Data Integration and Analysis
11 Distributed Data-Parallel Computation

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Announcements/Org

- **#1 Video Recording**
  - Link in **TeachCenter** & **TUbe** (lectures will be public)
- **#2 DIA Projects**
  - **13 Projects** selected (various topics)
  - **3 Exercises** selected (distributed data deduplication)
  - SystemDS: apps into `./scripts/staging/<your_project>`
- **#3 Exam**
  - **Feb 3, 1pm – Feb 5, 2pm**, remote exam possible
  - Oral exam, **45min slots**, first-come, first-serve
  - https://doodle.com/poll/ikzsffek2vhd85q4
- **#4 Course Evaluation**
  - Evaluation time frame: **Jan 14 – Feb 14** → feedback

If problems, **please ask for help**

Course Outline Part B:
Large-Scale Data Management and Analysis

12 Distributed Stream Processing [Jan 24]

13 Distributed Machine Learning Systems [Jan 31]

Compute/ Storage

11 Distributed Data-Parallel Computation [Jan 17]

10 Distributed Data Storage [Jan 10]

Infra

09 Cloud Resource Management and Scheduling [Dec 13]

08 Cloud Computing Fundamentals [Dec 06]

Agenda

- **Motivation and Terminology**
- **Data-Parallel Collection Processing**
- **Data-Parallel DataFrame Operations**
- **Data-Parallel Computation in SystemDS**

Motivation and Terminology

Recap: Central Data Abstractions

- **#1 Files and Objects**
  - **File:** Arbitrarily large sequential data in specific file format (CSV, binary, etc)
  - **Object:** binary large object, with certain meta data

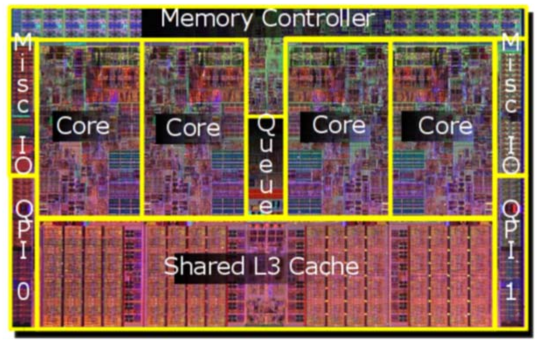
- **#2 Distributed Collections**
  - Logical multi-set (**bag**) of **key-value pairs** (**unsorted collection**)
  - Different physical representations
  - **Easy distribution** of pairs via horizontal partitioning (aka shards, partitions)
  - Can be created from single file, or directory of files (unsorted)

Key	Value
4	Delta
2	Bravo
1	Alpha
3	Charlie
5	Echo
6	Foxtrott
7	Golf

Excursus: Nehalem Architecture

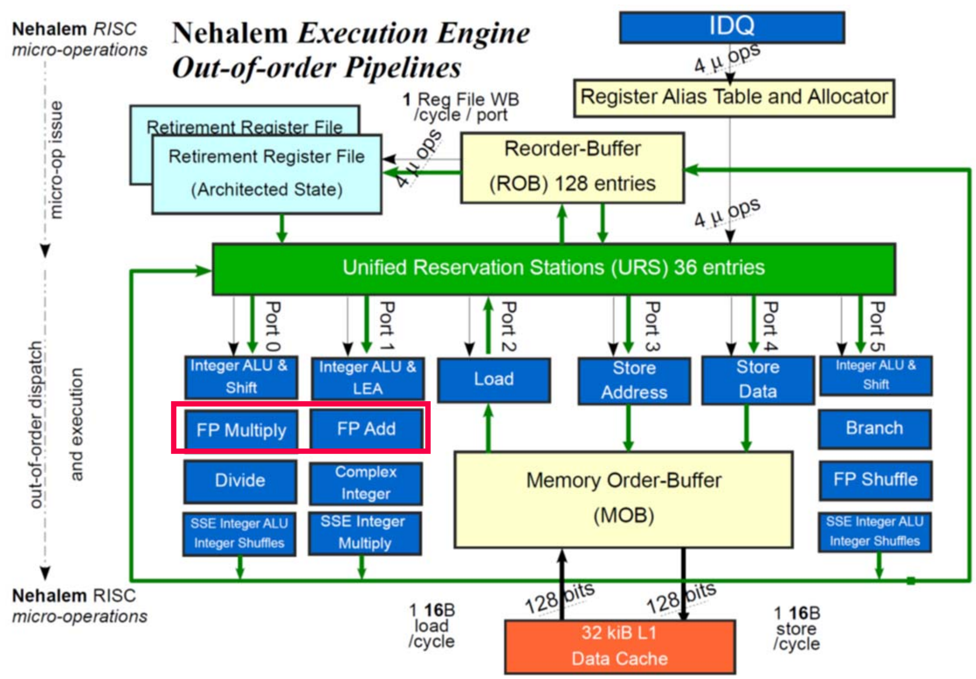
[Michael E. Thomadakis: The Architecture of the Nehalem Processor and Nehalem-EP SMP Platforms, Report, 2010]

- **Multi-core CPU**
  - 4 core w/ hyper-threading
  - **Per core:** L1i/L1d, L2 cache
  - **Per CPU:** L3 cache (8MB)
  - 3 memory channels (8B width, max 1.333Ghz)



QPI … Quick Path Interconnect

- **Pipeline**
  - **Frontend:** Instruction Fetch, Pre-Decode, and Decode
  - **Backend:** Rename/Allocate, Scheduler, Execute
- **Out-of-Order Execution Engine** (IPC=4)
  - 128b FP Multiply
  - 128b FP Add



Terminology

- **Flynn's Classification**
  - SISD, SIMD
  - (MISD), MIMD

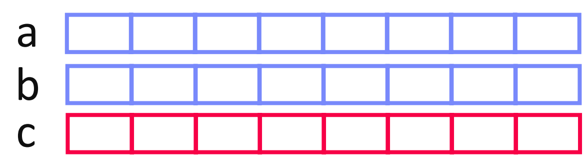
	Singe Data	Multiple Data
Singe Instruction	**SISD** (uni-core)	**SIMD** (vector)
Multiple Instruction	**MISD** (pipelining)	**MIMD** (multi-core)

[Michael J. Flynn, Kevin W. Rudd: Parallel Architectures. **ACM Comput. Surv. 28(1) 1996**]

- **Example: SIMD Processing**
  - Streaming SIMD Extensions (SSE)
  - Process the same operation on multiple elements at a time (**packed** vs scalar SSE instructions)
  - Data parallelism (aka: instruction-level parallelism)
  - Example: **`VFMADD132PD`**

2009 Nehalem: **128b** (2xFP64)
2012 Sandy Bridge: **256b** (4xFP64)
2017 Skylake: **512b** (8xFP64)

```
c = _mm512_fmadd_pd(a, b);
```



Terminology cont.

