



# Data Integration and Analysis 07 Data Provenane and Blockchain

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# Announcements/Org

#### #1 Video Recording

- Link in TUbe & TeachCenter (lectures will be public)
- Optional attendance (independent of COVID)
- Virtual lectures (recorded) until end of year https://tugraz.webex.com/meet/m.boehm









# Agenda

- Motivation and Terminology
- Data Provenance
- Blockchain Fundamentals





# Data Provenance Motivation and Terminology





## Excursus: FAIR Data Principles



#### #1 Findable

- Metadata and data have globally unique persistent identifiers
- Data describes w/ rich meta data; registered/indexes and searchable

#### #2 Accessible

- Metadata and data retrievable via open, free and universal comm protocols
- Metadata accessible even when data no longer available

### #3 Interoperable

- Metadata and data use a formal, accessible, and broadly applicable format
- Metadata and data use FAIR vocabularies and qualified references

#### #4 Reusable

- Metadata and data described with plurality of accurate and relevant attributes
- Clear license, associated with provenance, meets community standards





# Terminology

#### **Data Provenance**

Track and understand data origins and transformations of data (where?, when?, who?, why?, how?)



- Contains meta data, context, and modifications (transform, enrichment)
- Synonyms: data lineage, data pedigree
- **Data Catalogs** (curation/governance)
  - Directory of datasets including data provenance (meta data, artifacts)
  - Raw/original, curated datasets, derived data products

#### **Blockchain**

Data structure logging transactions in verifiable and permanent way





# **Applications and Goals**

### a) High-Level Goals

- #1 Versioning and Reproducibility (analogy experiments)
- #2 Explainability, Interpretability, Verification

### b) Low-Level Goals

- #3 Full and Partial Reuse of Intermediates
- #4 Incremental Maintenance of MatViews, Models, etc
- #5 Tape/log of Executed Operations → Auto Differentiation
- #6 Recomputation for Caching / Fault Tolerance
- #7 Debugging via Lineage Query Processing









# Data Provenance





## Overview Data Provenance

#### Def Data Provenance

Information about the origin and creation process of data

### Example

Debugging suspicious query results

SELECT Customer, sum(0.Quantity\*P.Price)
FROM Orders 0, Products P
WHERE 0.PID = P.PID
GROUP BY Customer

Customer	Sum
Α	7620
В	120
С	130
D	75

OID	Customer	Date	Quantity	PID
1	Α	2019-06-22	3	2
2	В	2019-06-22	1	3
3	Α	2019-06-22	101	4
4	С	2019-06-23	2	2
5	D	2019-06-23	1	4
6	С	2019-06-23	1	1

PID	Product	Price
1	Χ	100
2	Υ	15
4	Z	75
3	W	120





## Overview Data Provenance, cont.

#### An Abstract View

■ Data: schema, structure → data items

[Boris Glavic: CS595 Data Provenance – Introduction to Data Provenance, Illinois Institute of Technology, 2012]

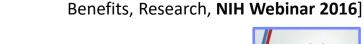
- Data composition (granularity): attribute, tuple, relation
- Transformation: consumes inputs, produces outputs



Additional: env context (OS, libraries, env variables, state), users

## Goal: Tracing of Derived Results

- Inputs and parameters
- Steps involved in creating the result
- → Store and query data & provenance
- General Data Protection Regulation (GDPR)?



[Zachary G. Ives: Data Provenance: Challenges,



- 1. Read file1
- 2. Sort rows
- 3. Compute median
- 4. Write to file2







## Classification of Data Provenance

#### Overview

- Base query Q(D) = O with database D =  $\{R_1, ..., R_n\}$
- Forward lineage query: L<sub>f</sub>(R<sub>i</sub>", O') from subset of input relation to output
- Backward lineage query: L<sub>b</sub>(O', R<sub>i</sub>) from subset of outputs to base tables

### #1 Lazy Lineage Query Evaluation

- Rewrite (invert) lineage queries as relational queries over input relations
- No runtime overhead but slow lineage query processing

### #2 Eager Lineage Query Evaluation

- Materialize annotations (data/transforms) during base query evaluation
- Runtime overhead but fast lineage query processing
- Lineage capture: Logical (relational)
   vs physical (instrumented physical ops)

