \rightarrow \text{AddN}(A, B, C)$ ,  
 $\mathbf{t}(X) \%* \% (w * (X \%* \% v)) \rightarrow \text{MMCHAIN}$
- Sparsity exploiting fused ops:  
e.g.,  $\text{sum}(X * \log(U \%* \% \mathbf{t}(V) + \text{eps}))$

[Arash Ashari: On optimizing machine learning workloads via kernel fusion. **PPOPP 2015**]



[Matthias Boehm: SystemML: Declarative Machine Learning on Spark. **PVLDB 2016**]



## ■ 2<sup>nd</sup> Gen: **Fusion Heuristics**

- Automatic operator fusion via elementary ops
- Heuristics for replacing sub-DAGs w/ fused ops

[Tarek Elgamal et al: SPOOF: Sum-Product Optimization and Operator Fusion for Large-Scale Machine Learning. **CIDR 2017**]



## ■ 3<sup>rd</sup> Gen: **Optimized Fusion Plans**

- Greedy/exact fusion plan (sub-DAG) selection
- [Greedy/evolutionary kernel implementations]

[Matthias Boehm et al.: On Optimizing Operator Fusion Plans for Large-Scale ML in SystemML. **PVLDB 2018**]



# Automatic Operator Fusion System Landscape

JIT

System	Year	Approach	Sparse	Distr.	Optimization
BTO	2009	Loop Fusion	No	No	k-Greedy, cost-based
Tupeware	2015	Loop Fusion	No	Yes	Heuristic
Kasen	2016	Templates	(Yes)	Yes	Greedy, cost-based
<b>SystemML</b>	<b>2017</b>	<b>Templates</b>	<b>Yes</b>	<b>Yes</b>	<b>Exact, cost-based</b>
Weld	2017	Templates	(Yes)	Yes	Heuristic
Taco	2017	Loop Fusion	Yes	No	Manuel
Julia	2017	Loop Fusion	Yes	No	Manuel
Tensorflow XLA	2017	Loop Fusion	No	No	Manuel/Heuristic
Tensor Comprehensions	2018	Loop Fusion	No	No	Evolutionary, cost-based
TVM	2018	Loop Fusion	No	No	ML/cost-based
PyTorch	2019	Loop Fusion	No	No	Manual/Heuristic
JAX	2019	N/A	No	No	See TF XLA

# A Case for Optimizing Fusion Plans



- **Problem:** Fusion heuristics → **poor plans** for complex DAGs (cost/structure), sparsity exploitation, and local/distributed operations
- **Goal:** **Principled approach for optimizing fusion plans**

$$C = A + s * B$$

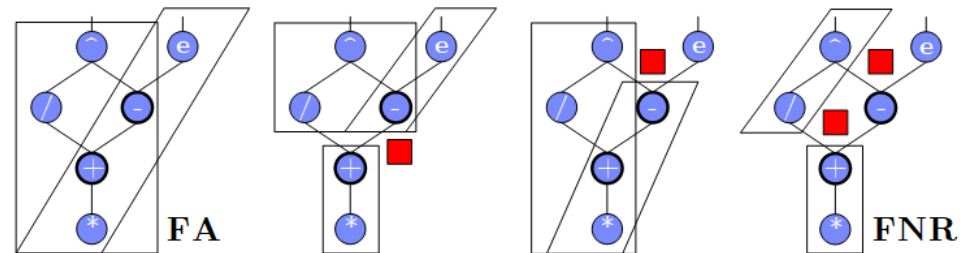
$$D = (C/2)^{(C-1)}$$

$$E = \exp(C-1)$$

- **#1 Materialization Points**  
(e.g., for multiple consumers)

- **#2 Sparsity Exploitation**  
(and ordering of sparse inputs)

- **#3 Decisions on Fusion Patterns**  
(e.g., template types)
- **#4 Constraints**  
(e.g., memory budget and block sizes)