- **Distributed, Data-Parallel Computation**

```
Y = X.map(x -> foo(x))
```

  - Parallel computation of function foo() ➔ **single instruction**
  - Collection X of data items (key-value pairs) ➔ **multiple data**
  - Data parallelism similar to **SIMD** but more coarse-grained notion of "instruction" and "data" ➔ **SPMD** (single program, multiple data)

[Frederica Darema: The SPMD Model : Past, Present and Future. **PVM/MPI 2001**]

- **Additional Terminology**
  - BSP: Bulk Synchronous Parallel (global barriers)
  - ASP: Asynchronous Parallel (no barriers, often with accuracy impact)
  - SSP: Stale-synchronous parallel (staleness constraint on fastest-slowest)
  - Other: Fork&Join, Hogwild!, event-based, decentralized

- **Beware:** data parallelism used in very different contexts (e.g., Param Server)

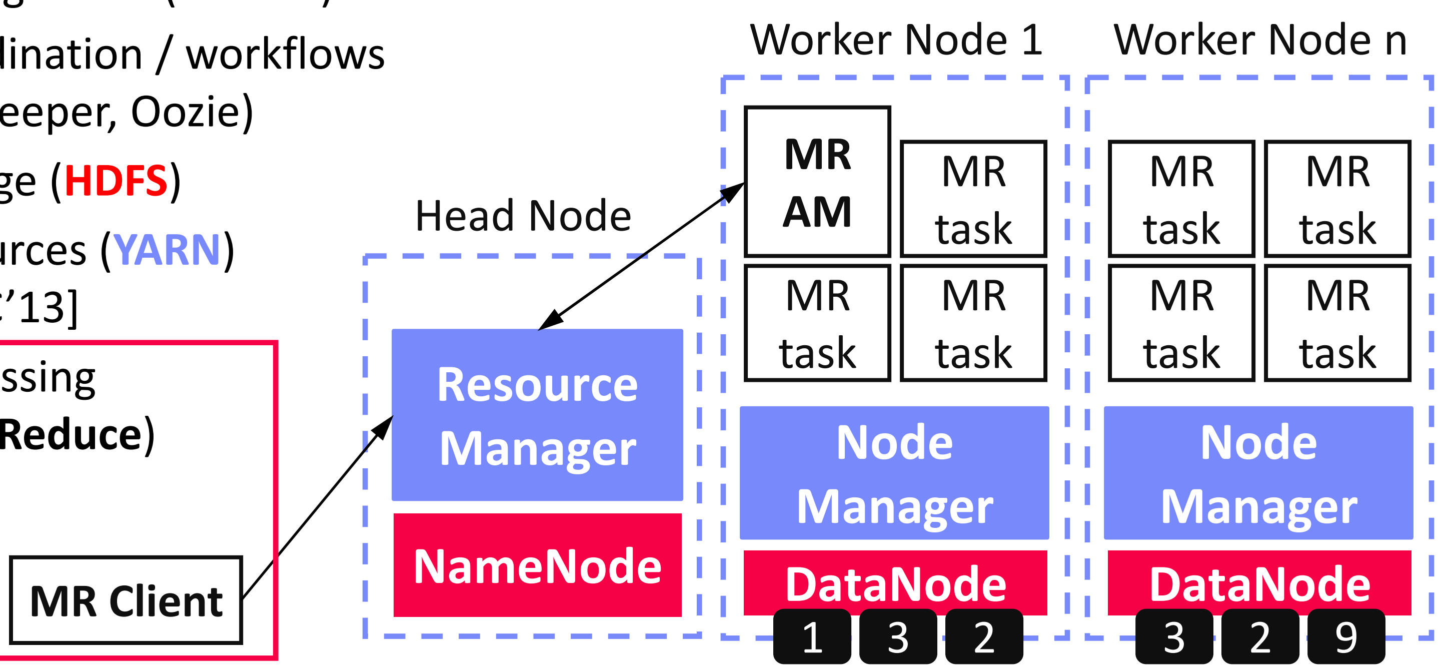
Data-Parallel Collection Processing

Hadoop History and Architecture

- **Recap: Brief History**
  - Google's GFS [SOSP'03] + MapReduce
    → Apache Hadoop (2006)
  - Apache Hive (SQL), Pig (ETL), Mahout (ML), Giraph (Graph)

[Jeffrey Dean, Sanjay Ghemawat: MapReduce: Simplified Data Processing on Large Clusters. **OSDI 2004**]

- **Hadoop Architecture / Eco System**
  - Management (Ambari)
  - Coordination / workflows (Zookeeper, Oozie)
  - Storage (**HDFS**)
  - Resources (YARN) [SoCC'13]
  - Processing (**MapReduce**)



MapReduce – Programming Model

- **Overview Programming Model**
  - Inspired by functional programming languages
  - **Implicit parallelism** (abstracts distributed storage and processing)
  - **Map** function: key/value pair → set of intermediate key/value pairs
  - **Reduce** function: merge all intermediate values by key

- **Example** `SELECT Dep, count(*) FROM csv_files GROUP BY Dep`

Name	Dep
X	CS
Y	CS
A	EE
Z	CS

Collection of key/value pairs

```
map(Long pos, String line) {
  parts ← line.split(“,”)
  emit(parts[1], 1)
}
```

