

Data Management

05 Query Languages (SQL)

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Announcements/Org

■ #1 Video Recording

- Link in **TeachCenter** & **TUbe** (lectures will be public)
- <https://tugraz.webex.com/meet/m.boehm>
- Corona traffic light **RED** until end of April



■ #2 Reminder Communication

- **Newsgroup:** <news://news.tugraz.at/tu-graz.lv.dbase>
- **Office hours:** Mo 12.30-1.30pm (<https://tugraz.webex.com/meet/m.boehm>)

■ #3 Exercise Submissions

- **Exercise 1:** Mar 30 + 7 late days, grading in progress
- **Exercise 2:** Apr 27, published Apr 07
- https://mboehm7.github.io/teaching/ss21_dbs/02_ExerciseQueriesAPIs.pdf

Agenda

- **Structured Query Language (SQL)**
- **Other Query Languages (XML, JSON)**
- **Exercise 2: Query Languages and APIs**

Structured Query Language (SQL)

What is a(n) SQL Query?

```

SELECT Firstname, Lastname, Affiliation, Location
FROM Participant AS R, Locale AS S
WHERE R.LID = S.LID
AND Location LIKE '%, GER'
    
```

#1 **Declarative:**
what not how



Firstname	Lastname	Affiliation	Location
Volker	Markl	TU Berlin	Berlin, GER
Thomas	Neumann	TU Munich	Munich, GER

#2 **Flexibility:**
closed → composability

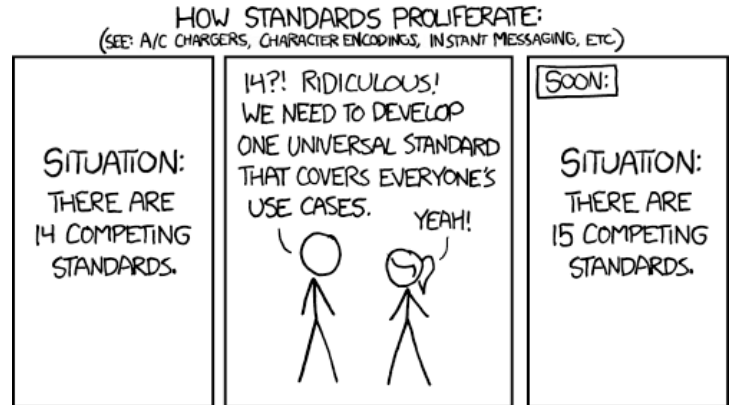
#3 **Automatic Optimization**

#4 **Physical Data Independence**

Why should I care?

SQL as a Standard

- Standards ensure **interoperability**, avoid **vendor lock-in**, and protect **application investments**
- Mature standard** with heavy industry support for decades
- Rich eco system** (existing apps, BI tools, services, frameworks, drivers, design tools, systems)



<https://xkcd.com/927/>

SQL is here to stay

- Foundation of mobile/server **application data management**
- Adoption of existing standard** by new systems (e.g., SQL on Hadoop, cloud DBaaS)
- Complemented by NoSQL abstractions, see lecture **10 NoSQL (key-value, document, graph)**



Overview SQL

- **Structured Query Language (SQL)**
 - Current Standard: ISO/IEC 9075:2016 (SQL:2016)
 - **Data Definition Language (DDL)** → Manipulate the database schema
 - **Data Manipulation Language (DML)** → Update and query database
 - **Data Control Language (DCL)** → Modify permissions

- **Dialects**

- Spectrum of system-specific dialects for **non-core features**
- Data types and size constraints
- Catalog, builtin functions, and tools
- Support for new/optional features
- Case-sensitive identifiers

Name	Examples
T-SQL	Microsoft, Sybase
PL/SQL	Oracle, (IBM)
PL/pgSQL	PostgreSQL, derived
Unnamed	Most systems

The History of the SQL Standard

[C. J. Date: A Critique of the
SQL Database Language.
SIGMOD Record 1984]

