

Data Management 10 NoSQL Systems

Matthias Boehm

Graz University of Technology, Austria
Computer Science and Biomedical Engineering
Institute of Interactive Systems and Data Science
BMK endowed chair for Data Management



Announcements/Org

■ #1 Video Recording

- Link in **TUbe** & **TeachCenter** (lectures will be public)
- Optional attendance (independent of COVID)
- **Virtual lectures** (recorded) until end of the year
<https://tugraz.webex.com/meet/m.boehm>



■ #2 Exercise Submissions

- **Exercise 2:** in process of begin graded (→ before Xmas)
- **Exercise 3:** due Dec 21, 11.59pm, provided data
- **Exercise 4:** extra credit, published Dec 11

Q&A

[https://mboehm7.github.io/teaching/ws2122_dbs/T3_data.zip]

■ #3 Course Evaluation and Exam

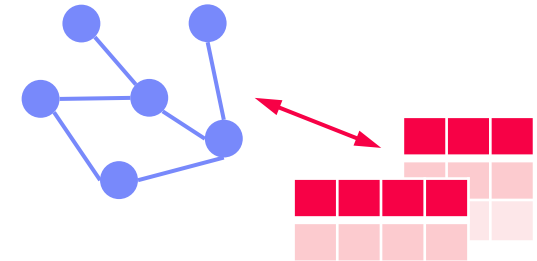
- Evaluation period: **Jan 01 – Feb 15**
- Exam dates: **Feb 04, 12.30pm** and **Feb 04, 5.30pm**
(90+min written exam, start 10 late)



SQL vs NoSQL Motivation

#1 Data Models/Schema

- **Non-relational:** key-value, graph, doc, time series (logs, social media, documents/media, sensors)
- Impedance mismatch / complexity
- **Pay-as-you-go/schema-free** (flexible/implicit)



#2 Scalability

- Scale-up vs simple scale-out
- Horizontal partitioning (sharding) and scaling
- **Commodity hardware, network, disks** (\$)



NoSQL Evolution

- Late 2000s: Non-relational, distributed, open source DBMSs
- Early 2010s: NewSQL: modern, distributed, relational DBMSs
- Not Only SQL: combination with relational techniques

➔ **RDBMS and specialized systems** (consistency/data models)



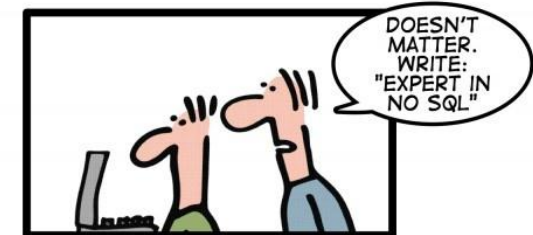
[Credit: <http://nosql-database.org/>]

Agenda

- Consistency and Data Models
- Key-Value Stores
- Document Stores
- Graph Processing
- Time Series Databases
- **Exercise 4: Large-Scale Data Analysis**

Lack of
standards
and imprecise
classification

HOW TO WRITE A CV



Leverage the NoSQL boom

<http://geek-and-poke.com/>



[Wolfram Wingerath, Felix Gessert, Norbert Ritter:
NoSQL & Real-Time Data Management in Research
& Practice. **BTW 2019**]

Consistency and Data Models

Recap: ACID Properties

■ Atomicity

- A transaction is executed atomically (**completely or not at all**)
- If the transaction fails/aborts no changes are made to the database (**UNDO**)

■ Consistency

- A successful transaction ensures that all **consistency constraints are met** (referential integrity, semantic/domain constraints)

■ Isolation

- Concurrent transactions are executed in isolation of each other
- **Appearance of serial transaction execution**

■ Durability

- **Guaranteed persistence** of all changes made by a successful transaction
- In case of system failures, the database is recoverable (**REDO**)

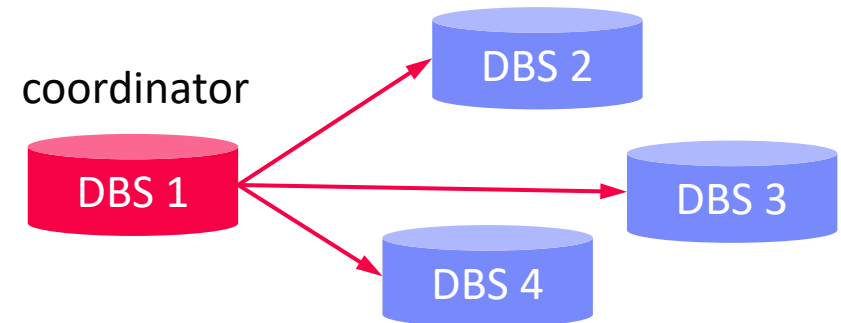
Two-Phase Commit (2PC) Protocol

■ Distributed TX Processing

- N nodes with logically related but physically distributed data (e.g., vertical data partitioning)
- **Distributed TX processing to ensure consistent view** (atomicity/durability)

■ Two-Phase Commit (via $2N$ msgs)

- **Phase 1 PREPARE:** check for successful completion, logging
- **Phase 2 COMMIT:** release locks, and other cleanups
- **Problem: Blocking protocol**



■ Excursus: Wedding Analogy

- Coordinator: marriage registrar
- **Phase 1:** Ask for willingness
- **Phase 2:** If all willing, declare marriage



8 CAP Theorem

Consistency

- **Visibility of updates** to distributed data (atomic or linearizable consistency)
- Different from ACIDs consistency in terms of integrity constraints

Availability

- **Responsiveness** of a services (clients reach available service, **read/write**)

Partition Tolerance

- Tolerance of temporarily **unreachable network partitions**
- System characteristics (e.g., latency) maintained

- **CAP Theorem** *"You can have **AT MOST TWO of these properties for a networked shared-data systems.**"*

[Eric A. Brewer: Towards robust distributed systems (abstract). **PODC 2000**]



Proof

[Seth Gilbert, Nancy A. Lynch: Brewer's conjecture and the feasibility of consistent, available, partition-tolerant web services. **SIGACT News 2002**]



CAP Theorem, cont.