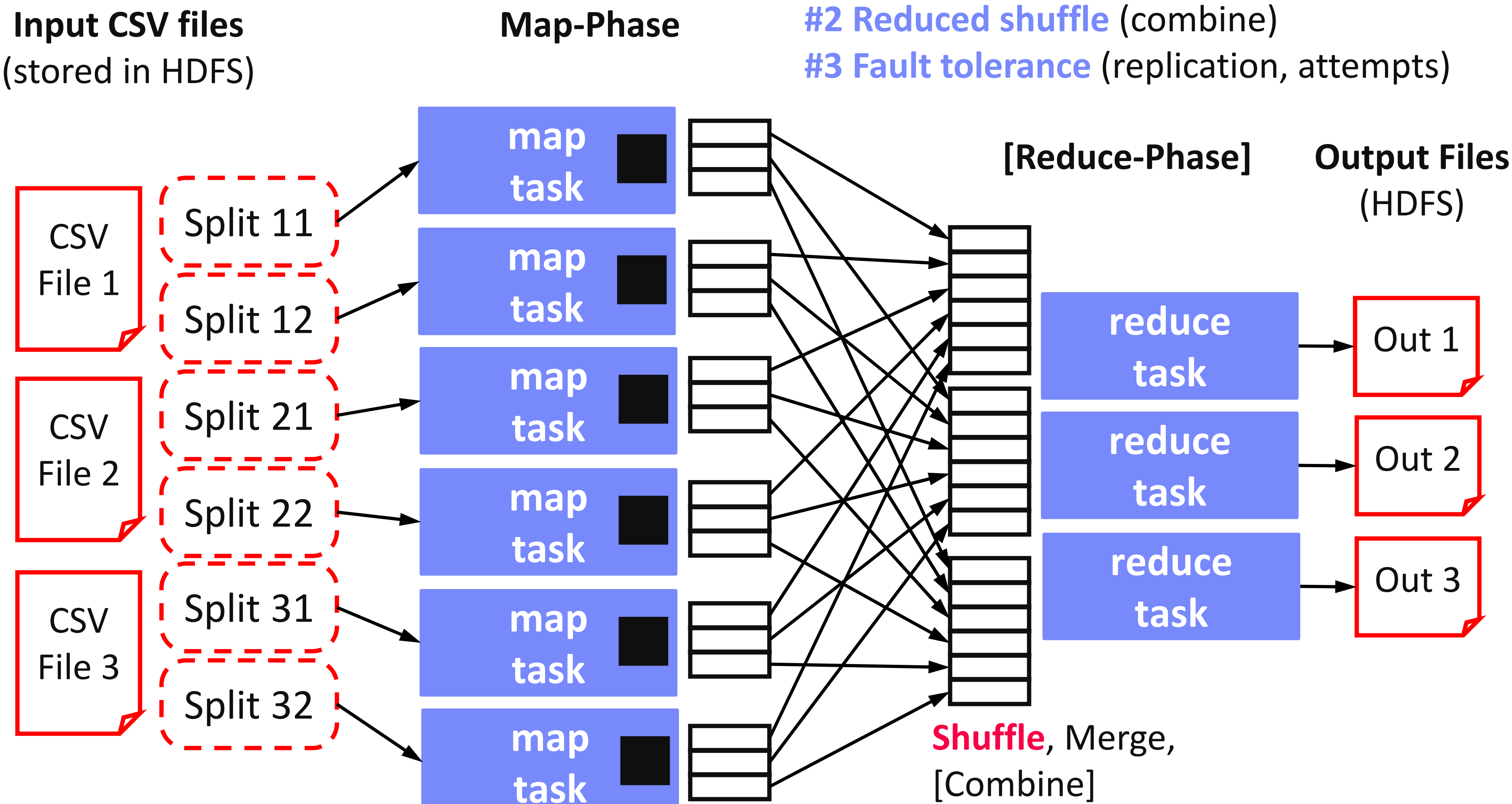
CS	1
CS	1
EE	1
CS	1

```
reduce(String dep,
       Iterator<Long> iter) {
  total ← iter.sum();
  emit(dep, total)
}
```

CS	3
EE	1

MapReduce – Execution Model

#1 Data Locality (delay sched., write affinity)
#2 Reduced shuffle (combine)
#3 Fault tolerance (replication, attempts)



Sort, [**Combine**], [Compress]

w/ #reducers = 3

MapReduce – Query Processing

- **Basic Unary Operations**
  - Selections (brute-force), projections
  - Ordering (e.g., **TeraSort**): Sample, pick k quantiles; shuffle-based partition sort
  - Additive and semi-additive aggregation with grouping, distinct
- **Binary Operations**
  - Set operations
    (union, intersect, difference) and joins
  - Different physical operators for R $\bowtie$ S
    - **Broadcast join**: broadcast S, build HT S, map-side HJOIN
    - **Repartition join**: shuffle (repartition) R and S, reduce-side MJOIN
    - **Improved repartition join:** avoid buffering via key-tag sorting
    - **Directed join** (pre/co-partitioned): map-only, R input, S read side-ways

[Spyros Blanas et al.: A comparison of join algorithms for log processing in MapReduce. **SIGMOD 2010**]

- **Hybrid SQL-on-Hadoop Systems** [VLDB'15]
  - E.g.: Hadapt (HadoopDB), Impala, IBM BigSQL, Presto, Drill, Actian

Spark History and Architecture

- **Summary MapReduce**
  - Large-scale & fault-tolerant processing w/ UDFs and files → **Flexibility**
  - Restricted functional APIs → **Implicit parallelism and fault tolerance**
  - **Criticism**: #1 **Performance**, #2 **Low-level APIs**, #3 **Many different systems**

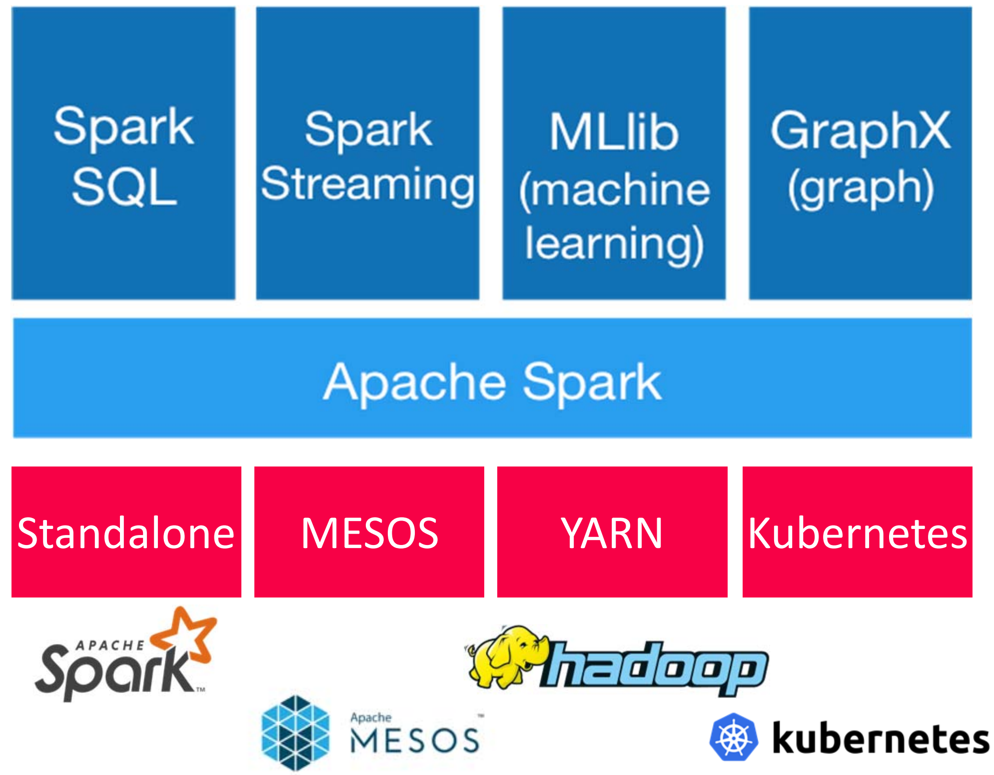
- **Evolution to Spark** (and Flink)
  - Spark [HotCloud'10] + RDDs [NSDI'12] → **Apache Spark** (2014)
  - **Design: standing executors with in-memory storage**, lazy evaluation, and fault-tolerance via RDD lineage
  - **Performance:** In-memory storage and fast job scheduling (100ms vs 10s)
  - **APIs:** Richer functional APIs and general computation DAGs, high-level APIs (e.g., DataFrame/Dataset), unified platform

→ **But many shared concepts/infrastructure**
  - **Implicit parallelism through dist. collections** (data access, fault tolerance)
  - Resource negotiators (YARN, Mesos, Kubernetes)
  - HDFS and object store connectors (e.g., Swift, S3)

Spark History and Architecture, cont.