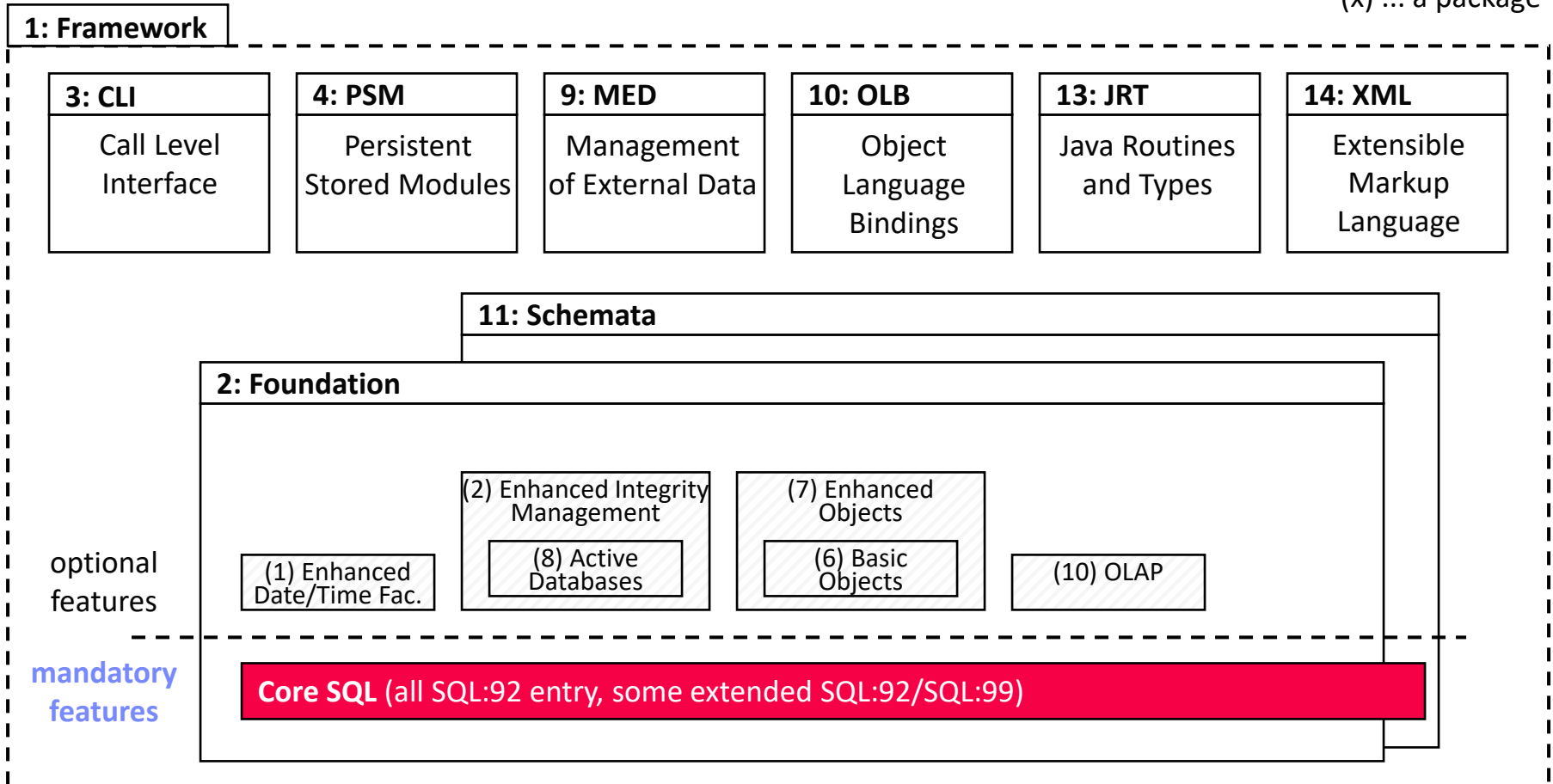


- **SQL:1986**
 - **Database Language SQL**, ANSI X3.135-1986, ISO-9075-1987(E)
 - '87 international edition
- **SQL:1989 (120 pages)**
 - **Database Language SQL with Integrity Enhancements**, ANSI X3.135-1989, ISO-9075-1989(E)
- **SQL:1992 (580 pages)**
 - **Database Language SQL**, ANSI X3-1992, ISO/IEC-9075 1992, DIN 66315
 - '95 SQL/CLI (part 3), '96 SQL/PSM (part 4)
- **SQL:1999 (2000 pages)**
 - **Information Technology – Database Language – SQL**, ANSI/ISO/IEC-9075 1999
 - Complete reorg, '00 OLAP, '01 SQL/MED, '01 SQL/OLB, '02 SQL/JRT
- **SQL:2003 (3764 pages)**
 - **Information Technology – Database Language – SQL**, ANSI/ISO/IEC-9075 2003

The History of the SQL Standard, cont.

Overview SQL:2003

x: ... a part
(x) ... a package



The History of the SQL Standard, cont.

Since SQL:2003 overall structure remained unchanged ...

- **SQL:2008 (???? pages)**
 - **Information Technology – Database Language – SQL**, ANSI/ISO/IEC-9075 2003
 - E.g., **XML** XQuery extensions, case/trigger extension
- **SQL:2011 (4079 pages)**
 - **Information Technology – Database Language – SQL**, ANSI/ISO/IEC-9075 2011
 - E.g., time periods, temporal constraints, time travel queries
- **SQL:2016 (???? pages)**
 - **Information Technology – Database Language – SQL**, ANSI/ISO/IEC-9075 2016
 - E.g., **JSON** documents and functions (optional)