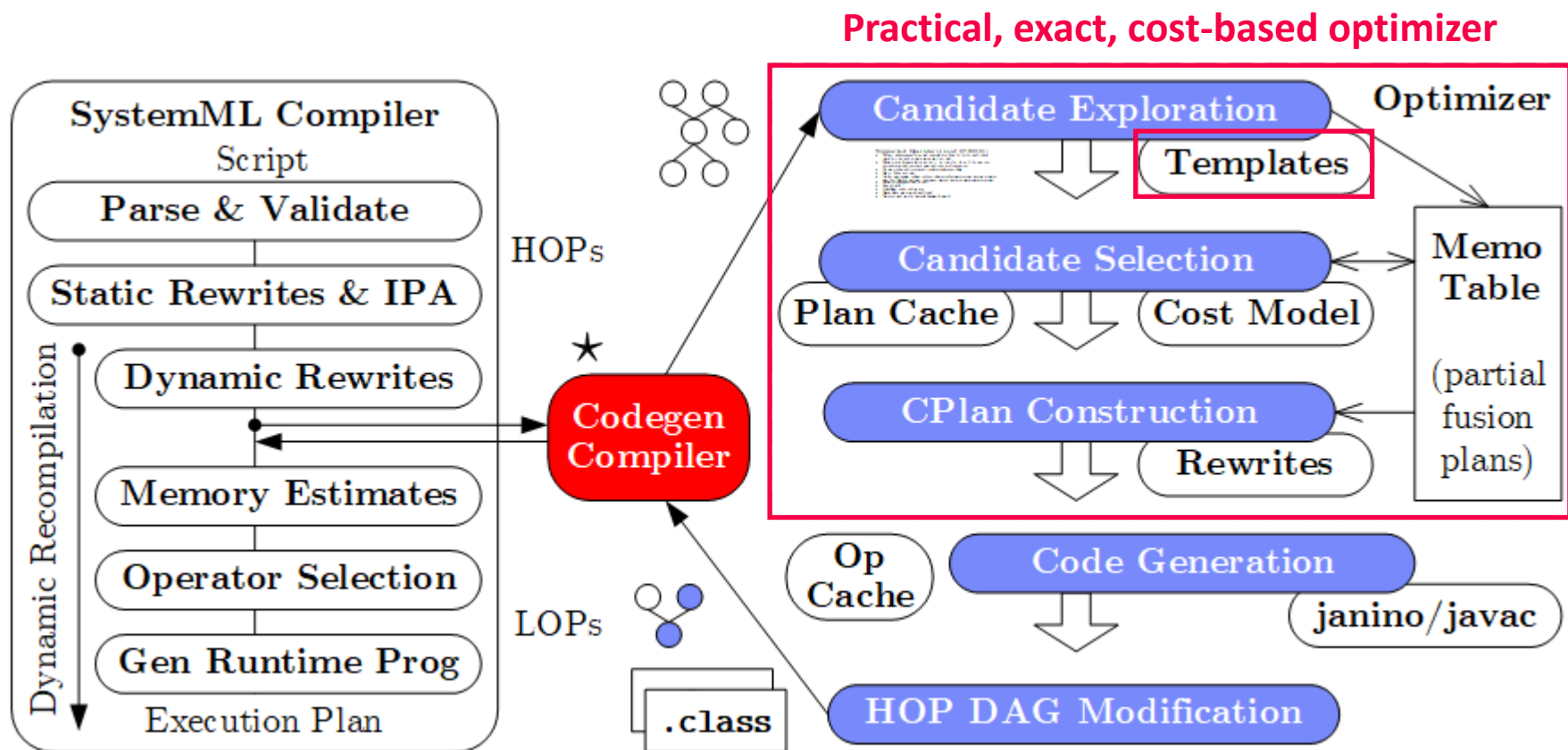


$$Y + \boxed{X * (U \% \% t(V))}$$

sparse-safe over X

→ Search Space that  
requires optimization

# System Architecture (Compiler & Codegen Architecture)



- CPlan representation/construction and codegen similar in TF XLA (HLO primitives, pre-clustering of nodes, caching, LLVM codegen)
- **Templates:** **Cell**, **Row**, **M**Agg, **O**uter w/ different data bindings