- **High-Level Architecture**
  - **Different language bindings**: Scala, Java, Python, R
  - **Different libraries**: SQL, ML, Stream, Graph
  - Spark core (incl RDDs)
  - **Different cluster managers**: Standalone, Mesos, **Yarn**, **Kubernetes**
  - Different file systems/ formats, and data sources: **HDFS**, **S3**, SWIFT, **DBs**, **NoSQL**

[https://spark.apache.org/]



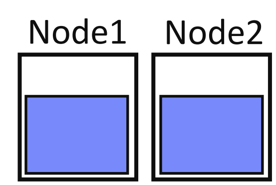
- **Focus on a unified platform for data-parallel computation** (Apache Flink w/ similar goals)

Spark Resilient Distributed Datasets (RDDs)

- **RDD Abstraction** `JavaPairRDD<MatrixIndexes,MatrixBlock>`
  - **Immutable**, partitioned **collections of key-value pairs**
  - **Coarse-grained** deterministic operations (transformations/actions)
  - Fault tolerance via lineage-based re-computation
- **Operations**
  - Transformations: define new RDDs
  - Actions: return result to driver

Type	Examples
Transformation (**lazy**)	`map, hadoopFile, textFile, flatMap, filter, sample, join, groupByKey, cogroup, reduceByKey, cross, sortByKey, mapValues`
Action	`reduce, save, collect, count, lookupKey`

- **Distributed Caching**
  - Use fraction of worker **memory for caching**
  - Eviction at granularity of individual partitions
  - **Different storage levels** (e.g., mem/disk x serialization x compression)



Spark Resilient Distributed Datasets (RDDs), cont.

- **Lifecycle of an RDD**
  - Note: can't broadcast an RDD directly

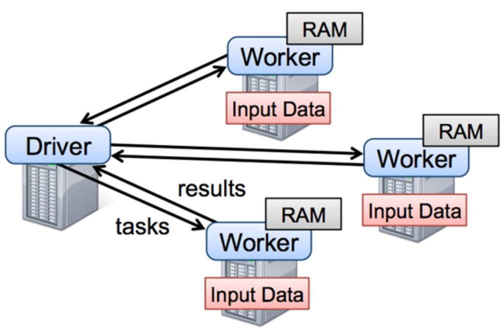
```
X.filter(foo())
X.mapValues(foo())
X.reduceByKey(foo())
X.cache()
```

Local Data (value, collection) → `sc.parallelize(lst)` → **Distributed Collection**

Distributed Collection → `lst = X.collect()`, `v = X.reduce(foo())` → **Local Data**

Distributed Collection → `X.saveAsObjectFile(f)`, `X.saveAsTextFile(f)` → **File on DFS**

File on DFS → `sc.hadoopFile(f)`, `sc.textFile(f)` → **Distributed Collection**



Spark Partitions and Implicit/Explicit Partitioning

- **Spark Partitions**
  - Logical key-value collections are split into **physical partitions**
  - Partitions are granularity of **tasks, I/O, shuffling, evictions**

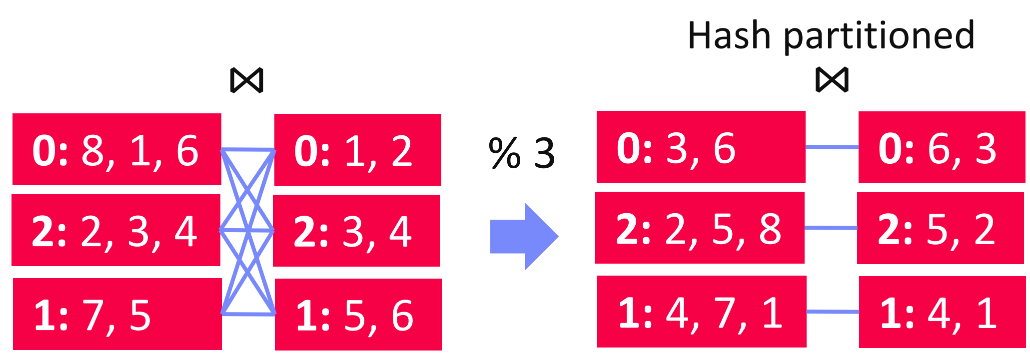
~128MB

- **Partitioning via Partitioners**
  - Implicitly on every data shuffling
  - Explicitly via `R.repartition(n)`

Example Hash Partitioning:
For all (k,v) of R:
pid = hash(k) % n

- **Partitioning-Preserving**
  - All operations that are guaranteed to keep keys unchanged (e.g. `mapValues()`, `mapPartitions()` w/ preservesPart flag)

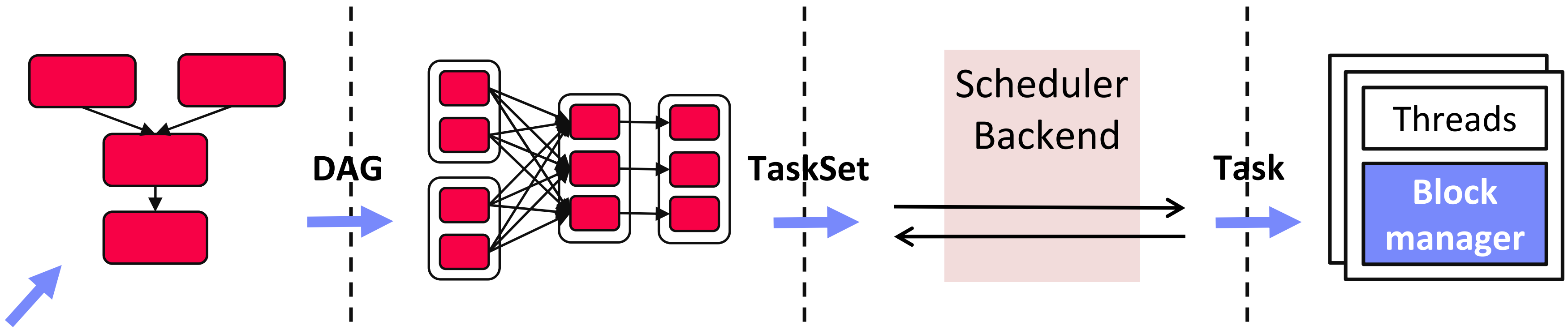
- **Partitioning-Exploiting**
  - Join: R3 = R1.join(R2)
  - Lookups:
    `v = C.lookup(k)`