■ CA: Consistency & Availability (ACID single node)

- Network partitions cannot be tolerated
- Visibility of updates (**consistency**) in conflict with **availability** → **no distributed systems**

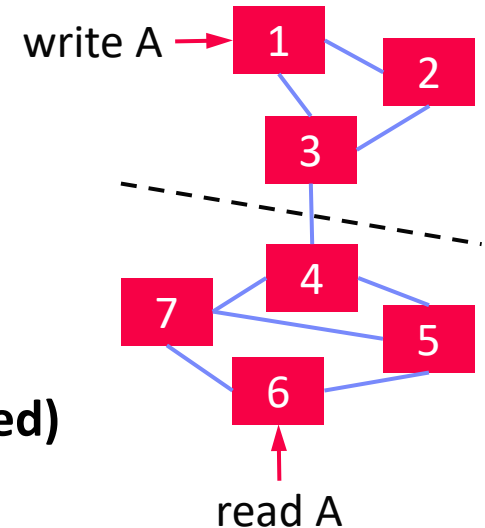
■ CP: Consistency & Partition Tolerance (ACID distributed)

- Availability cannot be guaranteed
- **On connection failure, unavailable**
(wait for overall system to become consistent)

■ AP: Availability & Partition Tolerance (BASE)

- Consistency cannot be guaranteed, use of optimistic strategies
- Simple to implement, main concern: availability to ensure revenue (\$\$\$)

➔ **BASE consistency model**





BASE Properties

- **Basically Available**

- **Major focus on availability**, potentially with outdated data
- No guarantee on global data consistency across entire system

- **Soft State**

- Even without explicit state updates, the **data might change** due to asynchronous propagation of updates and nodes that become available

- **Eventual Consistency**

- Updates eventually propagated, system would reach consistent state if no further updates, and network partitions fixed
- No temporal guarantees on changes are propagated

Eventual Consistency

[Peter Bailis, Ali Ghodsi: Eventual consistency today: limitations, extensions, and beyond. **Commun. ACM** 2013]



Basic Concept

- Changes made to a copy eventually migrate to all
- If update activity stops, replicas will **converge to a logically equivalent state**
- Metric:** time to reach consistency (probabilistic bounded staleness)

Amazon SimpleDB	500ms
Cassandra	200ms
Amazon S3	12s

#1 Monotonic Read Consistency

- After reading data object A, the client never reads an older version

#2 Monotonic Write Consistency

- After writing data object A, it will never be replaced with an older version

#3 Read Your Own Writes / Session Consistency

- After writing data object A, a client never reads an older version

#4 Causal Consistency

- If client 1 communicated to client 2 that data object A has been updated, subsequent reads on client 2 return the new value

Key-Value Stores

Motivation and Terminology

■ Motivation

- **Basic key-value mapping via simple API** (more complex data models can be mapped to key-value representations)
- **Reliability at massive scale on commodity HW** (cloud computing)

■ System Architecture

- **Key**-value maps, where values can be of a variety of data types
- APIs for CRUD operations (create, read, update, delete)
- Scalability via sharding (horizontal partitioning)

users:1:a

“Inffeldgasse 13, Graz”

users:1:b

“[12, 34, 45, 67, 89]”

users:2:a

“Mandellstraße 12, Graz”

users:2:b

“[12, 212, 3212, 43212]”

■ Example Systems

- **Dynamo** (2007, AP) → **Amazon DynamoDB** (2012)
- **Redis** (2009, CP/AP)



redis



[Giuseppe DeCandia et al:
Dynamo: amazon's highly
available **key-value store**.
SOSP 2007]



Example Systems

■ Redis Data Types



- Redis is not a plain KV-store, but “data structure server” with persistent log (**appendfsync no/everysec/always**)
- **Key:** ASCII string (max 512MB, common key schemes: comment:1234:reply.to)
- **Values:** strings, lists, sets, sorted sets, hashes (map of string-string), etc

■ Redis APIs

- **SET/GET/DEL:** insert a key-value pair, lookup value by key, or delete by key
- **MSET/MGET:** insert or lookup multiple keys at once
- **INCRBY/DECBY:** increment/decrement counters
- Others: EXISTS, LPUSH, LPOP, LRANGE, LTRIM, LLEN, etc

■ Other systems

- Classic KV stores (AP): **Riak**, **Aerospike**, **Voldemort**, **LevelDB**, **RocksDB**, **FoundationDB**, **Memcached**
- Wide-column stores: **Google BigTable** (CP), **Apache HBase** (CP), **Apache Cassandra** (AP)



LEVELDB



Log-structured Merge Trees

[Patrick E. O'Neil, Edward Cheng, Dieter Gawlick, Elizabeth J. O'Neil: The Log-Structured Merge-Tree (LSM-Tree). *Acta Inf.* 1996]

