



# Data Management 14 Q&A and Exam Preparation

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# **Exam Preparation**

Basic focus: fundamental concepts and ability to apply learned techniques to given problems





## **Exam Logistics**

### Timing

- Exam starts 10min after official start
- 90min working time (plenty of time to think about answers)
- Write into the worksheet if possible, additional paper allowed
- Grading will happen Feb 1 / Feb 2 → use exam Feb 6 as replacement

### Covered Content

- Must-have: Data modeling/normalization, SQL query processing
- Relational algebra, physical design, query and transaction processing
- NoSQL, distributed storage and computation, streaming

### Past Exams

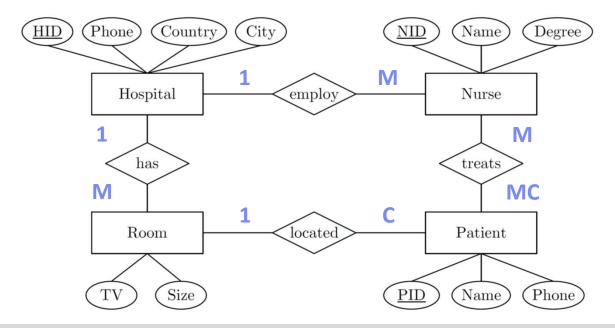
- 3x Data Management (previously known as Databases)
- 3x Databases (previously known as Databases 1)
- https://mboehm7.github.io/teaching/ss19 dbs/index.htm





# #1 Data Modeling

- Task 1a: Specify the cardinalities in Modified Chen notation (8 Points)
  - A hospital employs at least 4 nurses and has at least 8 patient rooms.
  - A nurse works in exactly one hospital and treats up to 16 patients.
  - A patient is treated by at least one but potentially many nurses.
  - Every patient has a room, a rooms belongs to exactly one hospital, and rooms are never shared by multiple patients.

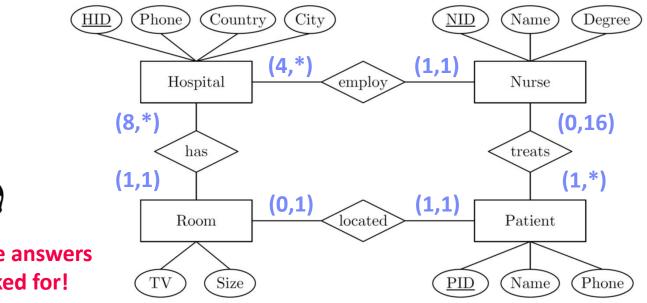






# #1 Data Modeling, cont.

- Task 1b: Specify the cardinalities in (min, max) notation (4 Points)
  - A hospital employs at least 4 nurses and has at least 8 patient rooms.
  - A nurse works in exactly one hospital and treats up to 16 patients.
  - A patient is treated by at least one but potentially many nurses.
  - Every patient has a room, a rooms belongs to exactly one hospital, and rooms are never shared by multiple patients.



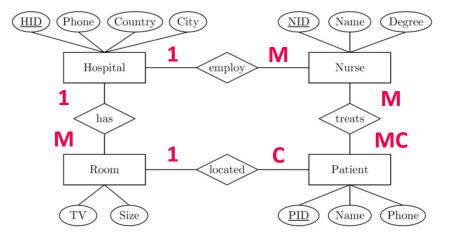






## #1 Data Modeling, cont.

- Task 1c: Map the given ER diagram into a relational schema (10 points)
  - Including data types, primary keys, and foreign keys



```
    Hospitals(
        HID:int, phone:char(16), Country:varchar(64), City:varchar(64))
    Nurses(
        NID:int, Name:varchar(64), Degree:varchar(32), HID<sup>FK</sup>:int)
    Patient(
        PID:int, Name:varchar(64), Phone:char(16), RID<sup>FK</sup>:int)
    Room(
        RID:int, TV:boolean, Size:int, HID<sup>FK</sup>:int)
    Treated(
        NID<sup>FK</sup>:int, PID<sup>FK</sup>:int)
```





# #1 Data Modeling, cont.

- Task 1d: Bring your schema in 3<sup>rd</sup> normal form and explain why it is in 3NF (12 points)
  - Let Hospital.Phone and Patient.Phone be multi-valued attributes
  - Assume the functional dependency City → Country