# Codegen Example L2SVM (Cell/MAgg)

## ■ L2SVM Inner Loop

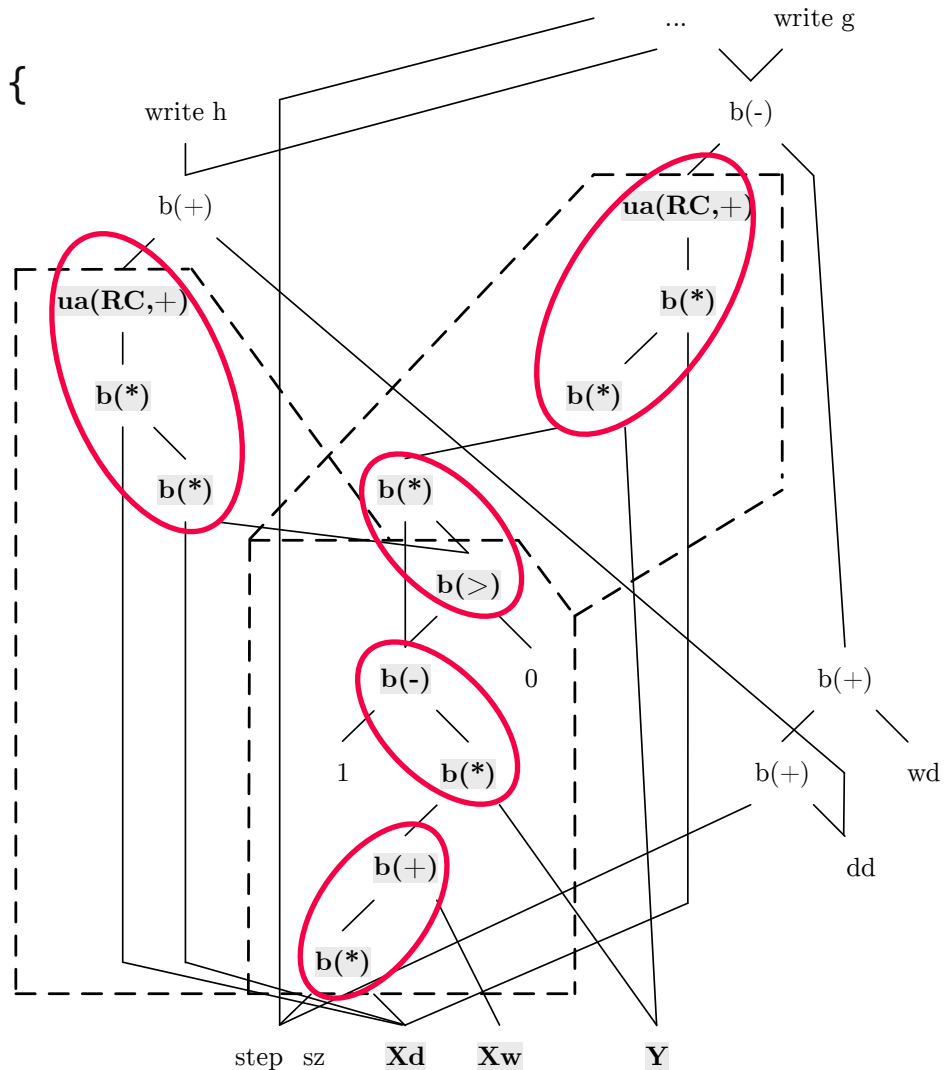
```

1: while(continueOuter & iter < maxi) {
2:   #...
3:   while(continueInner) {
4:     out = 1-Y* (Xw+step_sz*Xd);
5:     sv = (out > 0);
6:     out = out * sv;
7:     g = wd + step_sz*dd
        - sum(out * Y * Xd);
8:     h = dd + sum(Xd * sv * Xd);
9:     step_sz = step_sz - g/h;
10:  }} ...

```

## ■ # of Vector Intermediates

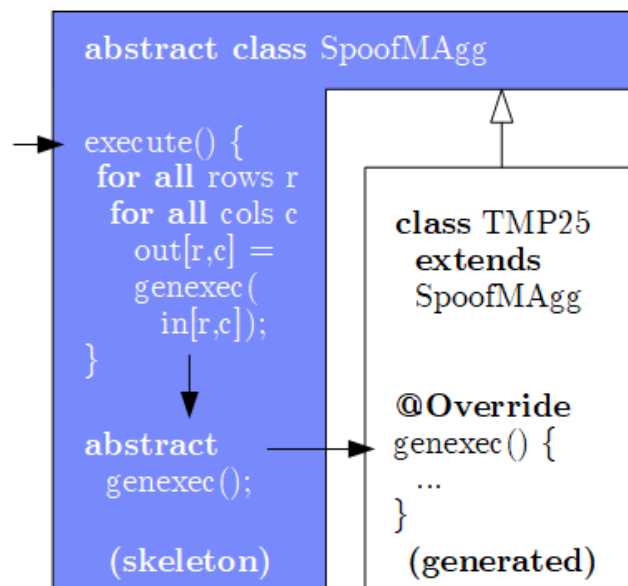
- Base (w/o fused ops): **10**
- Fused (w/ fused ops): **4**