➔ **Note:** We can only discuss common primitives

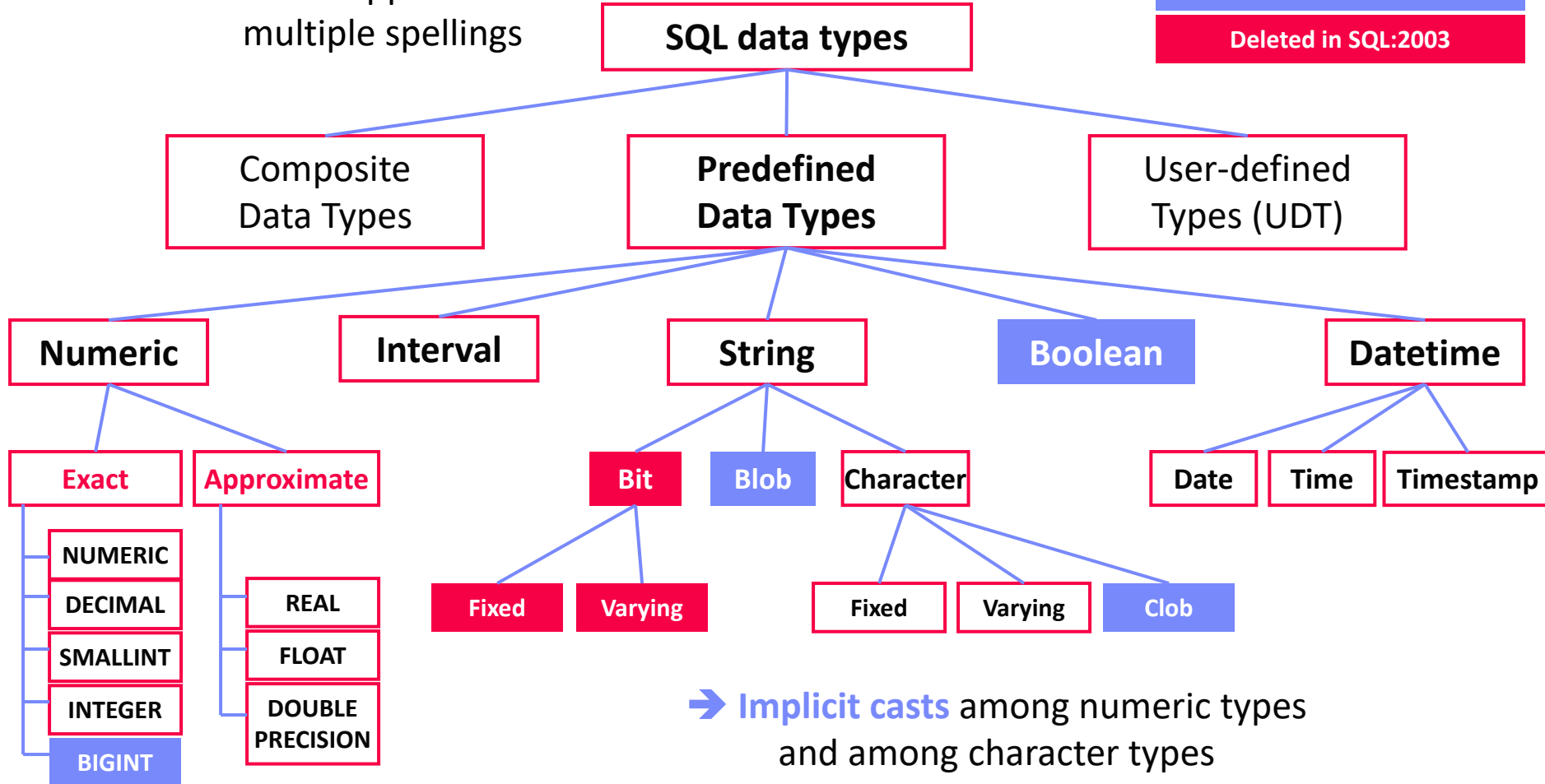
[Working Draft SQL:2011:
[https://www.wiscorp.com/
SQLStandards.html](https://www.wiscorp.com/SQLStandards.html)]

Data Types in SQL:2003

Large Variety of Types

- With support for multiple spellings

Added in SQL:1999 / SQL:2003
 Deleted in SQL:2003



Data Types in PostgreSQL

Appropriate, Brief, Complete

■ Strings

- **CHAR(n)** → fixed-length character sequence (padded to n)
- **VARCHAR(n)** → variable-length character sequence (n max)
- TEXT → variable-length character sequence

■ Numeric

- SMALLINT → 2 byte integer (signed short)
- **INT/INTEGER** → 4 byte integer (signed int)
- SERIAL → INTEGER w/ auto increment
- **NUMERIC(p, s)** → exact real with p digits and s after decimal point

■ Time

- **DATE** → date
- **TIMESTAMP/TIMESTAMPTZ** → date and time, timezone-aware if needed

■ JSON

- JSON → text JSON representation (requires reparsing)
- JSONB → binary JSON representation

Create, Alter, and Delete Tables

Templates in SQL
Examples in PostgreSQL

■ Create Table

- Typed attributes
- Primary and foreign keys
- **NOT NULL**, **UNIQUE** constraints
- **DEFAULT** values
- **CHECK** constraints

```
CREATE TABLE Students (
  SID INTEGER PRIMARY KEY,
  Fname VARCHAR(128) NOT NULL,
  Lname VARCHAR(128) NOT NULL,
  Mtime DATE DEFAULT CURRENT_DATE
);
```

```
CREATE TABLE Students AS SELECT ...;
```

■ Alter Table

- **ADD/DROP** columns
- **ALTER** data type, defaults, constraints, etc

```
ALTER TABLE Students ADD DoB DATE;
```

```
ALTER TABLE Students ADD CONSTRAINT
  PKStudent PRIMARY KEY(SID);
```

■ Delete Table

- Delete table
- **Note:** order of tables matters due to referential integrity

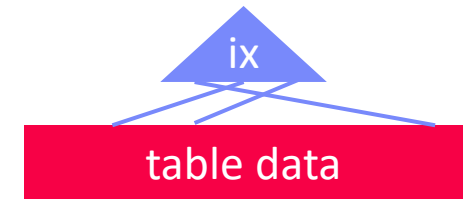
```
DROP TABLE Students; -- sorry
DROP TABLE Students CASCADE;
DROP TABLE IF EXISTS Countries,
  Cities, Airports, Airlines,
  Routes, Planes, Routes_Planes;
```

Create and Delete Indexes

■ Create Index

- Create a secondary (nonclustered) index on a set of attributes
- **Clustered:** tuples sorted by index
- **Non-clustered:** sorted attribute with tuple references
- Can specify uniqueness, order, and indexing method
- **PostgreSQL methods:** btree, hash, gist, and gin
 - ➔ see lecture **07 Physical Design and Tuning**

```
CREATE INDEX ixStudLname
ON Students USING btree
(Lname ASC NULLS FIRST);
```