- Phones(Number:char(16), HID<sup>FK</sup>:int, PID<sup>FK</sup>:int)
- Cities(City:varchar(64), Country:varchar(64))
- Hospitals(HID:int, City<sup>FK</sup>:varchar(64))
- 1st Normal Form: no multi-valued attributes
- 2<sup>nd</sup> Normal Form: 1NF + all non-key attributes fully functional dependent on PK
- 3<sup>rd</sup> Normal Form: 2NF + no dependencies among non-key attributes





# #2 Structured Query Language

Task 2a: Compute the results for the following queries (15 points)

#### **Orders**

OID	Customer	Date	Quantity	PID
1	A	'2019-06-25'	3	2
2	В	'2019-06-25'	1	3
3	A	'2019-06-25'	1	4
4	C	'2019-06-26'	2	2
5	D	'2019-06-26'	1	4
6	C	'2019-06-26'	1	1

**Q1: SELECT DISTINCT** Customer, Date

FROM Orders O, Products P

WHERE O.PID = P.PID AND Name IN('Y', 'Z')

Q2: SELECT Customer, count(\*) FROM Orders

**GROUP BY** Customer

ORDER BY count(\*) DESC, Customer ASC

Q3: SELECT Customer, sum(0.Quantity \* P.Price)

FROM Orders O, Products P

WHERE O.PID = P.PID

**GROUP BY** Customer

#### **Products**

PID	Name	Price
1	X	100
2	Y	15
4	Z	75
3	W	120

Customer	Date
Α	'2019-06-25'
С	'2019-06-26'
D	'2019-06-26'

Customer	Sum	
А	2	
С	2	
В	1	
D	1	

Customer	Sum
Α	120
В	120
С	130
D	75





# #2 Structured Query Language, cont.

**Orders** 

 Task 2b: Write SQL queries to answer the following Qs (15 points)

OID	Customer	Date	Quantity	PID
1	A	'2019-06-25'	3	2
2	В	'2019-06-25'	1	3
3	A	'2019-06-25'	1	4
4	C	'2019-06-26'	2	2
5	D	'2019-06-26'	1	4
6	С	'2019-06-26'	1	1

#### Products

PID	Name	Price
1	X	100
2	Y	15
4	Z	75
3	W	120

Q4: Which products where bought on 2019-06-25 (return the distinct product names)?

Q5: Which customers placed only one order?

Q6: How much revenue (sum( O.Quantity \* P.Price)) did products with a price less then 90 generate (return (product name, revenue))? FROM Orders O, Products P
WHERE O.PID = P.PID
AND Date = '2019-06-25'

SELECT Customer FROM Orders
GROUP BY Customer HAVING count(\*) = 1

FROM Orders O, Products P
WHERE O.PID = P.PID AND Price < 90
GROUP BY P.Name</pre>



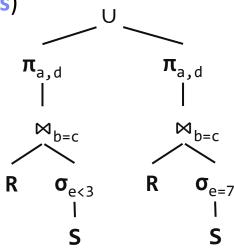


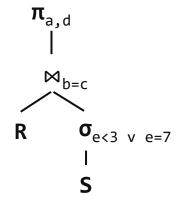
# **#3 Query Processing**

 Task 3a: Assume tables R(a,b), and S(c,d,e), draw a logical query tree in relational algebra for the following query: (5 points)

```
Q7: SELECT R.a, S.d FROM R, S
WHERE R.b = S.c AND S.e < 3
UNION ALL
SELECT R.a, S.d FROM R, S
WHERE R.b = S.c AND S.e = 7
```

 Task 3b: Draw an optimized logical query tree for the above query in relational algebra by eliminating the union operation (3 points)







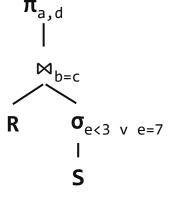


# #3 Query Processing, cont.

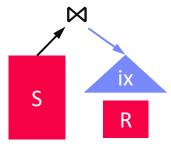
Task 3c: Given the schema and query above, which attribute or attributes are good candidates for secondary indexes and how could they be exploited during query processing? (4 points)

- Solution
  - S.e → index scan (lookup e=7, lookup e=3 and scan DESC)

table data



■ R.b (or S.c) → index nested loop join (for every S tuple s, loopup s.c in IX)