Spark Scheduling Process

[Tilmann Rabl: Big Data Systems, HPI WS2019/20]

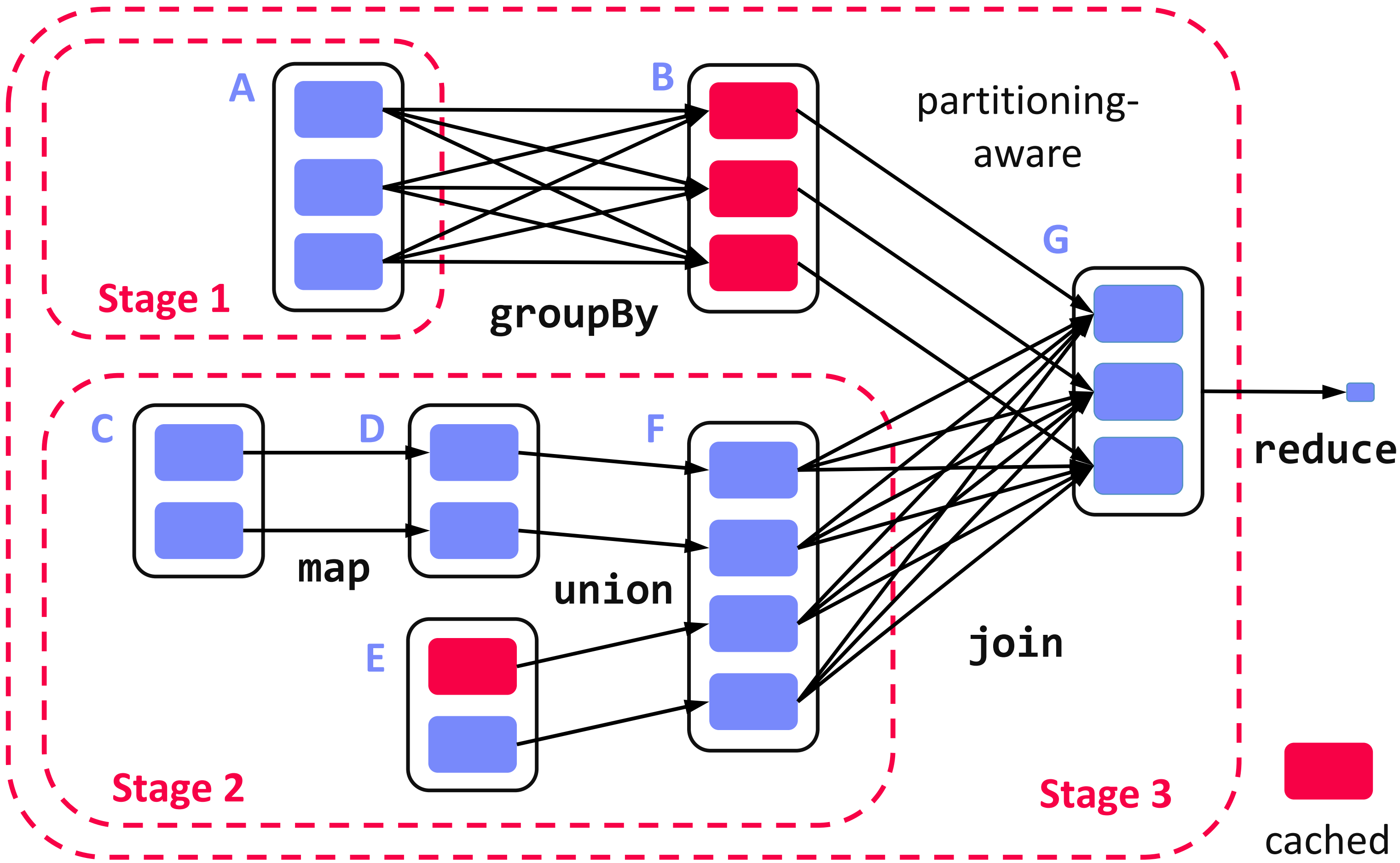
RDD Objects | **DAGScheduler** | **TaskScheduler** | **Workers**




```
rdd1.join(rdd2)
    .reduceByKey(…)
    .filter(…)
```

RDD Objects	DAGScheduler	TaskScheduler	Workers
build operator DAG	split graph into stages of tasks	launch tasks at workers	execute tasks
	submit each stage as ready	retry failed or straggling tasks	store and serve blocks

Spark Lazy Evaluation, Caching, and Lineage

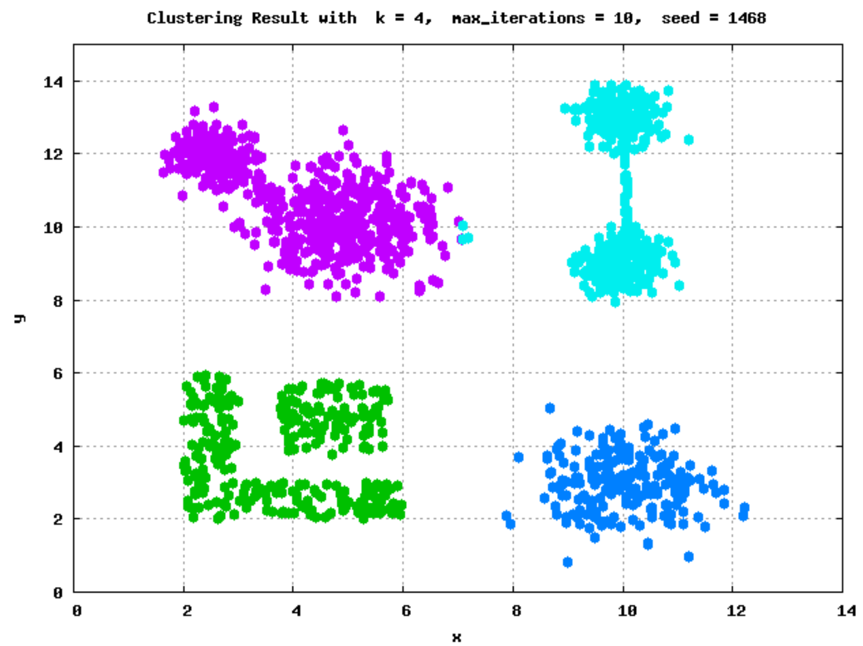


[Matei Zaharia, Mosharaf Chowdhury, Tathagata Das, Ankur Dave, Justin Ma, Murphy McCauly, Michael J. Franklin, Scott Shenker, Ion Stoica: Resilient Distributed Datasets: A Fault-Tolerant Abstraction for In-Memory Cluster Computing. **NSDI 2012**]