■ Delete Index

- Drop indexes by name

```
DROP INDEX ixStudLname;
```

■ Tradeoffs

- Indexes often automatically created for **primary keys** / **unique** attributes
- **Lookup/scan performance** vs **insert performance**

Database Catalog

[Meikel Poes: **TPC-H**. Encyclopedia of Big Data Technologies 2019]

■ Catalog Overview

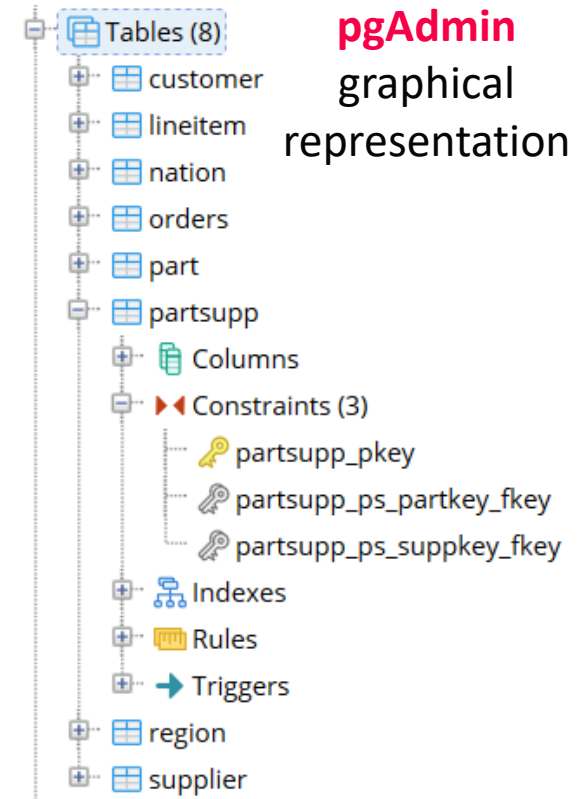
- **Meta data** of all database objects (tables, constraints, indexes) → **mostly read-only**
- **Accessible through SQL**
- Organized by schemas (**CREATE SCHEMA tpch;**)

■ SQL Information_Schema

- Schema with tables for all tables, views, constraints, etc
- **Example:** check for existence of accessible table

```
SELECT 1 FROM information_schema.tables
WHERE table_schema = 'tpch'
      AND table_name = 'customer'
```

(defined as views over PostgreSQL catalog tables)



Insert

■ Insert Tuple

- **Insert a single tuple** with implicit or explicit attribute assignment

```
INSERT INTO Students (SID, Lname, Fname, MTime, DoB)
VALUES (7, 'Boehm', 'Matthias', '2002-10-01', '1982-06-25');
```

- Insert attribute key-value pairs to use auto increment, defaults, NULLs, etc

```
INSERT INTO Students (Lname, Fname, DoB) SERIAL SID,
VALUES ('Boehm', 'Matthias', '1982-06-25'), DEFAULT MTime
(...), (...);
```

■ Insert Table

- **Redirect query result into INSERT** (append semantics)

```
INSERT INTO Students
SELECT * FROM NewStudents;
```

Analogy Linux redirect (append):
 cat NewStudents.txt >> Students.txt

Update and Delete

Update Tuple/Table

- **Set-oriented update** of attributes
- Update single tuple via predicate on **primary key**

```
UPDATE Students
  SET MTime = '2002-10-02'
  WHERE LName = 'Boehm';
```

Delete Tuple/Table

- **Set-oriented delete** of tuples
- Delete single tuple via predicate on **primary key**

```
DELETE FROM Students
  WHERE extract(year
    FROM mtime) < 2010;
```

Note: Time travel and multi-version concurrency control

- Deleted tuples might be just **marked as inactive**
- See lecture **09 Transaction Processing and Concurrency**

Basic Queries

Basic Query Template

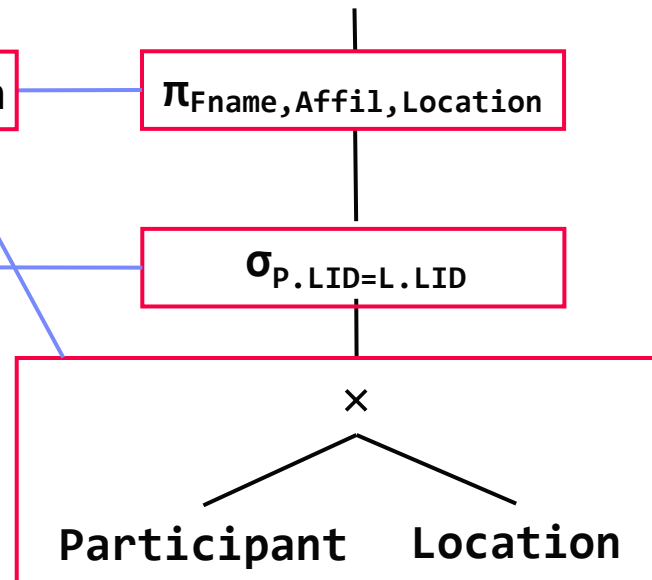
- **Select-From-Where**
- Grouping and Aggregation
- Having and ordering
- Duplicate elimination

```

SELECT [DISTINCT] <column_list>
FROM [<table_list> |
    <table1> [RIGHT | LEFT | FULL] JOIN
    <table2> ON <condition>]
[WHERE <predicate>]
[GROUP BY <column_list>]
[HAVING <grouping predicate>]
[ORDER BY <column_list> [ASC | DESC]]
    
```

Example

- **SELECT** Fname, Affil, Location
FROM Participant AS P,
 Locale AS L
WHERE P.LID = L.LID;



Basic Queries, cont.