# Codegen Example L2SVM, cont. (Cell/MAgg)

## Template Skeleton

- Data access, blocking
- Multi-threading
- Final aggregation



## # of Vector Intermediates

- Gen (codegen ops): 0

```

public final class TMP25 extends SpoofMAgg {
    public TMP25() {
        super(false, AggOp.SUM, AggOp.SUM);
    }
    protected void genexec(double a, SideInput[] b,
        double[] scalars, double[] c, ...) {
        double TMP11 = getValue(b[0], rowIndex);
        double TMP12 = getValue(b[1], rowIndex);
        double TMP13 = a * scalars[0];
        double TMP14 = TMP12 + TMP13;
        double TMP15 = TMP11 * TMP14;
        double TMP16 = 1 - TMP15;
        double TMP17 = (TMP16 > 0) ? 1 : 0;
        double TMP18 = a * TMP17;
        double TMP19 = TMP18 * a;
        double TMP20 = TMP16 * TMP17;
        double TMP21 = TMP20 * TMP11;
        double TMP22 = TMP21 * a;
        c[0] += TMP19;
        c[1] += TMP22;
    }
}
  
```

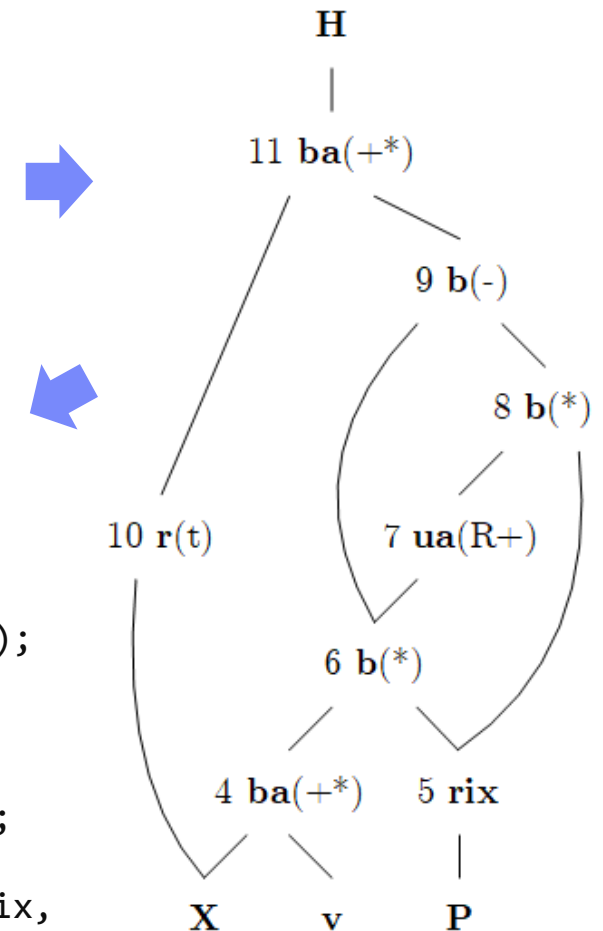
# Codegen Example MLogreg (Row)

## ■ MLogreg Inner Loop

(main expression on feature matrix X)

```
1: Q = P[, 1:k] * (X %%% v)
2: H = t(X) %%% (Q - P[, 1:k] * rowSums(Q))
```

```
public final class TMP25 extends SpoofRow {
  public TMP25() {
    super(RowType.COL_AGG_B1_T, true, 5);
  }
  protected void genexecDense(double[] a, int ai,
    SideInput[] b, double[] c, ..., int len) {
    double[] TMP11 = getVector(b[1].vals(rix),...);
    double[] TMP12 = vectMatMult(a, b[0].vals(rix),...);
    double[] TMP13 = vectMult(TMP11, TMP12, 0, 0,...);
    double TMP14 = vectSum(TMP13, 0, TMP13.length);
    double[] TMP15 = vectMult(TMP11, TMP14, 0,...);
    double[] TMP16 = vectMinus(TMP13, TMP15, 0, 0,...);
    vectOuterMultAdd(a, TMP16, c, ai, 0, 0,...);
  }
  protected void genexecSparse(double[] avals, int[] aix,
    int ai, SideInput[] b, ..., int len) {...}
}
```