[Fotis Psallidas, Eugene Wu: Smoke: Fine-grained Lineage at Interactive Speed. **PVLDB 2018**]







# Why-Provenance

[Boris Glavic: CS595 Data Provenance – Provenance Models and Systems, Illinois Institute of Technology, 2012]



### Overview Why

- Goal: Which input tuples contributed to an output tuple t in query Q
- Representation: Set of witnesses w for tuple t (set semantics!)
  - $w \subseteq I$  (subset of all tuples in instance I)
  - $t \in Q(w)$  (tuple in result of query over w)

## Example Witnesses

SELECT Customer, Product
FROM Orders O, Products P
WHERE O.PID=P.PID

	Customer	Date	PID
<b>o1</b>	Α	2019-06-22	2
<b>o2</b>	В	2019-06-22	3
о3	А	2019-06-22	2

	PID	Product
p1	1	X
<b>p2</b>	2	Υ
р3	4	Z
p4	3	W

Witnesses for t1:

$$w1 = \{o1,p2\}, w2 = \{o3,p2\},$$
  
 $w3 = \{o1,p2,p3\}, ..., wn = I$ 

Minimal witnesses for t1: w1 = {o1,p2}, w2 = {o3,p2}

	Customer	Product
t1	Α	Υ
t2	В	W

**Others:** Where/How Provenance





## How-Provenance

[Boris Glavic: CS595 Data Provenance – Provenance Models and Systems, Illinois Institute of Technology, 2012]



#### Overview

- Model how tuples where combined in the computation
- Alternative use: need one of the tuples (e.g., union/projection)
- Conjunctive use: need all tuples together (e.g., join)

### Provenance Polynomials

• Semiring annotations to model provenance ( $\mathbb{N}[I]$ , +,×, 0,1)

### Examples

$$q = \pi_a(R)$$

	а	b
r1	1	2
<b>r2</b>	1	3



$$r1 + r2$$

**Provenance** 

$$q = \pi_b(R \bowtie S)$$

,	а	b		С	а
r1	1	Р	s1	S	1
<b>r2</b>	2	G	<b>s2</b>	S	2
	3	M	<b>s3</b>	W	2







# Why Not?-Provenance

[Adriane Chapman, H. V. Jagadish: Why not? SIGMOD 2009]



#### Overview

- Why are items not in the results
- Example Problem:

"Window-display-books < \$20"

→ (Euripides, Medea).

>= 20\$?

→ Why not (Hrotsvit, Basilius)?

Bug in the query / system?

Not in book store?

Author	Title	Price	Publisher
	Epic of Gilgamesh	\$150	Hesperus
Euripides	Medea	\$16	Free Press
Homer	Iliad	\$18	Penguin
Homer	Odyssey	\$49	Vintage
Hrotsvit	Basilius	\$20	Harper
Longfellow	Wreck of the Hesperus	\$89	Penguin
Shakespeare	Coriolanus	\$70	Penguin
Sophocles	Antigone	\$48	Free Press
Virgil	Aeneid	\$92	Vintage

## Evaluation Strategies

- Given a user question (why no tuple satisfies predicate S), dataset D, result set R, and query Q, leverage why lineage
- #1 Bottom-Up: from leafs in topological order to find last op eliminating  $d \in S$
- #2 Top-Down: from result top down to find last op, requires stored lineage





# **Apache Atlas**

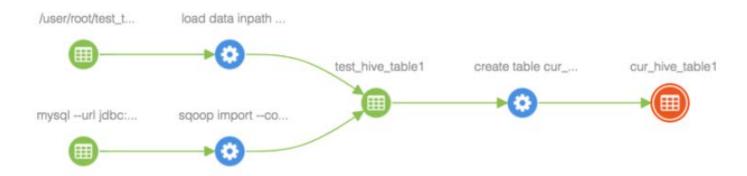


#### Apache Atlas Overview

- Metadata management and governance capabilities
- Build catalog (data classification, cross-component lineage)

## Example

- Configure Atlas hooks w/ Hadoop components
- Automatic tracking of lineage and side effects



[https://www.cloudera.com/tutorials/cross-component-lineage-with-apache-atlas-across-apache-sqoop-hive-kafka-storm/.html]





# Provenance for ML Pipelines (fine-grained)

#### DEX: Dataset Versioning

- Versioning of datasets, stored with delta encoding
- Checkout, intersection, union queries over deltas
- Query optimization for finding efficient plans