▪ Distinct and All

- **Distinct and all** alternatives
- Projection w/ **bag semantics** by default

```
SELECT DISTINCT Lname, Fname
FROM Students;
```

▪ Sorting

- Convert a **bag** into a **sorted list** of tuples; order lost if used in other ops
- Single order: (Lname, Fname) **DESC**
- Evaluated last in a query tree

```
SELECT * FROM Students
ORDER BY Lname DESC,
Fname DESC;
```

▪ Set Operations

- See **04 Relational Algebra and Calculus**
→ **UNION, INTERSECT, EXCEPT**
- Set operations **set semantics** by default
→ **DISTINCT** (set) vs **ALL** (bag)

```
(SELECT Firstname, Lastname
FROM Participant2018)
UNION DISTINCT
(SELECT Firstname, Lastname
FROM Participant2013)
```

Grouping and Aggregation

Grouping and Aggregation

- **Grouping:** determines the distinct groups
- **Aggregation:** compute aggregate $f(B)$ per group
- Column list can only contain **grouping columns**, **aggregates**, or **literals**
- **Having:** selection predicate on groups and aggregates

Example

- Sales (Customer, Location, Product, Quantity, Price)
- **Q: Compute number of sales sumQ and revenue per product sumQP**

```

SELECT Product,
       sum(Quantity) AS SumQ,
       sum(Quantity*Price) AS SumQP
FROM Sales
GROUP BY Product
    
```



Product	Quantity	Price
A	1	10
B	3	20
A	2	10
B	1	20



Product	SumQ	SumQP
A	3	30
B	4	80

BREAK (and Test Yourself)

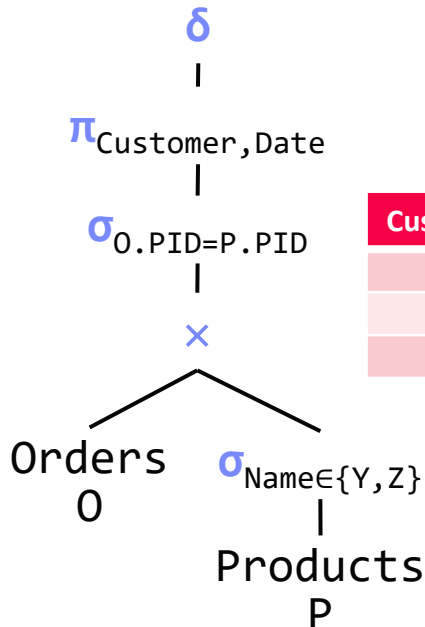
- Task: SQL queries for the following query trees.

Orders

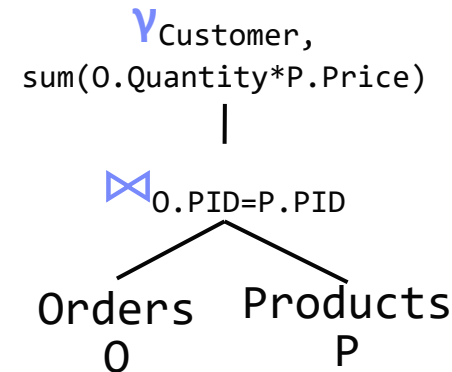
OID	Customer	Date	Quantity	PID
1	A	'2019-06-22'	3	2
2	B	'2019-06-22'	1	3
3	A	'2019-06-22'	1	4
4	C	'2019-06-23'	2	2
5	D	'2019-06-23'	1	4
6	C	'2019-06-23'	1	1

Products

PID	Name	Price
1	X	100
2	Y	15
4	Z	75
3	W	120



Customer	Date
A	'2019-06-22'
C	'2019-06-23'
D	'2019-06-23'



Customer	Sum
A	120
B	120
C	130
D	75

```

SELECT DISTINCT Customer, Date
FROM Orders O, Products P
WHERE O.PID = P.PID
AND Name IN('Y', 'Z')
    
```

```

SELECT Customer,
       sum(O.Quantity * P.Price)
FROM Orders O, Products P
WHERE O.PID = P.PID
GROUP BY Customer
    
```

Subqueries

■ Subqueries in Table List

- Use a subquery result like a base table
- Modularization with **WITH C AS (SELECT ...)**

```
SELECT S.Fname, S.Lname, C.Name
FROM Students AS S,
     (SELECT CID, Name FROM Country
      WHERE ...) AS C
WHERE S.CID=C.CID;
```

■ Subqueries w/ IN

- Check containment of values in result set of sub query

```
SELECT Product, Quantity, Price
FROM Sales
WHERE Product NOT IN(
     SELECT Product FROM Sales
     GROUP BY Product
     HAVING sum(Quantity*Price)>1e6)
```

■ Other subqueries

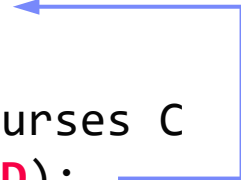
- **EXISTS**: existential quantifier $\exists x$ for correlated subqueries
- **ALL**: comparison (w/ universal quantifier $\forall x$)
- **SOME/ANY**: comparison (w/ existential quantifier $\exists x$)

Correlated and Uncorrelated Subqueries

■ Correlated Subquery

- **Evaluated subquery for every tuple** of outer query
- Use of attribute from table bound in outer query inside subquery

```
SELECT P.Fname, P.Lname
FROM Professors P,
WHERE NOT EXISTS(
    SELECT * FROM Courses C
    WHERE C.PID=P.PID);
```