# #3 Query Processing, cont.

 Task 3d: Describe the volcano (open-next-close) iterator model by example of a selection operator and discuss the space complexity of this selection operator. (6 points)

- Open, next, close calls propagate from root to leafs
- Open: operator initialization
- Next: compute next tuple (selection: call next of input until next qualifying tuple found)
- Close: cleanup resources
- Space complexity: O(1)

```
void open() { R.open(); }

void close() { R.close(); }

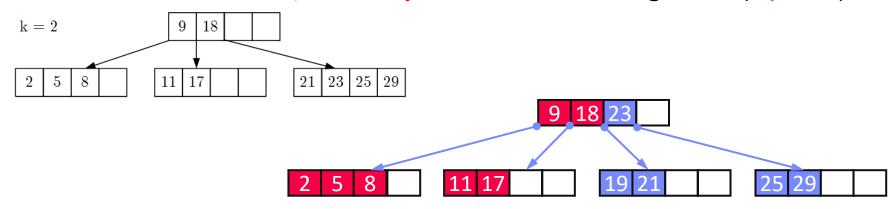
Record next() {
  while( (r = R.next()) != EOF )
    if( p(r) ) //A==7
      return r;
  return EOF;
}
```



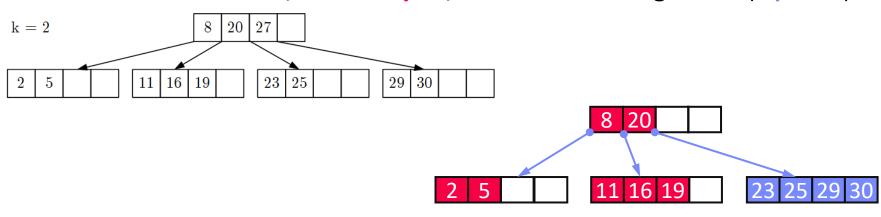


# #4 Physical Design – B-Trees

Task 4a: Given B-tree, insert key 19 and draw resulting B-tree (7 points)



■ Task 4b: Given B-tree, delete key 27, and draw resulting B-tree (8 points)



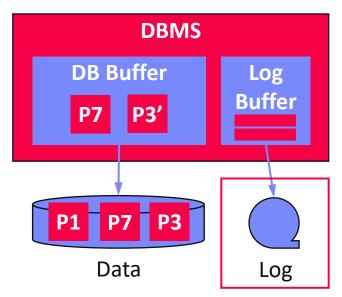




# **#5 Transaction Processing**

 Task 5a: Describe the concept of a database transaction log, and explain how it relates to the ACID properties Atomicity and Durability (7 points)

- Log: append-only TX changes, often on separate devices
- Write-ahead logging (log written before DB, forced-log on commit)
- Recovery: forward (REDO) and backward (UNDO) processing



- #1 Atomicity: A TX is executed atomically (completely or not at all); on failure/aborts no changes in DB (UNDO)
- #2 Durability: Guaranteed persistence of changes of successful TXs; in case of system failures, the database is recoverable (REDO)





### #6 NoSQL

Task 6a: Describe the concept and system architecture of a key-value store, including techniques for achieving high write throughput, and scale-out in distributed environments. Please focus specifically on aspects of physical design such as index structures, and distributed data storage. (10 points)

### Solution

- KV store: simple map of key-value pairs,
   w/ get/put interface, often distributed
- Index structure for high write throughput: Log-structured merge trees (LSM)
- Distributed data storage for scale-out:
   horizontal partitioning (sharding) via hash or range partitioning,
   partitioning via selection, reconstruction via union

eventual consistency for high availability and partition tolerance





# Remaining Questions & Answers

Course Content

Data Management in general





### Conclusions and Q&A

- Summary
  - 13 Data Stream Processing Systems
  - 14 Q&A and Exam Preparation
- Next Week: NO lecture (use time for exam prep)
  - Office hours Mo 3pm as usual
- Exams
  - Jan 30, 5.30pm Exam DM VO / DB VU, HS i13
  - Jan 31, 5.30pm Exam DM VO / DB VU, HS i13
  - Feb 6, 4pm Exam DM VO / DB VU, HS i13 (also as replacement exam)