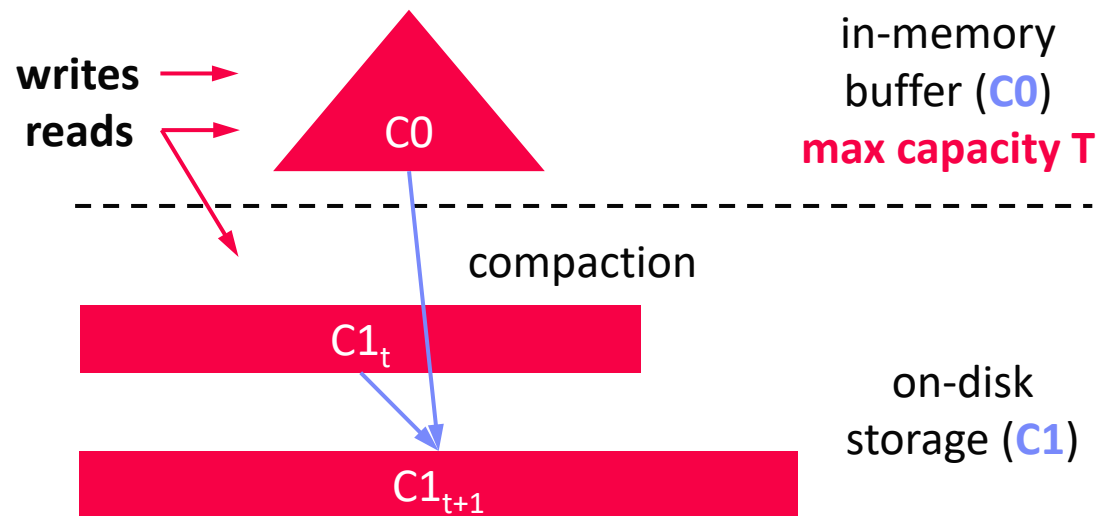


LSM Overview

- Many KV-stores rely on LSM-trees as their storage engine (e.g., **BigTable**, **DynamoDB**, **LevelDB**, **Riak**, **RocksDB**, **Cassandra**, **HBase**)
- Approach:** Buffers writes in memory, flushes data as sorted runs to storage, merges runs into larger runs of next level (compaction)

System Architecture

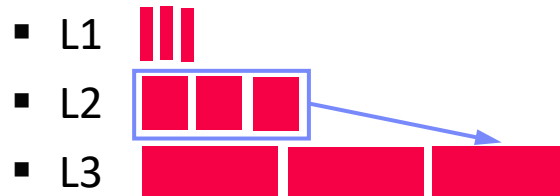
- Writes in C0
- Reads against C0 and C1 (w/ buffer for C1)
- Compaction (rolling merge): sort, merge, including **deduplication**



Log-structured Merge Trees, cont.

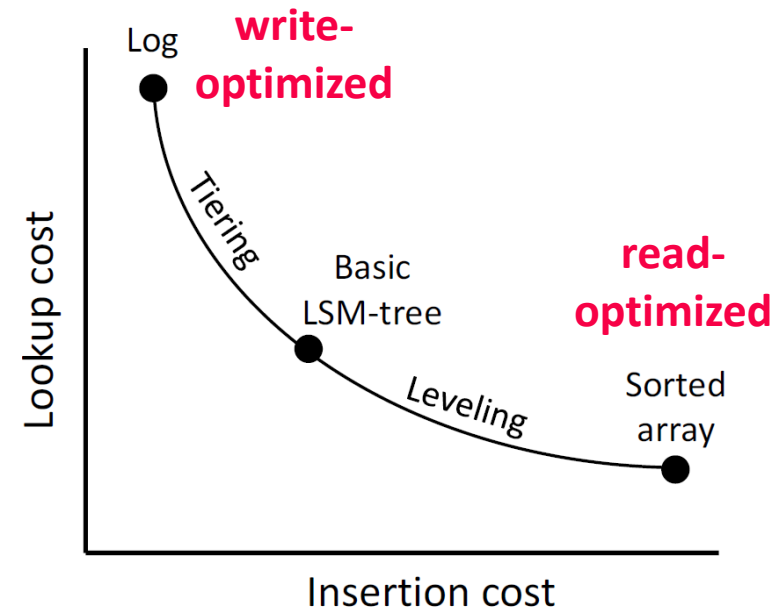
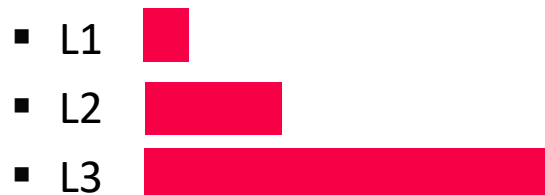
LSM Tiering

- Keep up to $T-1$ runs per level L
- Merge all runs of L_i into 1 run of L_{i+1}



LSM Leveling

- Keep 1 run per level L
- Merge run of L_i with L_{i+1}



[Niv Dayan: Log-Structured-Merge Trees, **Comp115** guest lecture, 2017]



[Stratos Idreos, Mark Callaghan: Key-Value Storage Engines (Tutorial), **SIGMOD 2020**]



Document Stores

Recap: JSON (JavaScript Object Notation)

■ JSON Data Model

- Data exchange format for **semi-structured data**
- **Not as verbose as XML** (especially for arrays)



- Popular format (e.g., Twitter)

```
{“students:”[
  {“id”: 1, “courses”: [
    {“id”: “INF.01017UF”, “name”: “DM”},
    {“id”: “706.550”, “name”: “AMLS”} ]},
  {“id”: 5, “courses”: [
    {“id”: “706.520”, “name”: “DIA”} ]},
]}
```

■ Query Languages

- **Most common: libraries** for tree traversal and data extraction
- **JSONiq**: XQuery-like query language
- **JSONPath**: XPath-like query language

JSONiq Example:

```
declare option jsoniq-version “...”;
for $x in collection(“students”)
  where $x.id lt 10
  let $c := count($x.courses)
  return {“sid”: $x.id, “count”: $c}
```

[<http://www.jsoniq.org/docs/JSONiq/html-single/index.html>]

Motivation and Terminology

■ Motivation

- Application-oriented management of **structured, semi-structured, and unstructured information** (pay-as-you-go, schema evolution)
- Scalability via parallelization on commodity HW (cloud computing)

■ System Architecture

- Collections of (**key**, **document**)
- Scalability via sharding (horizontal partitioning)
- Custom SQL-like or functional query languages

1234

```
{customer:"Jane Smith",
  items:[{name:"P1",price:49},
         {name:"P2",price:19}]}
```