# Candidate Exploration (by example MLogreg)

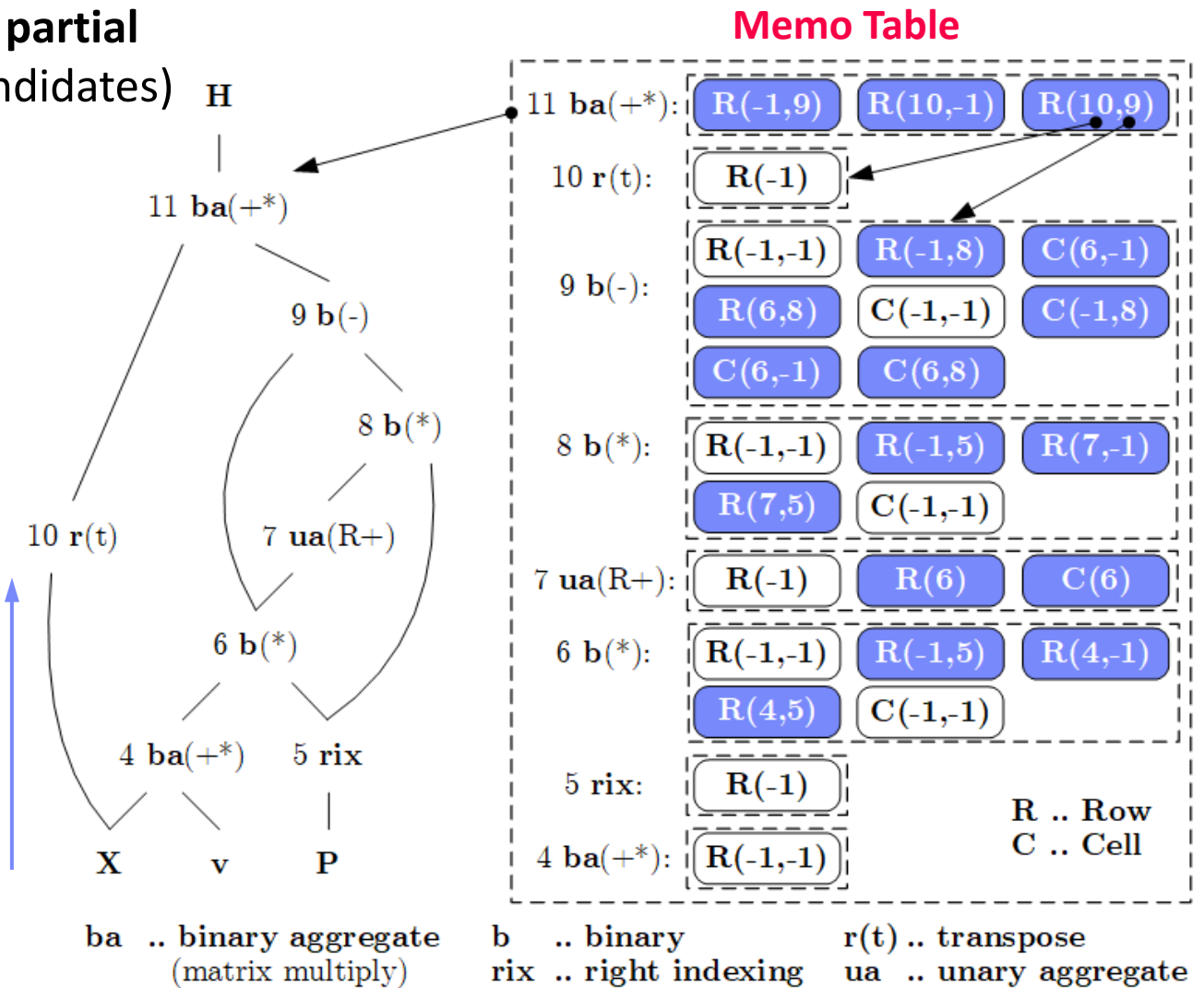
- Memo Table for partial fusion plans (candidates)

- OFMC Template Fusion API

- Open
- Fuse, Merge
- Close

- OFMC Algorithm

- **Bottom-up Exploration** (single-pass, template-agnostic)
- Linear space and time