■ Uncorrelated Subquery

- Evaluate subquery just once
- No attribute correlations between subquery and outer query

```
SELECT P.Fname, P.Lname
FROM Professors P,
WHERE P.PID NOT IN(
    SELECT PID FROM Courses);
```

■ Query Unnesting (de-correlation)

- Rewrite during query compilation
- See lecture [08 Query Processing](#)

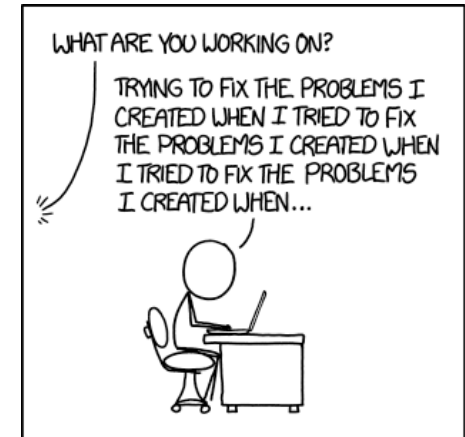
[Thomas Neumann, Alfons Kemper: Unnesting Arbitrary Queries. **BTW 2015**]



Recursive Queries

Approach

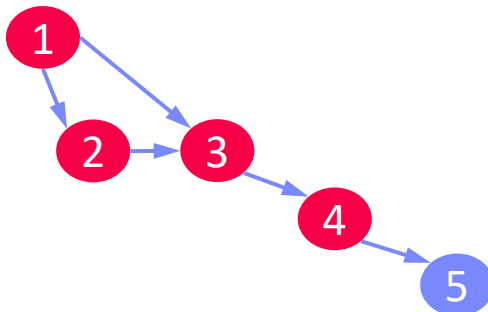
- **WITH RECURSIVE** <name> (<arguments>)
- Compose recursive table from **non-recursive term**, **union all/distinct**, and **recursive term**
- Terminates when recursive term yields empty result



<https://xkcd.com/1739/>

Example

- Courses(CID, Name),
Precond(pre REF CID, suc REF CID)
- Dependency graph (pre → suc)



```

WITH RECURSIVE rPrereq(p,s) AS(
    (SELECT pre, suc
     FROM Precond WHERE suc=5)
    UNION DISTINCT
    (SELECT B.pre, B.suc
     FROM Precond B, rPrereq R
     WHERE B.suc = R.p)
)
SELECT DISTINCT p FROM rPrereq
    
```

4
3
1
2

Procedures and Functions

Overview Procedures and Functions

- Stored programs, written in **PL/pgSQL** or other languages
- Control flow (loops, branches) and SQL queries**

(Stored) Procedures

- Can be called standalone via **CALL** <proc_name>(<args>);
- Procedures return no outputs

```

CREATE PROCEDURE prepStud(a INT)
LANGUAGE PLPGSQL AS $$
BEGIN
    DELETE FROM Students;
    INSERT INTO Students
        SELECT * FROM NewStudents;
END; $$;
    
```

Functions

- Can be called standalone or inside queries
- Functions are value mappings
- Table functions can return sets of records with multiple attributes

```

CREATE FUNCTION sampleProp(FLOAT)
RETURNS FLOAT
AS 'SELECT $1 * (1 - $1);'
LANGUAGE SQL;
    
```

Triggers

Overview Trigger

- Similar to stored procedure but register ON INSERT, DELETE, or UPDATE
- Allows complex check constraints and active behavior such as replication, auditing, etc (good and bad)

Trigger **CREATE TRIGGER** <triggername>

Template

```

BEFORE | AFTER | INSTEAD OF
INSERT | DELETE | (UPDATE OF <column_list>)
ON <tablename>
    
```

```

[REFERENCING <old_new_alias_list>]
    
```

```

[FOR EACH {ROW | STATEMENT}]
    
```

```

[WHEN (<search condition>)]
    
```

```

<SQL procedure statement> |
    
```

```

BEGIN ATOMIC
    
```

```

    {<SQL Procedure statement>;}...
    
```

```

END
    
```

Event

Condition

Action

Not supported in
PostgreSQL
(need single UDF)

Views and Authorization

■ Creating Views

- **Create a logical table from a query**
- Inserts can be propagated back to base relations only in special cases
- **Allows authorization** for subset of tuples

```

CREATE VIEW TeamDM AS
SELECT * FROM
    Employee E, Employee M
WHERE E.MgrID = M.EID
AND M.login = 'mboehm';
  
```

■ Access Permissions Tables/Views

- **Grant** query/modification rights on database objects for specific users, roles
- **Revoke** access rights from users, roles (recursively revoke permissions of dependent views via **CASCADE**)

```

GRANT SELECT
ON TABLE TeamDM
TO mboehm;
  
```

```

REVOKE SELECT
ON TABLE TeamDM
FROM mboehm;
  