1756

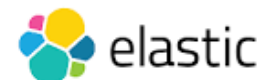
```
{customer:"John Smith", ...}
```

989

```
{customer:"Jane Smith", ...}
```

■ Example Systems

- **MongoDB** (C++, 2007, **CP**) → **RethinkDB**, **Espresso**, **Amazon DocumentDB** (Jan 2019)
- **CouchDB** (Erlang, 2005, **AP**) → **CouchBase**



Example MongoDB

[Credit: <https://api.mongodb.com/python/current>]

■ Creating a Collection

```
import pymongo as m
conn = m.MongoClient("mongodb://localhost:123/")
db = conn["dbs19"]           # database dbs19
cust = db["customers"]      # collection customers
```

■ Inserting into a Collection

```
mdict = {
    "name": "Jane Smith",
    "address": "Inffeldgasse 13, Graz"
}
id = cust.insert_one(mdict).inserted_id
# ids = cust.insert_many(mlist).inserted_ids
```

■ Querying a Collection

```
print(cust.find_one({"_id": id}))

ret = cust.find({"name": "Jane Smith"})
for x in ret:
    print(x)
```

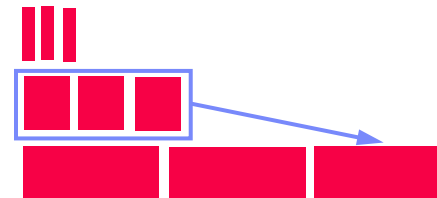
BREAK (and Test Yourself)

■ NoSQL Systems (10/100 points)

- Describe the concept and system architecture of a **key-value store**, including techniques for achieving **high write throughput**, and **scale-out** in distributed environments. [...]

■ Solution

- Key-value store system architecture [4]
- Write-throughput via **LSM** (log-structured merge tree) [3]
- Horizontal partitioning** [3] (see **07 Physical Design**)



$$R1 = \sigma_{k \leq 5}(R)$$

$$R2 = \sigma_{k > 5 \wedge k \leq 10}(R)$$

$$R3 = \sigma_{k > 10 \wedge k \leq 15}(R)$$

$$R = (R1 \cup R2) \cup R3$$

R	
<u>k</u>	v
1	Blob1
2	Blob2
4	Blob4
7	Blob7
15	Blob15
9	Blob9
14	Blob14
8	Blob8

Graph Processing

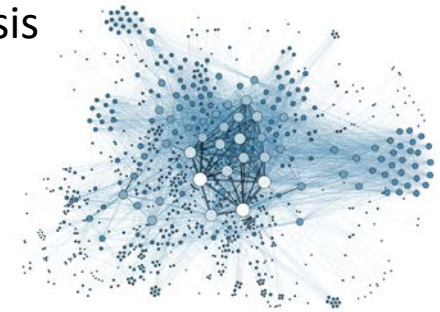
Motivation and Terminology

■ Ubiquitous Graphs

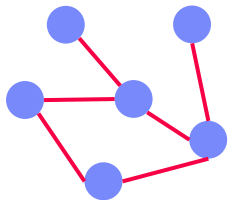
- **Domains:** social networks, open/linked data, knowledge bases, bioinformatics
- **Applications:** influencer analysis, ranking, topology analysis

■ Terminology

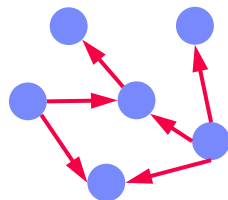
- Graph $G = (V, E)$ of vertices V (set of nodes) and edges E (set of links between nodes)
- **Different types of graphs**



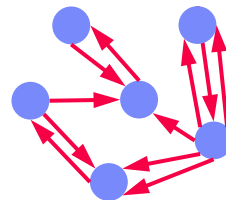
Undirected
Graph



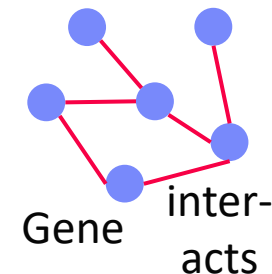
Directed
Graph



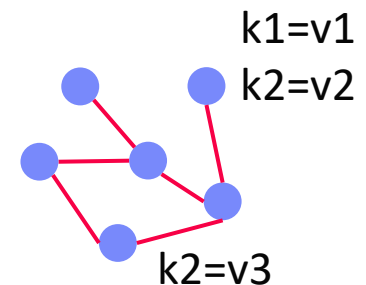
Multi
Graph



Labeled
Graph



Data/Property
Graph



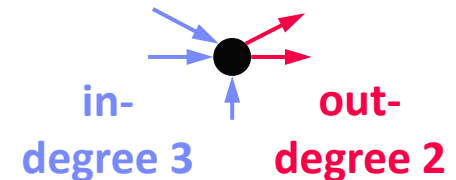
Terminology and Graph Characteristics

■ Terminology, cont.

- **Path:** Sequence of edges and vertices (**walk:** allows repeated edges/vertices)
- **Cycle:** Closed walk, i.e., a walk that starts and ends at the same vertex
- **Clique:** Subgraph of vertices where every two distinct vertices are adjacent

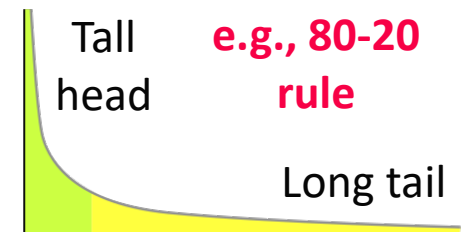
■ Metrics

- **Degree** (in/out-degree): number of incoming/outgoing edges of that vertex
- **Diameter:** Maximum distance of pairs of vertices (longest shortest-path)



■ Power Law Distribution

- Degree of most real graphs follows a power law distribution



Vertex-Centric Processing

[Grzegorz Malewicz et al: **Pregel**:
a system for large-scale graph
processing. **SIGMOD 2010**
SIGMOD 2020 Test of Time Award]