Example: k-Means Clustering

- **k-Means Algorithm**
  - Given dataset D and number of clusters k, find cluster centroids ("mean" of assigned points) that minimize within-cluster variance
  - Euclidean distance: **sqrt**(**sum**((**a**-**b**)^2))
- **Pseudo Code**

```
function Kmeans(D, k, maxiter) {
  C' = randCentroids(D, k);
  C = {};
  i = 0; //until convergence
  while( C' != C & i<=maxiter ) {
    C = C';
    i = i + 1;
    A = getAssignments(D, C);
    C' = getCentroids(D, A, k);
  }
  return C'
}
```



Example: K-Means Clustering in Spark

```
// create spark context (allocate configured executors)
JavaSparkContext sc = new JavaSparkContext();

// read and cache data, initialize centroids
JavaRDD<Row> D = sc.textFile("hdfs:/user/mboehm/data/D.csv")
  .map(new ParseRow()).cache(); // cache data in spark executors
Map<Integer,Mean> C = asCentroidMap(D.takeSample(false, k));

// until convergence
while( !equals(C, C2) & i<=maxiter ) {
  C2 = C; i++;
  // assign points to closest centroid, recompute centroid
  Broadcast<Map<Integer,Row>> bC = sc.broadcast(C)
  C = D.mapToPair(new NearestAssignment(bC))
       .foldByKey(new Mean(0), new IncComputeCentroids())
       .collectAsMap();
}

return C;
```

Note: Existing library algorithm
[https://github.com/apache/spark/blob/master/mllib/src/main/scala/org/apache/spark/mllib/clustering/KMeans.scala]

Data-Parallel DataFrame Operations

Origins of DataFrames

- **Recap: Data Preparation Problem**
  - **80% Argument:** 80-90% time for finding, integrating, cleaning data
  - Data scientists prefer scripting languages and in-memory libraries
- **R and Python DataFrames**
  - R `data.frame/dplyr` and Python `pandas` DataFrame for seamless data manipulations (most popular packages/features)
  - DataFrame: **table with a schema**
  - Descriptive stats and basic math, reorganization, joins, grouping, windowing
  - **Limitation:** Only in-memory, single-node operations
- **Example Pandas**

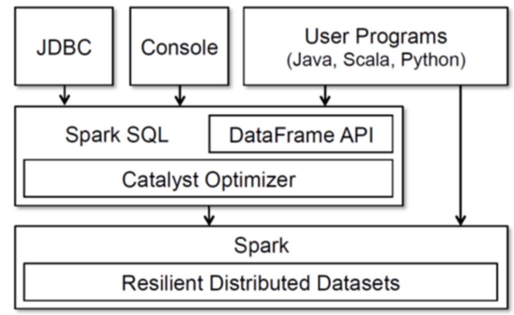
```
import pandas as pd

df = pd.read_csv('data/tmp1.csv')
df.head()

df = pd.concat(df, df[['A', 'C']], axis=0)
```

Spark DataFrames and DataSets

- **Overview Spark DataFrame**
  - DataFrame is **distributed collection of rows** with named/typed columns
  - **Relational operations** (e.g., projection, selection, joins, grouping, aggregation)
  - **DataSources** (e.g., json, jdbc, parquet, hdfs, s3, avro, hbase, csv, cassandra)



- **DataFrame and Dataset APIs**

```
DataFrame = Dataset[Row]
```

  - DataFrame was introduced as basis for Spark SQL
  - DataSets allow **more customization** and compile-time analysis errors (Spark 2)

- **Example DataFrame**

```
logs = spark.read.format("json").open("s3://logs")
logs.groupBy(logs.user_id).agg(sum(logs.time))
  .write.format("jdbc").save("jdbc:mysql//...")
```

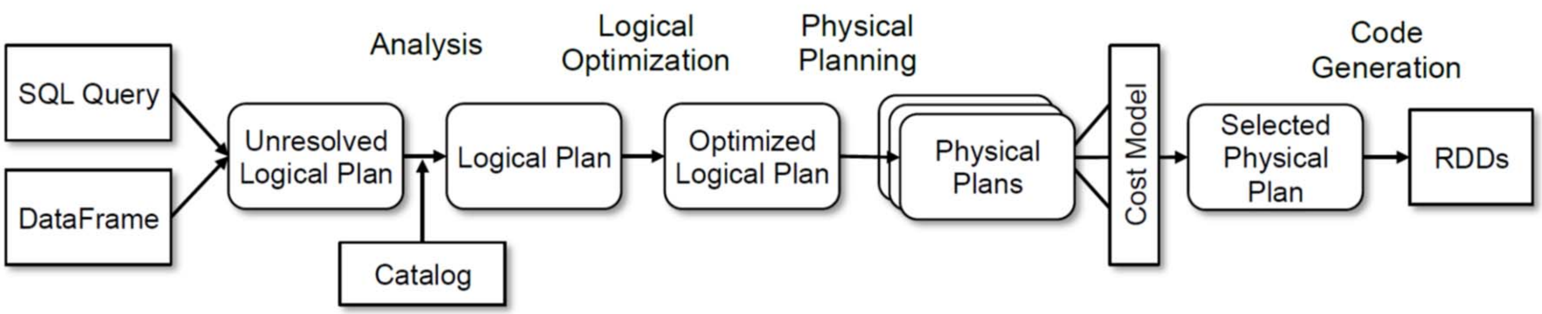
[Michael Armbrust: Structuring Apache Spark – SQL, DataFrames, Datasets, and Streaming, **Spark Summit 2016**]

→ **PySpark**

SparkSQL and DataFrame/Dataset

[Michael Armbrust et al.: Spark SQL: Relational Data Processing in Spark. **SIGMOD 2015**]

- **Overview SparkSQL**
  - Shark (~2013): academic prototype for SQL on Spark
  - SparkSQL (~2015): reimplementation from scratch
  - Common IR and compilation of SQL and DataFrame operations
- **Catalyst: Query Planning**

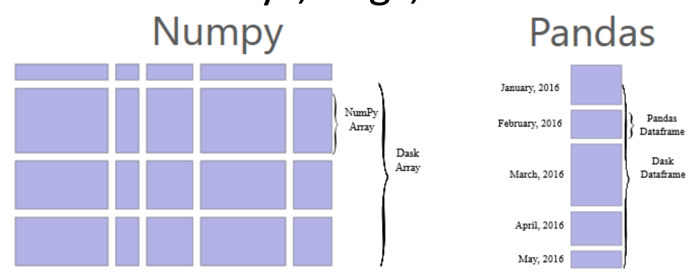


- **Performance features**
  - **#1** Whole-stage code generation via Janino
  - **#2** Off-heap memory (`sun.misc.Unsafe`) for caching and certain operations
  - **#3** Pushdown of selection, projection, joins into data sources (+ join ordering)

Dask DASK

[Matthew Rocklin: Dask: Parallel Computation with Blocked algorithms and Task Scheduling, **Python in Science 2015**]
[Dask Development Team: Dask: Library for dynamic task scheduling, 2016, https://dask.org]