[Amit Chavan, Amol Deshpande: DEX: Query Execution in a Deltabased Storage System. SIGMOD 2017]



### MISTIQUE: Intermediates of ML Pipelines

- Capturing, storage, querying of intermediates
- Lossy deduplication and compression
- Adaptive querying/materialization for finding efficient plans

### Linear Algebra Provenance

- Provenance propagation by decomposition
- Annotate parts w/ provenance polynomials (identifiers of contributing inputs + impact)

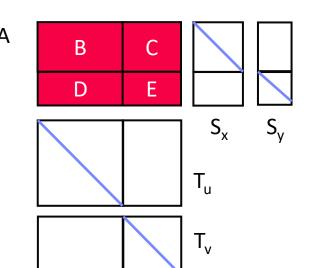
$$A = S_x B T_u + S_x C T_v + S_y D T_u + S_y E T_v$$



[Zhepeng Yan, Val Tannen, Zachary G. Ives: Fine-grained Provenance for Linear Algebra Operators. **TaPP 2016**]

[Manasi Vartak et al: MISTIQUE: A System to Store and Query Model Intermediates for Model Diagnosis. **SIGMOD 2018**]







## Provenance for ML Pipelines (coarse-grained)

#### MLflow

- Programmatic API for tracking parameters, experiments, and results
- autolog() for specific params

```
import mlflow
mlflow.log_param("num_dimensions", 8)
mlflow.log_param("regularization", 0.1)
mlflow.log_metric("accuracy", 0.1)
mlflow.log_artifact("roc.png")
```

## Flor (on Ground)

- DSL embedded in python for managing the workflow development phase of the ML lifecycle
- DAGs of Actions, Artifacts, and Literals
- Data context generated by activities in Ground

[Credit: <a href="https://rise.cs.berkeley.edu/">https://rise.cs.berkeley.edu/</a> projects/jarvis/ ]

[Credit: https://databricks.com/

[Joseph M. Hellerstein et al: Ground: A Data Context Service. **CIDR 2017**]



## Dataset Relationship Management

- Reuse, reveal, revise, retarget, reward
- Code-to-data relationships (data provenance)
- Data-to-code relationships (potential transforms)

[Zachary G. Ives, Yi Zhang, Soonbo Han, Nan Zheng,: Dataset Relationship Management. **CIDR 2019**]









# Provenance for ML Pipelines (coarse-grained), cont.

#### HELIX

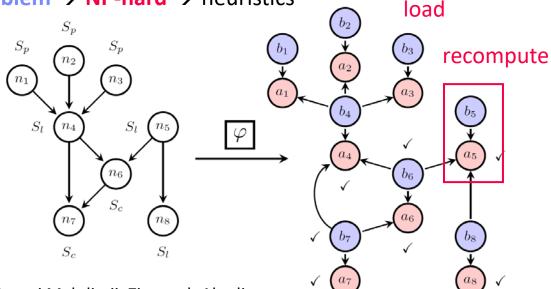
- Goal: focus on iterative development
   w/ small modifications (trial & error)
- Caching, reuse, and recomputation

[Doris Xin, Stephen Macke, Litian Ma, Jialin Liu, Shuchen Song, Aditya G. Parameswaran: Helix: Holistic Optimization for Accelerating Iterative Machine Learning. **PVLDB 2018**]



■ Reuse as Max-Flow problem → NP-hard → heuristics

 Materialization to disk for future reuse



CollaborativeOptimizer



[Behrouz Derakhshan, Alireza Rezaei Mahdiraji, Ziawasch Abedjan, Tilmann Rabl, Volker Markl: Optimizing Machine Learning Workloads in Collaborative Environments. **SIGMOD 2020**]





# Lineage Tracing & Reuse in SystemDS





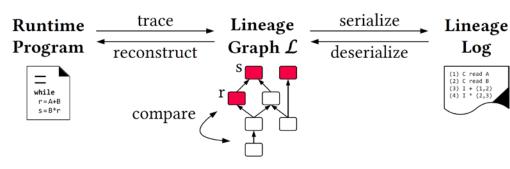
#### Problem

- Exploratory data science (data preprocessing, model configurations)
- Reproducibility and explainability of trained models (data, parameters, prep)
- → Lineage/Provenance as Key Enabling Technique:

  Model versioning, reuse of intermediates, incremental maintenance, auto differentiation, and debugging (query processing over lineage)

### Efficient Lineage Tracing

- Tracing of inputs, literals, and non-determinism
- Trace lineage of logical operations
- Deduplication for loops/functions
- Program/output reconstruction



[Arnab Phani, Benjamin Rath, Matthias Boehm: LIMA: Fine-grained Lineage Tracing and Reuse in Machine Learning Systems, SIGMOD 2021]







# Lineage Tracing & Reuse in SystemDS



- Multi-level, Lineage-based Reuse
  - Lineage trace uniquely identifies intermediates
  - Reuse intermediates at function / block / operation level
- Full Reuse of Intermediates
  - Before executing instruction, probe output lineage in cache Map<Lineage, MatrixBlock>
  - Cost-based/heuristic caching and eviction decisions (compiler-assisted)
- Partial Reuse of Intermediates
  - Problem: Often partial result overlap
  - Reuse partial results via dedicated rewrites (compensation plans)
  - Example: steplm

```
X
```

t(X)

```
for( i in 1:numModels )
    R[,i] = lm(X, y, lambda[i,], ...)

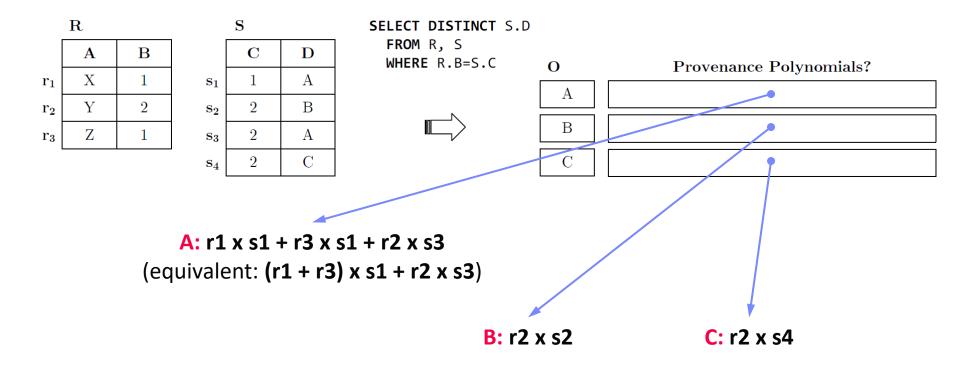
m_lmDS = function(...) {
    l = matrix(reg,ncol(X),1)
    A = t(X) %*% X + diag(1)
    b = t(X) %*% y
    beta = solve(A, b) ...}
```

```
m_steplm = function(...) {
  while( continue ) {
    parfor( i in 1:n ) {
       if( !fixed[1,i] ) {
          Xi = cbind(Xg, X[,i])
          B[,i] = lm(Xi, y, ...)
       }
       # add best to Xg
       # (AIC)
    }
}
```



# BREAK (and Test Yourself)

• Given below tables R and S (with tuples r<sub>i</sub> and s<sub>i</sub>), query Q and results O, specify the provenance polynomials for every tuple in O. (3 points)





# **Blockchain Fundamentals**





# Recap: Database (Transaction) Log

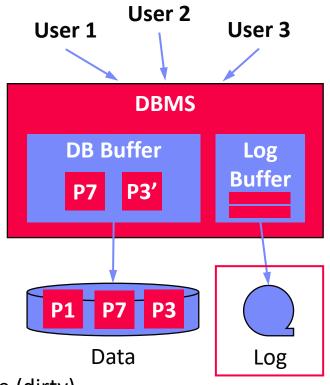
#### Database Architecture

- Page-oriented storage on disk and in memory (DB buffer)
- Dedicated eviction algorithms
- Modified in-memory pages marked as dirty, flushed by cleaner thread
- Log: append-only TX changes
- Data/log often placed on different devices and periodically archived (backup + truncate)

## Write-Ahead Logging (WAL)

- The log records representing changes to some (dirty)
   data page must be on stable storage before the data page (UNDO atomicity)
- Force-log on commit or full buffer (REDO durability)
- Recovery: forward (REDO) and backward (UNDO) processing
- Log sequence number (LSN)