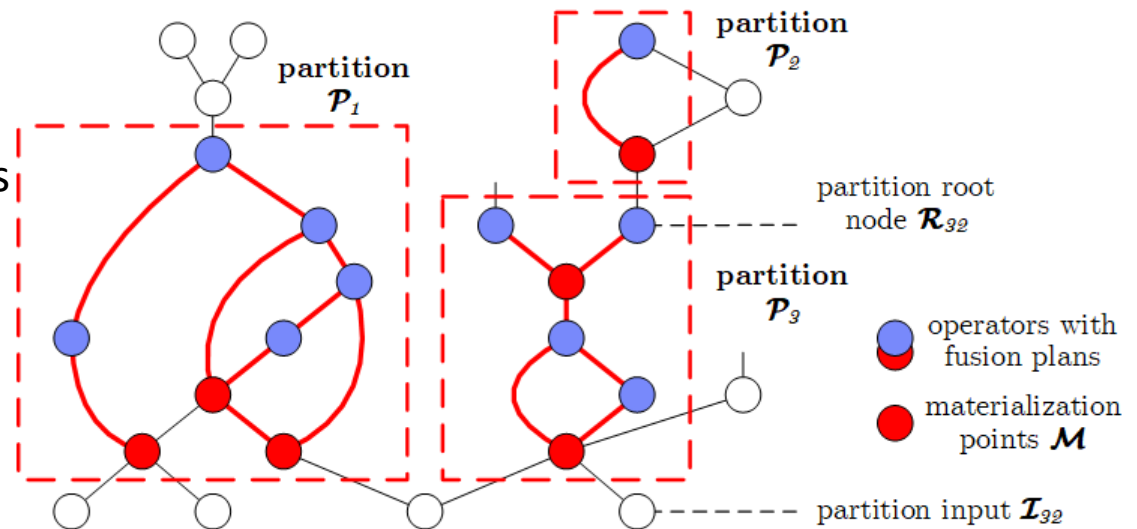


# Candidate Selection (Partitions and Interesting Points)

## #1 Determine Plan Partitions

- Materialization Points  $\mathcal{M}$
- Connected components of fusion references
- Root and input nodes

→ Optimize partitions independently



## #2 Determine Interesting Points

- Materialization Point Consumers:** Each data dependency on materialization points considered separately
- Template / Sparse Switches:** Data dependencies where producer has templates that are non-existing for consumers

→ Optimizer considers all  $2^{|\mathcal{M}'|}$  plans (with  $|\mathcal{M}'| \geq |\mathcal{M}_i|$ ) per partition

# Candidate Selection, cont. (Costs and Constraints)

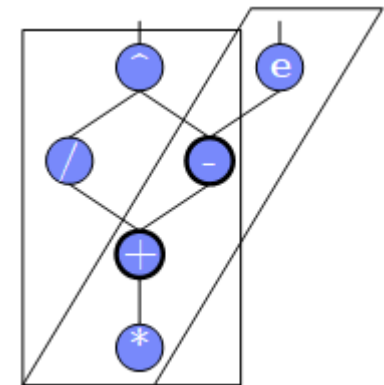
## Overview Cost Model

- Cost partition with **analytical cost model** based on peak memory and compute bandwidth
- Plan comparisons / fusion errors don't propagate / dynamic recompilation

$$C(\mathcal{P}_i|\mathbf{q}) = \sum_{p \in \mathcal{P}_i|\mathbf{q}} \left( \hat{T}_p^w + \max \left( \hat{T}_p^r, \hat{T}_p^c \right) \right)$$

## #3 Evaluate Costs

- #1: Memoization of already processed sub-DAGs
  - #2: Account for shared reads and CSEs within operators
  - #3: Account for redundant computation (overlap)
- **DAG traversal** and **cost vectors** per fused operator (with memoization of pairs of operators and cost vectors)



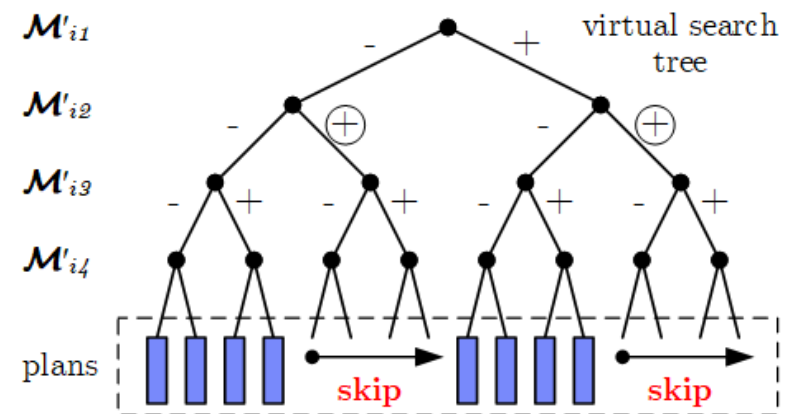
## #4 Handle Constraints

- Prefiltering** violated constraints (e.g., row template in distributed ops)
- Assign **infinite costs for violated constraints** during costing