```

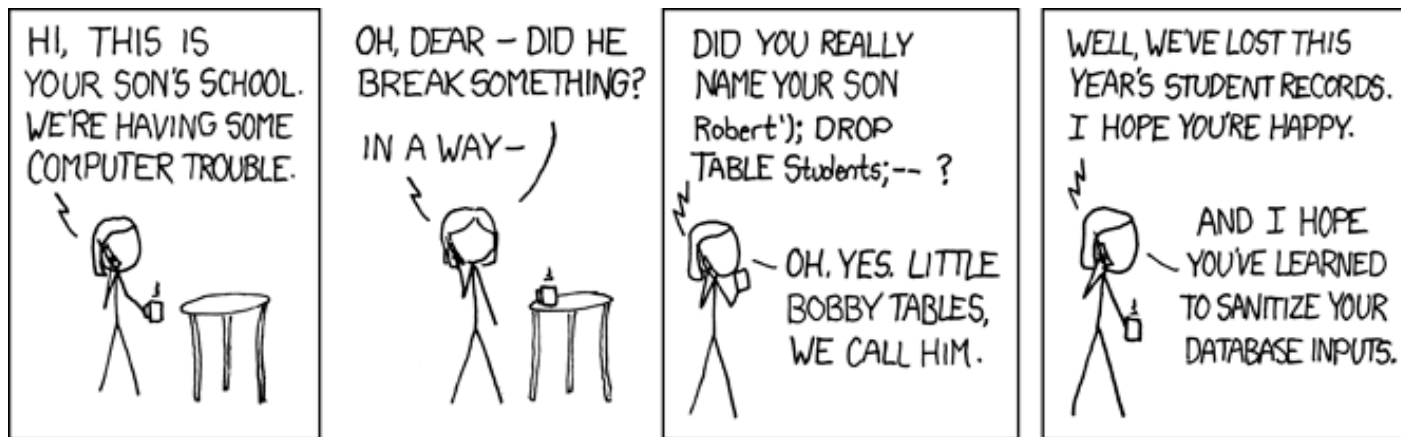
Beware of SQL Injection



■ Problematic SQL String Concatenation

```
INSERT INTO Students (Lname, Fname)
VALUES ('"+ @lname +"', '"+ @fname +"' );";
```

■ Possible SQL-Injection Attack



<https://xkcd.com/327/>

```
INSERT INTO Students (Lname, Fname) VALUES ('Smith', 'Robert');
DROP TABLE Students; --');
```

Other Query Languages (XML, JSON)

No really, why should I care?

- **Semi-structured XML and JSON**
 - **Self-contained documents** for representing nested data
 - **Common data exchange formats** without redundancy of flat files
 - Human-readable formats → often used for SW configuration
- **Goals**
 - **Awareness of XML and JSON** as data models
 - Query languages and embedded querying in SQL

XML (Extensible Markup Language)

■ XML Data Model

- Meta language to define specific **exchange formats**
- Document format for **semi-structured data**
- Well formedness
- XML schema / DTD

```
<?xml version="1.0" encoding="UTF-8"?>
<data>
  <student id="1">
    <course id="INF.01017UF" name="DM"/>
    <course id="706.550" name="AMLS"/>
  </student>
  <student id="5">
    <course id="706.520" name="DIA"/>
  </student>
</data>
```

■ XPath (XML Path Language)

- Query language for **accessing collections of nodes** of an XML document
- Axis specifies for ancestors, descendants, siblings, etc

```
/data/student[@id='1']/course/@name
```

↓
"DM"
"AMLS"

■ XSLT (XML Stylesheet Language Transformations)

- Schema mapping (transformation) language for XML documents

■ XQuery

- Query language to extract, transform, and analyze XML documents

XML in PostgreSQL, cont.

■ Overview XML in PostgreSQL

- Data types **TEXT** or **XML** (well-formed, type-safe operations)
- ISO/IEC 9075-14 XML-related specifications (SQL/XML)

■ Creating XML

- Various **built-in functions** to parse documents, and create elements/attributes
- XMLPARSE(<xml_document>) → **XML type**
- XMLELEMENT / XMLATTRIBUTES

```
INSERT INTO Students
(Fname, Lname, Doc)
VALUES ('John', 'Smith',
xmlparse(<source_doc>));
```

■ Processing XML

- Execute **XPath** expressions on XML types
- XMLEXIST with **XPath instead of XQuery**
- XPATH with optional namespace handling

```
SELECT Fname, Lname,
xpath('/student/@id', Doc)
FROM Students
```


JSON (JavaScript Object Notation)

■ JSON Data Model

- Data exchange format for **semi-structured data**
- **Not as verbose as XML** (especially for arrays)
- Popular format (e.g., Twitter)



```
{“students:”[
  {“id”: 1, “courses”:[
    {“id”:“INF.01017UF”, “name”:“DM”},
    {“id”:“706.550”, “name”:“AMLS”}]}],
  {“id”: 5, “courses”:[
    {“id”:“706.520”, “name”:“DIA”}]}],
]}
```

■ Query Languages

- **Most common: libraries** for tree traversal and data extraction
- **JSONiq**: XQuery-like query language
- **JSONPath**: XPath-like query language

JSONiq Example:

```
declare option jsoniq-version “...”;
for $x in collection(“students”)
  where $x.id lt 10
  let $c := count($x.courses)
  return {“sid”:$x.id, “count”:$c}
```

[<http://www.jsoniq.org/docs/JSONiq/html-single/index.html>]

JSON in PostgreSQL, cont.