Google **Pregel**

- Name: Seven Bridges of Königsberg (Euler 1736)
- “**Think-like-a-vertex**” computation model
- Iterative processing in super steps, comm.: message passing

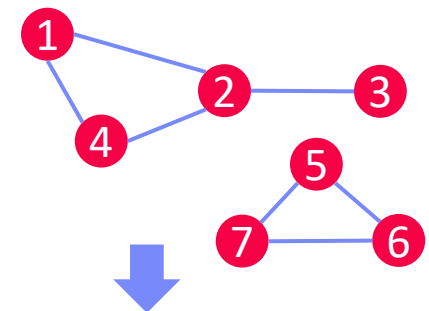


Programming Model

- Represent graph as collection of vertices w/ edge (adjacency) lists
- Implement algorithms via Vertex API
- Terminate if all vertices halted / no more msgs

```
public abstract class Vertex {
    public String getID();
    public long superstep();
    public VertexValue getValue();

    public compute(Iterator<Message> msgs);
    public sendMsgTo(String v, Message msg);
    public void voteToHalt();
}
```

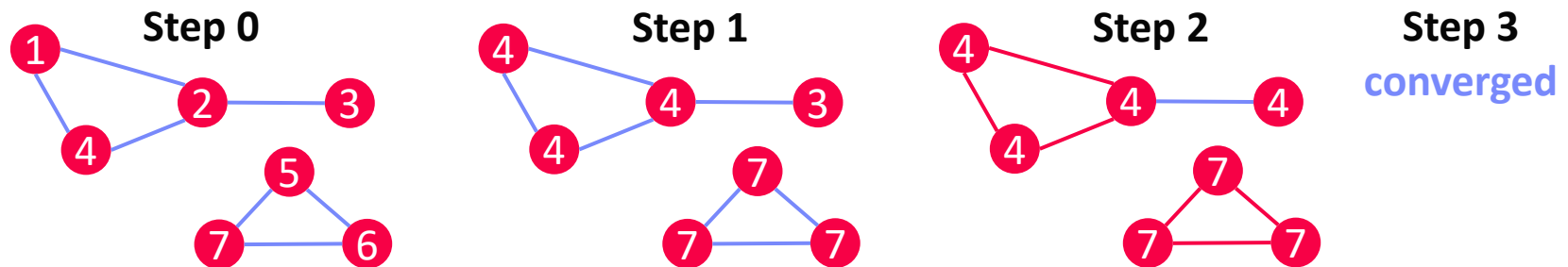


2	[1, 3, 4]	
7	[5, 6]	Worker
4	[1, 2]	1
1	[1, 2, 4]	
<hr/>		
5	[6, 7]	
3	[2]	Worker
6	[5, 7]	2

Vertex-Centric Processing, cont.

Example 1: Connected Components

- Determine connected components of a graph (subgraphs of connected nodes)
- Propagate $\max(\text{current}, \text{msgs})$ if $\neq \text{current}$ to neighbors, terminate if no msgs

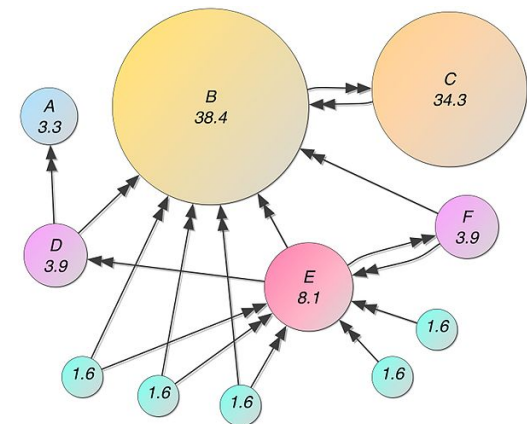


Example 2: Page Rank

- Ranking of webpages by importance / impact
- #1: **Initialize vertices** to $1/\text{numVertices}()$
- #2: **In each super step**
 - Compute current vertex value:

$$\text{value} = 0.15/\text{numVertices}() + 0.85 * \text{sum}(\text{msg})$$
 - Send to all neighbors:

$$\text{value}/\text{numOutgoingEdges}()$$



[Credit: <https://en.wikipedia.org/wiki/PageRank>]

Graph-Centric Processing

■ Motivation

- Exploit graph structure for algorithm-specific optimizations (number of network messages, scheduling overhead for super steps)
- **Large diameter / average vertex degree**

■ Programming Model

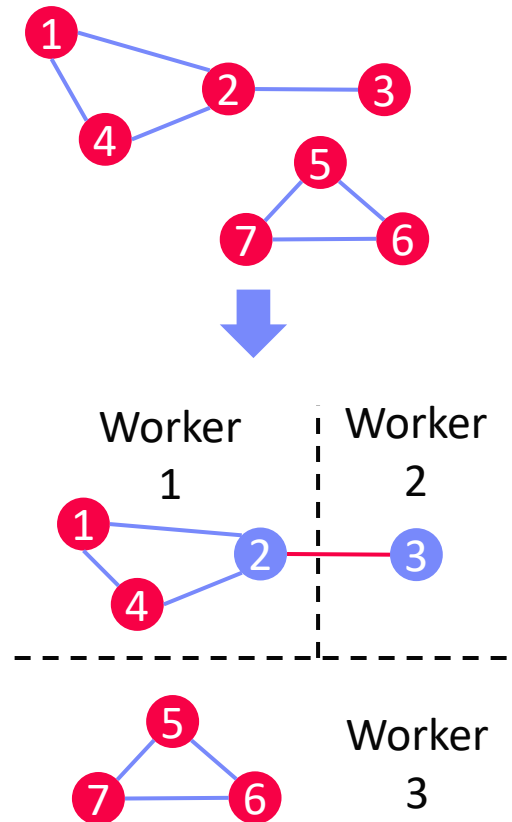
- **Partition graph into subgraphs** (block/graph)
- Implement algorithm directly against subgraphs (internal and boundary nodes)
- Exchange messages in super steps only between boundary nodes → **faster convergence**