# Candidate Selection, cont. (MPSkipEnum and Pruning)

## #5 Basic Enumeration

- Linearized search space: from - to \*
- ```
for( j in 1:pow(2, |M'_i|) )
  q = createAssignment(j)
  C = getPlanCost(P_i, q)
  maintainBest(q, C)
```

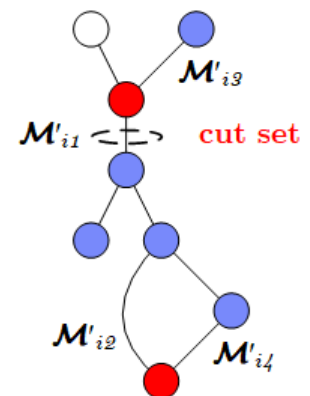


## #6 Cost-Based Pruning

- Upper bound:** cost  $C^U$  of best plan  $q^*$  (monotonically decreasing)
- Opening heuristic:** evaluate FA and FNR heuristics first
- Lower bound:**  $C^{LS}$  (read input, write output, min compute) + dynamic  $C^{LD}$  (materialize intermediates  $q$ )  $\rightarrow$  **skip subspace** if  $C^U \leq C^{LS} + C^{LD}$

## #7 Structural Pruning

- Observation:** Assignments can create independent sub problems
- Build **reachability graph** to determine **cut sets**
- During enum: probe cut sets, recursive enum, combine, and skip



# Ahead-of-Time Compilation

## TensorFlow `tf.compile`

- Compile entire TF graph into binary function w/ low footprint
- **Input:** Graph, config (feeds+fetches w/ fixed shape sizes)
- **Output:** x86 binary and C++ header (e.g., inference)
- **Specialization for frozen model and sizes**



[Chris Leary, Todd Wang:  
XLA – TensorFlow, Compiled!,  
TF Dev Summit 2017]

## PyTorch Compile



- Compile Python functions into ScriptModule/ScriptFunction
- Lazily collect operations, optimize, and JIT compile
- Explicit `jit.script` call or `@torch.jit.script`

```
a = torch.rand(5)
def func(x):
    for i in range(10):
        x = x * x # unrolled into graph
    return x
```

```
jitfunc = torch.jit.script(func) # JIT
jitfunc.save("func.pt")
```



[Vincent Quenneville-Bélair:  
How PyTorch Optimizes  
Deep Learning Computations,  
Guest Lecture Stanford 2020]



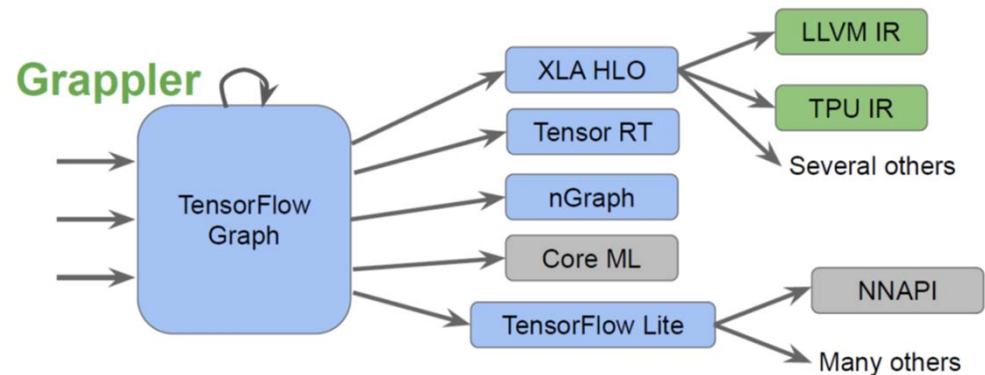
# Excursus: MLIR

[Rasmus Munk Larsen, Tatiana Shpeisman:  
TensorFlow Graph Optimizations,  
Guest Lecture Stanford 2019]