- **Overview Dask**
  - Multi-threaded and distributed operations for arrays, bags, and dataframes
  - `dask.array:` list of numpy n-dim arrays
  - `dask.dataframe:` list of pandas data frames
  - `dask.bag:` unordered list of tuples (second order functions)
  - Local and distributed schedulers: threads, processes, YARN, Kubernetes, containers, HPC, and cloud, GPUs



- **Execution**
  - Lazy evaluation
  - Limitation: requires **static size inference**
  - Triggered via `compute()`

```
import dask.array as da

x = da.random.random(
  (10000,10000), chunks=(1000,1000))
y = x + x.T
y.persist() # cache in memory
z = y[::2, 5000:].mean(axis=1)
ret = z.compute() # returns NumPy array
```

Data-Parallel Operations in SystemDS

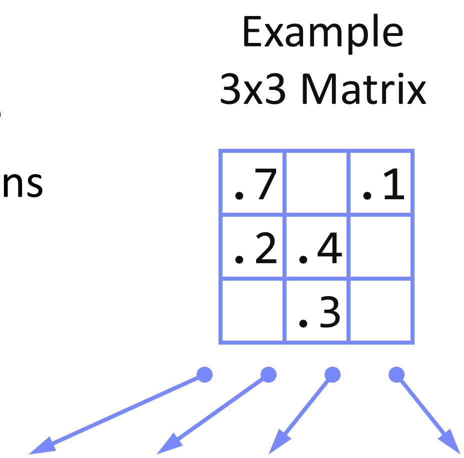
[Matthias Boehm et al.: SystemDS: A Declarative Machine Learning System for the End-to-End Data Science Lifecycle. **CIDR 2020**]

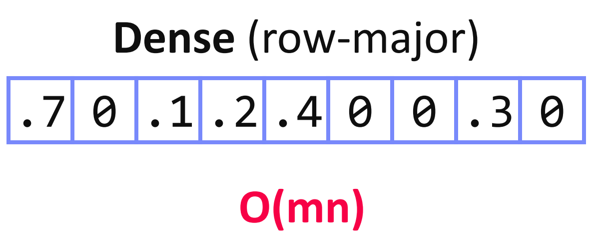
[Matthias Boehm et al.: SystemML: Declarative Machine Learning on Spark. **PVLDB 9(13) 2016**]

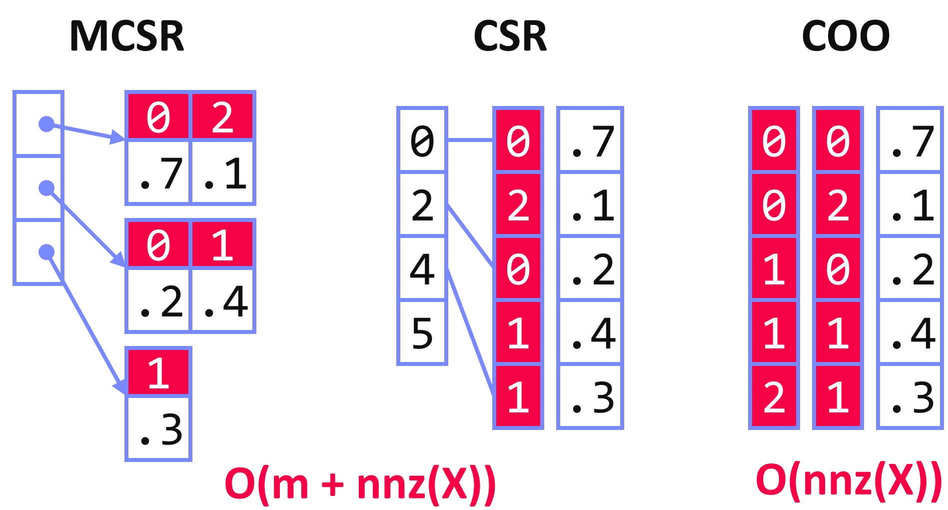
[Amol Ghoting et al.: SystemML: Declarative Machine Learning on MapReduce. **ICDE 2011**]

Background: Matrix Formats

- **Matrix Block** (m x n)
  - A.k.a. tiles/chunks, most operations defined here
  - Local matrix: single block, different representations
- **Common Block Representations**
  - Dense (linearized arrays)
  - MCSR (modified CSR)
  - CSR (compressed sparse rows), CSC
  - COO (Coordinate matrix)







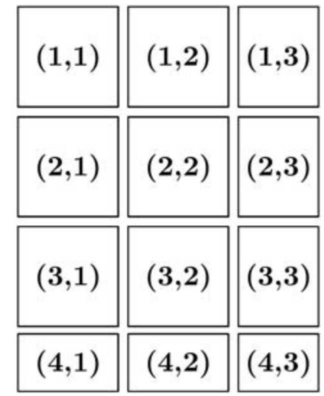
Distributed Matrix Representations

- **Collection of "Matrix Blocks" (and keys)**
  - Bag semantics (duplicates, unordered)
  - Logical (Fixed-Size) Blocking

    + join processing / independence

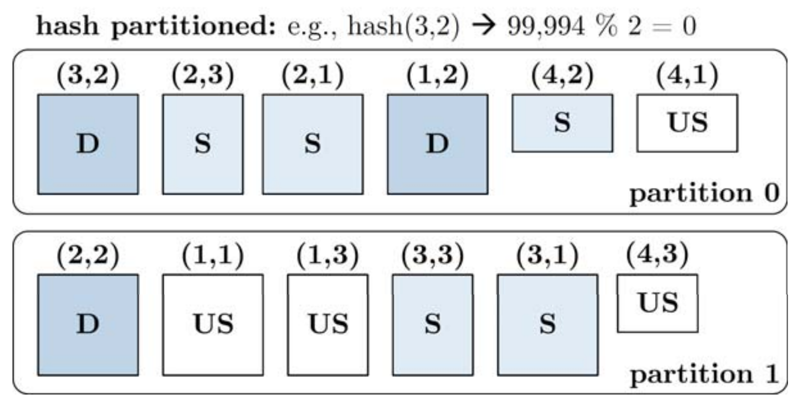
    - (sparsity skew)
  - E.g., SystemML/SystemDS on Spark: `JavaPairRDD<MatrixIndexes,MatrixBlock>`
  - Blocks encoded independently (dense/sparse)

Logical Blocking 3,400x2,700 Matrix (w/ B_c=1,000)



- **Partitioning**
  - Logical Partitioning (e.g., row-/column-wise)
  - Physical Partitioning (e.g., hash / grid)
  - PartitionPruning for Indexing