[**Yuanyuan Tian**, Andrey Balmin, Severin Andreas Corsten, Shirish Tatikonda, John McPherson: From "Think Like a Vertex" to "Think Like a Graph". **PVLDB 2013**]



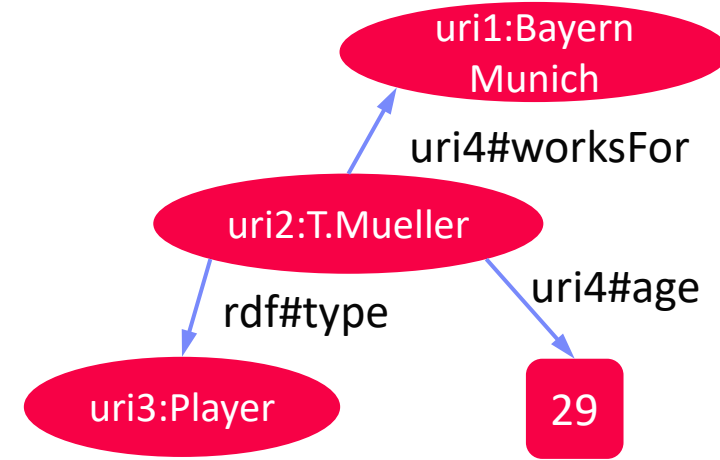
[Da Yan, James Cheng, Yi Lu, Wilfred Ng: Blogel: A Block-Centric Framework for Distributed Computation on Real-World Graphs. **PVLDB 2014**]



Resource Description Framework (RDF)

■ RDF Data

- Data and meta data description via triples
- **Triple: (subject, predicate, object)**
- Triple components can be URIs or literals
- Formats: e.g., RDF/XML, RDF/JSON, Turtle
- RDF graph is a directed, labeled multigraph



■ Querying RDF Data

- SPARQL (SPARQL Protocol And RDF Query Language)
- Subgraph matching

```
SELECT ?person
WHERE {
    ?person rdf:type uri3:Player ;
    uri4:worksFor uri1:"Bayern Munich" .
}
```

■ Selected Example Systems

AWS Amazon Neptune



AllegroGraph

*Sparksee



Excursus: Example Systems



■ Understanding Use in Practice

- Types of graphs user have
- Graph computations run
- Types of graph systems used



[Siddhartha Sahu, Amine Mhedhbi, Semih Salihoglu, Jimmy Lin, M. Tamer Özsu: The Ubiquity of Large Graphs and Surprising Challenges of Graph Processing. **PVLDB 2017**]

Technology	Software	# Users	
Graph Database System	ArrangoDB [3]	40	233
	Caley [8]	14	
	DGraph [14]	33	
	JanusGraph [35]	32	
	Neo4j [48]	69	
	OrientDB [53]	45	
RDF Engine	Apache Jena [38]	87	115
	Sparksee [64]	5	
	Virtuoso [67]	23	
Distributed Graph Processing Engine	Apache Flink (Gelly) [17]	24	39
	Apache Giraph [21]	8	
	Apache Spark (GraphX) [27]	7	
Query Language	Gremlin [28]	82	82
Graph Library	Graph for Scala [22]	4	97
	GraphStream [24]	8	
	Graphtool [25]	28	
	NetworkKit [50]	10	
	NetworkX [51]	27	
	SNAP [62]	20	
Graph Visualization	Cytoscape [13]	93	116
	Elasticsearch (X-Pack Graph) [16]	23	
Graph Representation	Conceptual Graphs [11]	6	6

■ Summary of State of the Art Runtime Techniques

[Da Yan, Yingyi Bu, Yuanyuan Tian, Amol Deshpande, James Cheng: Big Graph Analytics Systems. **SIGMOD 2016**]



Excursus: Future Graph Processing Systems

Community Perspective

- 2019 Dagstuhl Seminar on Big Graph Processing Systems
- Opportunities and challenges:
Abstractions, Ecosystems, and Performance