## ■ Motivation TF Compiler Ecosystem

- Different IRs and compilation chains for runtime backends
- **Duplication of infrastructure** and fragile error handling



## ■ MLIR (Multi-level, Machine Learning IR)

- SSA-based IR, similar to LLVM
- Hierarchy of modules, functions, blocks, and operations (text-based)
- **Dialects for different backends** (defined ops, customization)
- **Systematic lowering**

```
func @testFunction(%arg0: i32) {
    %x = call @thingToCall(%arg0)
        : (i32) -> i32
    br ^bb1
^bb1:
    %y = addi %x, %x : i32
    return %y : i32
}
```

# Discussion

## Programming Projects

# Example Projects

## APIS and Algorithms

- **#1 Extended Python and Java Language Bindings** (addition ops)
- **#2 Built-in Functions for Regression, Classification, Clustering**
- **#3 Built-in Functions for Time Series Missing Value Imputation**
- **#4 Neural Collaborative Filtering** (see MLPerf benchmark)
- **#5 DL-based Entity Resolution Primitives** (baseline implementation)
- **#6 Model Selection Primitives** (BO, multi-armed bandit, hyperband)

## Documentation, Tutorials, and Tests

- **#7 Documentation and Tutorials** (for different target users)
- **#8 Extended Test Framework** (comparisons, caching, remove redundancy)
- **#9 New Optimizer Test Framework** (rewrites, optimization passes)
- **#10 SLAB Benchmark** (benchmark driver, summary)
- **#11 Performance Testsuite** (extend algorithm-level suite)

# Example Project, cont.

## Tools and Experimental

- **#12 ONNX Graph Importer/Exporter** (DML script / HOP DAG generation)
- **#13 Auto Differentiation** (builtin function and compiler)
- **#14 SLIDE Operators and Runtime Integration** (Sub-Linear DL Engine)
- **#15 Quantum Neural Networks** (Grover's Quantum Search, Qiskit/TFQ)

## Compiler Features

- **#16 Loop Vectorization Rewrites** (more general framework)
- **#17 Canonicalization Rewrite Framework** (refactoring, new rewrites)
- **#18 Extended CSE & Constant Folding** (commutativity, one-shot)
- **#19 Extended Matrix Multiplication Chain Opt** (sparsity, rewrites)
- **#20 Extended Update In-Place Framework** (reference counting)
- **#21 LLVM Code Generator Framework** (extension CPU native)
- **#22 Operator Scheduling Algorithms** (baselines)

# Example Projects, cont.

## Runtime Features

- **#23 Feature Transform: Equi-Height/Custom Binning** (local, distributed)
- **#24 Federated Feature Transformations** (recoding, one-hot encoding)
- **#25 Selected N-Dimensional Tensor Operations**
- **#26 Compression Planning Extensions** (co-coding search algorithm)
- **#27 Extended Intel MKL-DNN Runtime Operations** (beyond conv2d)
- **#28 Selected Dense and Sparse GPU Operations** (libs, custom)

## I/O Subsystem

- **#29 Lineage-Exploitation in Buffer Pool** (for recomputation)
- **#30 Multi-threaded Buffer Pool Eviction** (multi-part/multi-disk)
- **#31 Extended I/O Framework for Other Formats** (e.g., NetCDF, HDF5, Arrow)
- **#32 Protobuf reader/writer into Data Tensor** (local, distributed)

# Conclusions

## Summary

- Motivation and Terminology
- Runtime Adaptation
- Operator Fusion & JIT
- Discussion Programming Projects

## Recommended Reading

[Chris Leary, Todd Wang: XLA – TensorFlow, Compiled!, TF Dev Summit 2017, [https://www.youtube.com/watch?time\\_continue=1541&v=kAOanJczHA0&feature=emb\\_logo](https://www.youtube.com/watch?time_continue=1541&v=kAOanJczHA0&feature=emb_logo)]



➔ Impact of Size Inference and Costs (lecture 03)

➔ Ubiquitous Rewrite, Fusion, and Codegen/JIT Opportunities

## Next Lectures (Runtime Aspects)

- 05 Data- and Task-Parallel Execution (batch/prog) [Apr 03]
- 06 Parameter Servers (mini-batch) [Apr 24]
- 07 Hybrid Execution and HW Accelerators [May 08]
- 08 Caching, Partitioning, Indexing and Compression [May 15]