[C. Mohan, Donald J. Haderle, Bruce G. Lindsay, Hamid Pirahesh, Peter M. Schwarz: ARIES: A Transaction Recovery Method Supporting Fine-Granularity Locking and Partial Rollbacks Using Write-Ahead Logging. **TODS 1992**]







## Bitcoin and Blockchain

[Satoshi Nakamoto: Bitcoin: A Peer-to-Peer Electronic Cash System, White paper 2008]



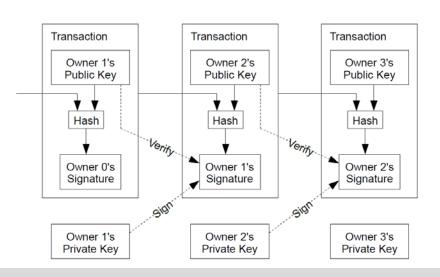
#### Motivation

- Peer-to-peer (decentralized, anonymous) electronic cash/payments
- Non-reversible transactions w/o need for trusted third party

1.16 Megabytes	303,921 Transactions	11,304,890 Bytes	.blockchain. com/charts]
1.29	310,424	19,920,773	
	Megabytes	1.29 310,424	1.29 310,424 19,920,773

#### Transaction Overview

- Electronic coin defined as chain of digital signatures
- Transfer by signing hash of previous
   TX and public key of next owner
- Double-spending problem (without global verification)







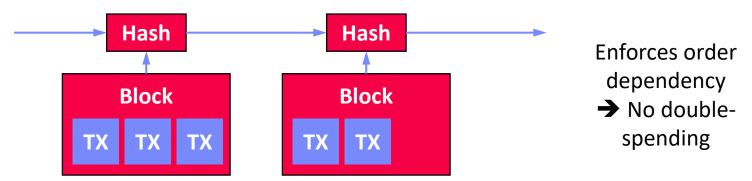
## Blockchain Data Structure

[Satoshi Nakamoto: Bitcoin: A Peer-to-Peer Electronic Cash System, White paper 2008]



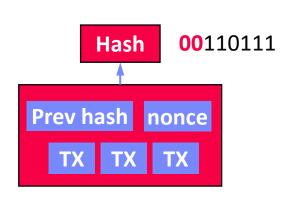
#### Timestamp Server

Decentralized timestamp server: chain of hashes 
 public ledger



#### Proof-of-Work

- Scanning for value (nonce) which SHA-256 hash begins with a number of zero bits
   → exponential in number of zeros
- # zero bits determined by MA of avg blocks/hour
- Hard to recompute for chain, easy to check
- Majority decision: CPU time, longest chain



Merkel tree (hash tree)





# Blockchain Data Structure, cont.

## Bitcoin Mining

- HW: from CPU to GPUs/FPGAs/ASICs (10-70 TH/s @ 2-3KW)
- Usually mining pools → "mining cartels"







@Malaysia

#### Hash Rate of Bitcoin Network

Blockchain.com

2019

~10 min per block (144 blocks per day)



2020

Nov 12, 2021: ~160EH

[https://www.blockchain.com/en/charts/hash-rate?daysAverageString=7&timespan=180days, Nov 12 2021]

2021





## **Blockchain Communication**

[Satoshi Nakamoto: Bitcoin: A Peer-to-Peer Electronic Cash System, White paper 2008]



## **Networking Protocol**

- New TXs are broadcast to all nodes
- Fach node collects new TXs into a block
- Each node works on finding **proof-of-work** for its block
  - → Incentive: 1st TX in block new coin (halves every 210k blocks) for the block creator + TX fees

- 2008: 50BTC
- 2012: 25BTC
- 2016: 12.5BTC 2020: 6.25BTC
- When a node finds a proof-of-work, broadcast the block to all nodes
- Nodes accept the block if all TXs are valid (double spending)
- Nodes express acceptance by working on next block in the chain, using the hash of the accepted block as the previous hash

#### Fault Tolerance

- TX broadcasts: no need to reach all but many → next block contains it
- **Block broadcast:** no need to reach all  $\rightarrow$  next block references it





# **Smart Contracts (Ethereum)**

#### Motivation

- Problem: Bitcoin TXs for transferring X \$BTC from A to B (exchange as assets)
- Goals: voting, auctions, games, bets, legal agreements (notary)