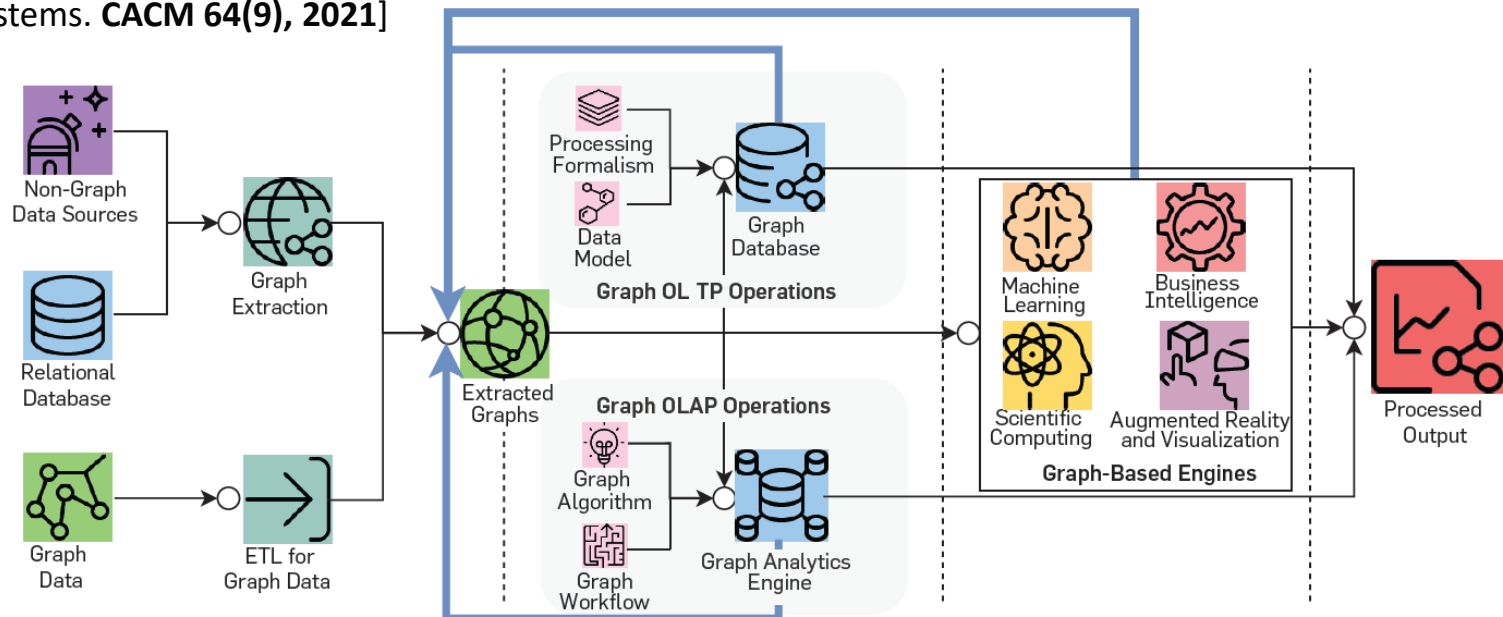


[Credit:

https://commons.wikimedia.org/wiki/File:Dagstuhl_DSC02285.jpg]



[Sherif Sakr et al: The future is big graphs: a community view on graph processing systems. **CACM 64(9), 2021**]



Time Series Databases

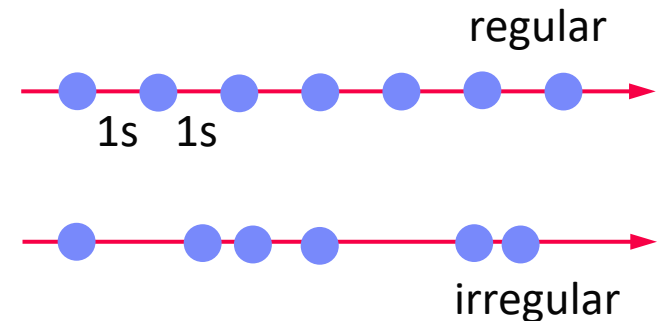
Motivation and Terminology

Ubiquitous Time Series

- **Domains:** Internet-of-Things (IoT), sensor networks, smart production/planet, telemetry, stock trading, server/application metrics, event/log streams
- **Applications:** monitoring, anomaly detection, time series forecasting
- Dedicated storage and analysis techniques → Specialized systems

Terminology

- Time series X is a sequence of data points x_i for a specific measurement identity (e.g., sensor) and time granularity
- **Regular** (equidistant) time series (x_i)
vs **irregular** time series (t_i, x_i)



Example InfluxDB

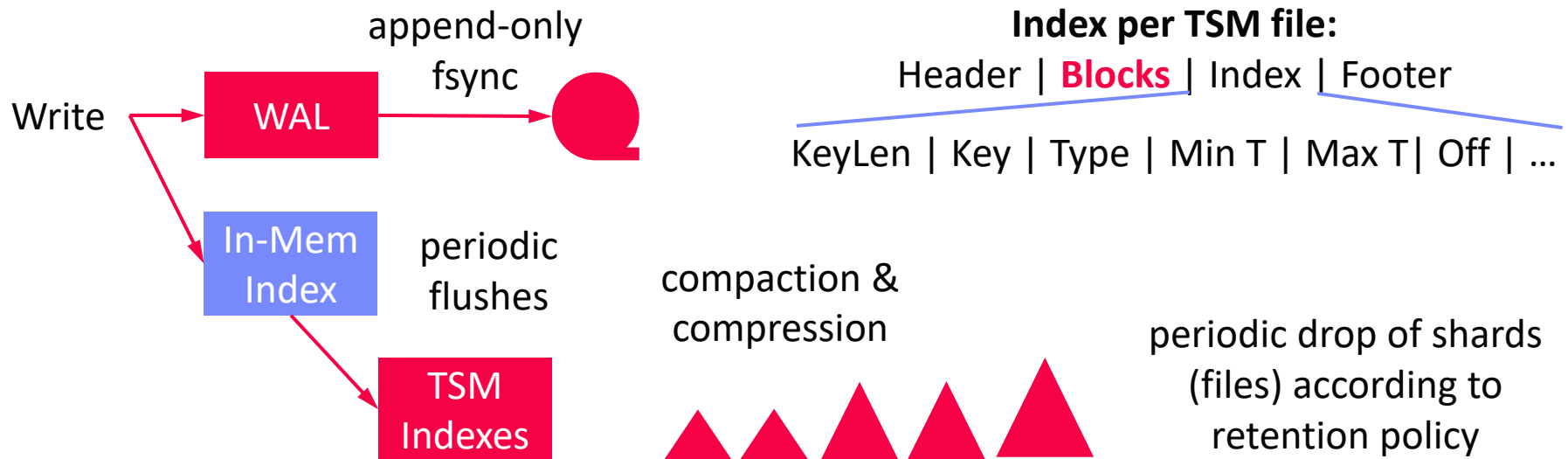


[Paul Dix: InfluxDB
Storage Engine Internals,
CMU Seminar, 09/2017]

- Input Data `cpu,region=west,host=A` **Measurement**
`user=85,sys=2,idle=10` **Fields (values)**
`1443782126` **Time**
`region=west,host=A` **Tags**

System Architecture

- Written in Go, originally **key-value store**, now **dedicated storage engine**
- Time Structured Merge Tree (TSM)**, similar to LSM
- Organized in shards, TSM indexes and inverted index for reads



Example InfluxDB, cont.