Overview JSON in PostgreSQL

- Alternative data types: **JSON** (text), **JSONB** (binary, with restrictions)
- Implements RFC 7159, built-ins for conversion and access

Creating JSON

- Built-in functions for creating JSON from tables and tables from JSON input

```
SELECT row_to_json(t) FROM
      (SELECT Fname, Lname
       FROM Students) t
```

Processing JSON

- Specialized operators for **tree traversal and data extraction**
- > operator**: get JSON array element/object
- >> operator**: get JSON array element/object as text
- Built-in functions for extracting json (e.g., json_each)

```
SELECT Fname, Lname,
      Doc->students->>id
FROM Students
```

Exercise 2:

Query Languages and APIs

Published: **Apr 07, 2021**

(updated: data Apr 07, task description Apr 08)

Deadline: **Apr 27, 2021**

Exercises: Summer Olympics

Dataset

- Past Summer Olympics (games, teams, athletes, events, medals)
- Clone or download your copy from <https://github.com/tugraz-isds/datasets.git> (summer_olympics)
- **Finished data cleaning by Apr 07**



AthletesEvents.csv
HostCities.csv
NOCRegions.csv

Exercises

- **01** Data modeling (relational schema)
- **02** Data ingestion and SQL query processing
- **03** Physical design tuning, query processing, and transaction processing
- **04** Large-scale data analysis (distributed query processing and ML model training)



Task 2.1: Schema Creation via SQL (3/25 points)

■ Schema creation via SQL

- Relies on lectures [04 Relational Algebra](#) and [05 Query Languages \(SQL\)](#)
- Setup DBMS PostgreSQL, and start pgAdmin (UI), or psql (terminal)
- Create database db<studentID> and setup relational schema, including primary keys, foreign keys, NOT NULL, UNIQUE

■ Recommended Schema

- Feel free to use and submit the provided schema
- https://mboehm7.github.io/teaching/ss21_dbs/CreateSchema.sql

■ Partial Results

- CreateSchema.sql

Task 2.2 Data Ingestion via CLI (10/25 points)

■ Data Ingestion Program via ODBC/JDBC

- Relies on lectures [05 Query Languages \(SQL\)](#) and [06 APIs \(ODBC, JDBC\)](#)
- Write a program that performs **deduplication and data ingestion**
- Programming language of your choosing (Python, Java, C#, C++ recommended)

■ Data Ingestion Process

- Data: <https://github.com/tugraz-isds/datasets/tree/master/movies>
- Invoke your ingestion program as follows → script to compile and run

```
./runIngestData.sh ./AthleteEvents.csv ./HostCities.csv \  
./NOCRRegions.csv <host> <port> <database> <user> <password>  
(e.g., localhost 5432 db1234567 postgres postgres)
```

■ Partial Results

- Source code `IngestData.*`, and
- Script `runIngestData.sh`

Task 2.3: SQL Query Processing (10/25 points)

■ SQL Query Processing

- Relies on lecture [05 Query Languages \(SQL\)](#)
- Expected results: https://mboehm7.github.io/teaching/ss21_dbs/Results.zip

■ List of Queries

- **Q01:** Obtain the detailed information of the athlete Usain St. Leo Bolt. (return name, gender, day of birth, height)
- **Q02:** Compute the distinct cities that hosted the Olympic games. (return distinct {city name, country name} pairs, **sorted** ascending by country name)
- **Q03:** List the athletes of team Austria 2012. (return name, gender, day of birth, **sorted** ascending by gender, day of birth)
- **Q04:** Compute for each sport type, the number of distinct event types as well as its last (i.e., maximum) year of occurrence. (return sport name, count, last occurrence, **sorted** descending by count)
- **Q05:** Determine the years in which female athletes whose names contain Ledecky won (in total) greater or equal than five gold or silver medals. (return year)

Task 2.3: SQL Query Processing (10/25 points)

■ List of Queries, cont.

- **Q06:** How many medals of each type did Michael Fred Phelps; II win? (return medal, count, **sorted** descending by count)
- **Q07:** Determine the top-10 athletes that participated between 1948 and 2016 in the most event occurrences but never won a single medal. (return name, day of birth, participation count, **sorted** descending by count, name).
- **Q08:** Compute the athlete-centric Olympic medal table of 2016, which counts every medal awarded (i.e., multiple medals in team events). (return country name, NOC, #gold, #silver, #bronze, #total, **sorted** descending by {#gold, #silver, #bronze, #total}) **// be careful about years**

■ Partial Results

- SQL Script for each query: Q01.sql, Q02.sql, ..., Q08.sql

Task 2.4: Query Plans (2/25 points)

■ Explain Query Plans

- Relies on lecture [04 Relational Algebra](#) and [05 Query Languages \(SQL\)](#)
- Obtain and **analyze execution plans** of Q05

■ Example

EXPLAIN VERBOSE

```
SELECT L.location, count(*)
  FROM Participant P,
       Locale L
 WHERE P.lid = L.lid
 GROUP BY L.location
 HAVING count(*)>1
```

```
"HashAggregate (...)" // grouping
" Output: l.location, count(*)"
" Group Key: l.location"
" Filter: (count(*) > 1)" // selection
" -> Hash Join (...)" // join
"   Output: l.location" // projection
"   Hash Cond: (l.lid = p.lid)"
"   -> Seq Scan on Locale l (...)"
"     Output: l.lid, l.location"
"   -> Hash (...)"
"     Output: p.lid" // projection
"     -> Seq Scan on Participant p (...)"
"       Output: p.lid"
```

■ Partial Results

- ExplainQ05.sql

Conclusions and Q&A

■ Summary

- History and fundamentals of the **Structured Query Language (SQL)**
- Awareness of **XML and JSON** (data model and querying)

■ Exercise Submissions

- **Exercise 1:** Mar 30 + 7 late days, grading in progress
- **Exercise 2:** Apr 27, published Apr 07

■ Next Lectures

- **06 APIs (ODBC, JDBC, OR frameworks)** [Apr 19] → videos available
- **07 Physical Design and Tuning** [Apr 26]
- **08 Query Processing** [May 03]
- **09 Transaction Processing and Concurrency** [May 10]