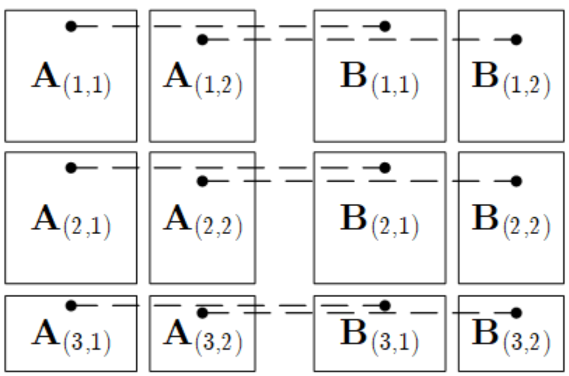
Physical Blocking and Partitioning



Distributed Matrix Operations

Elementwise Multiplication
(Hadamard Product)

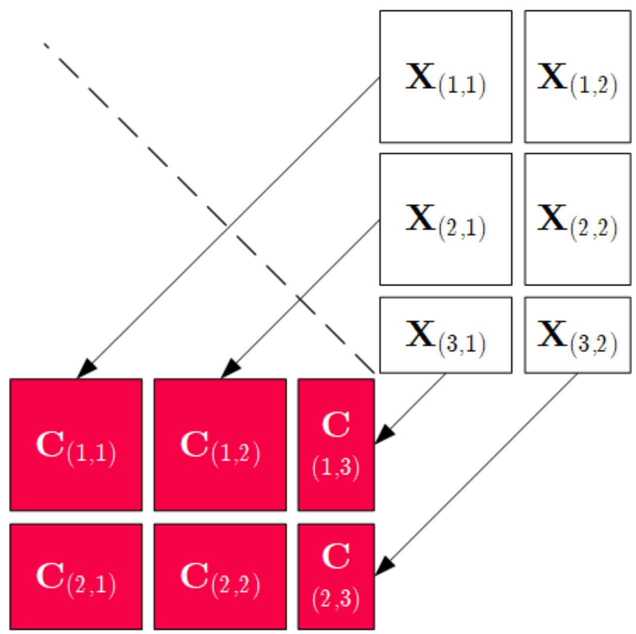
`C = A * B`



Note: also with row/column vector rhs

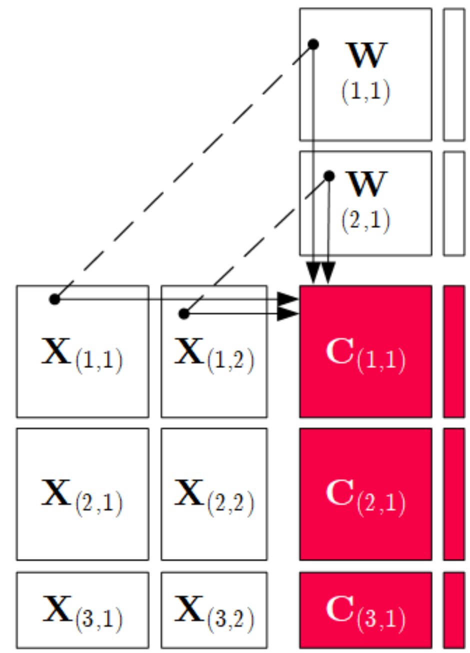
Transposition

`C = t(X)`



Matrix Multiplication

`C = X %*% W`



Note: 1:N join

Partitioning-Preserving Operations

- **Shuffle is major bottleneck for ML on Spark**
- **Preserve Partitioning**
  - Op is partitioning-preserving if keys unchanged (guaranteed)
  - Implicit: Use restrictive APIs (`mapValues()` vs `mapToPair()`)
  - Explicit: Partition computation w/ declaration of partitioning-preserving
- **Exploit Partitioning**
  - Implicit: Operations based on `join`, `cogroup`, etc
  - Explicit: Custom operators (e.g., zipmm)
- **Example: Multiclass SVM**
  - Vectors fit neither into driver nor broadcast
  - $ncol(X) \leq B_c$

repart, chkpt X MEM_DISK

```
parfor(iter_class in 1:num_classes) {
   Y_local = 2 * (Y == iter_class) - 1
   g_old = t(X) %*% Y_local
   ...
   while( continue ) {
      Xd = X %*% s
      ... inner while loop (compute step_sz)
      Xw = Xw + step_sz * Xd;
      out = 1 - Y_local * Xw;
      out = (out > 0) * out;
      g_new = t(X) %*% (out * Y_local) ...
```

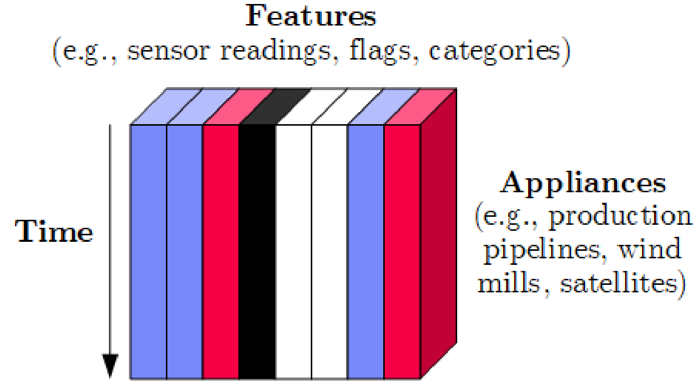
chkpt y_local MEM_DISK

chkpt Xd, Xw MEM_DISK

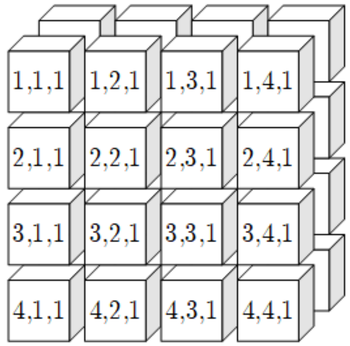
zipmm

SystemDS Data Model: Heterogeneous Tensors

- **Basic Tensor Block**
  - **BasicTensorBlock:** homogeneous tensors (FP32, FP64, INT32, INT64, BOOL, STRING/JSON)
  - **DataTensorBlock:** composed from basic TBs
  - Represents local tensor (CPU/GPU)



- **Distributed Tensor Representation**
  - Collection of fix-sized tensor blocks
  - Squared blocking schemes in n-dim space (e.g., 1024^2, 128^3, 32^4, 16^5, 8^6, 8^7)
  - PairRDD<TensorIndex,TensorBlock>



- **Federated Tensor Representation**
  - Collection of meta data handles to `TensorObjects`, each of which might refer to data on a different worker instance (local or distributed)
  - Generalizes to federated tensors of CPU and GPU data objects

Summary and Q&A

- **Motivation and Terminology**
- **Data-Parallel Collection Processing**
- **Data-Parallel DataFrame Operations**
- **Data-Parallel Computation in SystemDS**

- **Projects and Exercises**
  - **13 projects + 3 exercises**
  - In case of problem: ask for help + problem scaling possible

- **Next Lectures**
  - **12 Distributed Stream Processing** [Jan 24]
  - **13 Distributed Machine Learning Systems** [Jan 31]