#### Ethereum

- Decentralized platform that allows creation, management, and execution of smart contracts
- Ether cryptocurrency, block mining rate: seconds (5 ETH/block)

# 1 • <contract> 2 3 4 5 6 7 19 20 • </contract>

[Credit: Shana Hutchison]

#### Smart Contract

- Store smart contract (turing-complete programs) on the blockchain
- On transfer/trigger: run smart contract (in ) in Ethereum Virtual Machine
- Language: Serpent/Solidity (deterministic, w/ control flow and function calls)
  - → Problem: while(true) → EVM gas and fees (start gas, gas price)
- Immutability guarantees persistence





## Consensus Mechanisms

"means of showing that one invested a non-trivial amount of effort related to some statement"

#### Proof of Work (PoW)

- Validation by performing work, existence of HW resources
- High HW cost of attacks
- Wasted work, resources, energy (only first block, no real outcome, e-waste)

## Proof of Stake (PoS)

- Validation by stake-weighted random node selection
- Intrinsic coin cost, less HW resources/energy
- Untested attack mitigation?

#### Ethereum 2.0

→ PoS/sharding over time

## Proof of Space/Capacity

 Upfront creation of "plot files", store nonces+hashes, find solutions, occasional validation

- HW costs of attacks, use of unused space
- Moderate adoption

[https://www.chia.net/]

[Stefan Dziembowski, Sebastian Faust, Vladimir Kolmogorov, Krzysztof Pietrzak: Proofs of Space. IACR Cryptol. 2013]







# Permissioned/Private Blockchains



#### Private Setup

Business Networks connect businesses

Participants with Identity

Assets flow over business networks

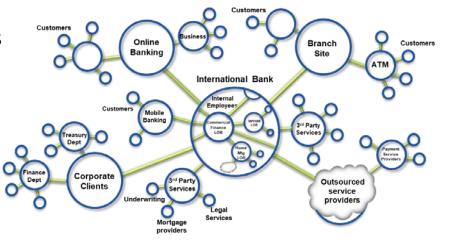
 Transactions describe exchange or change of states of assets

- Smart contracts underpin transactions
- Blockchain as shared, replicated, permissioned ledger (TX log):

consensus, provenance, immutability

[C. Mohan: State of Permissionless and Permissioned Blockchains: Myths and

Reality, **2019**]



## Hyperledger Fabric (<a href="https://github.com/hyperledger/">https://github.com/hyperledger/</a>)

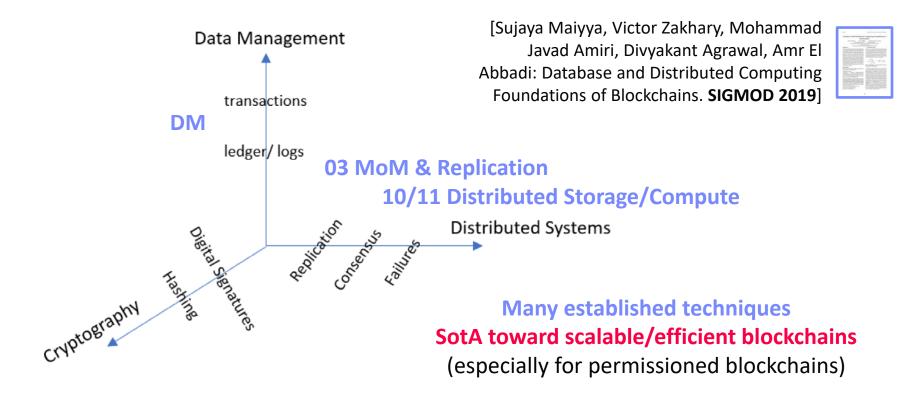


- IBM, Oracle, Baidu, Amazon, Alibaba, Microsoft, JD, SAP, Huawei, Tencent
- Blockchain-as-a-Service (BaaS) offerings: distributed ledger, libs, tools





## **Discussion Blockchain**



→ Recommendation: Investigate business requirements/context, decide on technical properties and acceptable trade-offs





# Summary and Q&A

### Summary

- Motivation and Terminology
- Data Provenance
- Blockchain Fundamentals

## Next Lectures (Part B)

- 08 Cloud Computing Foundations [Nov 27]
- 09 Cloud Resource Management and Scheduling [Dec 04]
- 10 Distributed Data Storage [Dec 11]