■ Compression (of blocks)

- **Compress up to 1000 values per block** (Type | Len | Timestamps | Values)
- **Timestamps:** Delta + Run-length encoding for regular time series; Simple8B or uncompressed for irregular
- **Values:** double delta for FP64, bits for Bool, double delta + zig zag for INT64, Snappy for strings

■ Query Processing

- SQL-like and functional APIs for filtering (e.g., range) and aggregation
- Inverted indexes

```
SELECT percentile(90, user)
FROM cpu WHERE time>now()-12h
AND "region"='west'
GROUP BY time(10m), host
```

Measurement to fields:

cpu → [user,sys,idle]

host → [A, B]

Region → [west, east]

Posting lists:

cpu → [1,2,3,4,5,6]

host=A → [1,2,3]

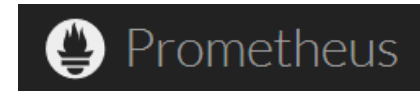
host=B → [4,5,6]

region=west → [1,2,3]

Other Systems

■ Prometheus

- Metrics, high-dim data model, sharding and federation custom storage and query engine, implemented in Go



■ OpenTSDB

- TSDB on top of HBase or Google BigTable, Hadoop



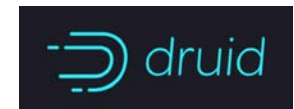
■ TimescaleDB

- TSDB on top of PostgreSQL, standard SQL and reliability



■ Druid

- Column-oriented storage for time series, OLAP, and search



■ IBM Event Store

- HTAP system for high data ingest rates, and data-parallel analytics via Spark
- Shard-local logs → groomed data

[Ronald Barber et al: Evolving Databases for New-Gen Big Data Applications. **CIDR 2017**]

[Christian Garcia-Arellano et al: **Db2 Event Store**: A Purpose-Built IoT Database Engine. **PVLDB 13(12) 2020**]



Exercise 4: Large-Scale Data Analysis

Published: Dec 11

Deadline: Jan 18

Entire Exercise is Extra Credit

Task 4.1 Apache Spark Setup

3/25
points

■ #1 Pick your Spark Language Binding

- Java, Scala, Python

■ #2 Install Dependencies

- Java: Maven
`spark-core, spark-sql`
- Python:
`pip install pyspark`

```
<dependency>
  <groupId>org.apache.spark</groupId>
  <artifactId>spark-core_2.11</artifactId>
  <version>2.4.7</version>
</dependency>
<dependency>
  <groupId>org.apache.spark</groupId>
  <artifactId>spark-sql_2.11</artifactId>
  <version>2.4.7</version>
</dependency>
```

■ (#3 Win Environment)

- Download <https://github.com/steveloughran/winutils/tree/master/hadoop-2.7.1/bin/winutils.exe> (or <https://github.com/cdarlint/winutils>)
- Create environment variable HADOOP_HOME=“<some-path>/hadoop”

Task 4.2 Query Processing via Spark RDDs

11/25
points

■ #1 Spark Context Creation

- Create a spark context sc w/ local master (local[*])

■ #2 Implement Q09 via RDD Operations

- Implement Q09 in self-contained executeQ09RDD()
- All reads should use `sc.textFile(fname)`
- RDD operations only → stdout

<https://spark.apache.org/docs/latest/rdd-programming-guide.html>

- **Note:** Query will be shared by Dec 28

Task 4.3 Query Processing via Spark SQL

5/25
points

■ #1 Spark Session Creation

- Create a spark session via a spark session builder and w/ local master (`local[*]`)

→ SQL processing of high importance in modern data management

■ #2 Implement Q09 via Dataset Operations

- Implement Q09 self-contained in `executeQ09Dataset()`
- All reads should use `sc.read().format("csv")`
- SQL or Dataset operations only → `out07.json`

- **WebUI** INFO Utils: Successfully started service 'SparkUI' on port 4040.
INFO SparkUI: Bound SparkUI to [...] <http://192.168.108.220:4040>

Task 4.4 Graph Processing

6/25
points

Input Co-author graph

AuthPapersCOO.csv

(coordinate format)

AuthPapersCSR.csv

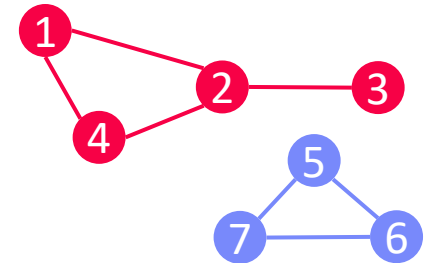
(compressed sparse row)

```
1 author,co-author
2 1001634,70215
3 1001634,519925
```

```
1 author,co-authors
2 1001634,70215:519925:1444319:2383440
3 1243968,76416:323847:407298:688292:918500:1198961:1231227:1256611:1377989
```

#1 Compute Connected Components

- Leverage Spark to compute assignment of vertices to components
- Write output to text file, print #components to stdout
- APIs up to you (e.g., Spark RDDs, Spark SQL, Spark GraphX)



- Example
Apache
SystemDS

```
37 # initialize state with vertex ids
38 c = seq(1,nrow(G));
39 diff = Inf;
40 iter = 1;
41
42 # iterative computation of connected components
43 while( diff > 0 & (maxi==0 | iter<=maxi) ) {
44     u = max(rowMaxs(G * t(c)), c);
45     diff = sum(u != c)
46     c = u; # update assignment
```


Conclusions and Q&A

- **Summary 10 NoSQL Systems**
 - Consistency and Data Models
 - Key-Value and Document Stores
 - Graph and Time Series Databases

- **Next Lectures (Part B: Modern Data Management)**
 - 11 Distributed Storage and Data Analysis [Jan 10]
 - 12 Data stream processing systems, Q&A [Jan 17]